Coarse-grained and fine-grained locking Niklas Fors 2013-12-05

Slides borrowed from: http://cs.brown.edu/courses/cs176course_information.shtml



Art of Multiprocessor Programming

Topics discussed

- Coarse-grained locking

 One lock
- Fine-grained locking
 - More than one lock



Abstract Data Types

• Concrete representation:

- Abstract Type:
 - -{a, b}



Abstract Data Types

 Meaning of rep given by abstraction map

$-S(\square \rightarrow a \rightarrow b \rightarrow \square) = \{a, b\}$















• Easy, same as synchronized methods – "One lock to rule them all ..."



- Easy, same as synchronized methods – "One lock to rule them all ..."
- Simple, clearly correct
 - Deserves respect!
- Works poorly with contention



Fine-grained Locking

- Requires careful thought
- Split object into pieces
 - Each piece has own lock
 - Methods that work on disjoint pieces need not exclude each other











12













Removing a Node





Removing a Node



Removing a Node





Removing a Node a С O Why lock victim node? remove(b) Ο 0



Art of Multiprocessor Programming





















Uh, Oh Bad news, c not removed a С O remove(c) remove(b) 0 0 Art of Multiprocessor Programming 32








Hand-Over-Hand Again



Hand-Over-Hand Again





Removing a Node C b d a remove(c) remove(b) Ο 0 Art of Multiprocessor Programming 40































```
public boolean remove(T item) {
  int key = item.hashCode();
  Node pred, curr;
  try {
    ...
  } finally {
    curr.unlock();
    pred.unlock();
  }}
```





Key used to order node





Predecessor and current nodes











```
try {
  pred = head;
  pred.lock();
  curr = pred.next;
  curr.lock();
  ...
} finally { ... }
```















```
while (curr.key <= key) {</pre>
  if (item == curr.item) {
   pred.next = curr.next;
   return true;
  pred.unlock();
  pred = curr;
  curr = curr.next;
  curr.lock();
```



return false;
































Why does this work?

- To remove node e
 - Must lock e
 - Must lock e's predecessor
- Therefore, if you lock a node
 - It can't be removed
 - And neither can its successor



Adding Nodes

- To add node e
 - Must lock predecessor
 - Must lock successor
- Neither can be deleted



Same Abstraction Map

- S(head) =
 - { x | there exists a such that
 - a reachable from head and
 - a.item = x



Rep Invariant

- Easy to check that
 - tail always reachable from head
 - Nodes sorted, no duplicates



Drawbacks

- Better than coarse-grained lock
 Threads can traverse in parallel
- Still not ideal
 - Long chain of acquire/release
 - Inefficient





This work is licensed under a

Creative Commons Attribution-ShareAlike 2.5 License.

- You are free:
 - to Share to copy, distribute and transmit the work
 - **to Remix** to adapt the work
- Under the following conditions:
 - Attribution. You must attribute the work to "The Art of Multiprocessor Programming" (but not in any way that suggests that the authors endorse you or your use of the work).
 - Share Alike. If you alter, transform, or build upon this work, you
 may distribute the resulting work only under the same, similar or a
 compatible license.
- For any reuse or distribution, you must make clear to others the license terms of this work. The best way to do this is with a link to
 - http://creativecommons.org/licenses/by-sa/3.0/.
- Any of the above conditions can be waived if you get permission from the copyright holder.
- Nothing in this license impairs or restricts the author's moral rights.

