Lunds Universitet
Datavetenskapliga institutionen

DAT 401 – DM

# Prov på delkurs Diskret Matematik

## 20 mars 2002 kl 8–13 Gasque

**Tillåtna hjälpmedel:** Litteratur, egna anteckningar

Tentamen består av 3 uppgifter, varje uppgift kan ge 10 poäng. Deluppgifternas poäng anges inom [hakparenteser]. Observera att poänggivningen inte nödvändigtvis avspeglar deluppgiftens svårhetsgrad. Uppgifterna kan besvaras på svenska eller engelska.

Behandla högst en uppgift per papper (det går bra att behandla deluppgifter på samma papper). Skriv bara på ena sidan och markera varje sida med dina initialer. Skriv läsligt.

► a)[0] Hur många uppgifter får man behandla per papper?

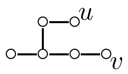
► b)[0] Får man skriva på baksidan?

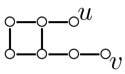
Det går bra att referera till kursboken av Rosen eller Husfeldts anteckningar. Referenser måste ha formen [R: *något*] eller [H: *något*], där *något* är ett avsnittsnummer, en kodbit, en räknad uppgift, etc. Skriv till exempel "genom att använda Lemma [R:7.4]".

# Exercise 1 (Graphs, Probability)

All graphs $G = (V, E)$ in this exercise are undirected.

The *distance* $d(u, v)$ between two vertices in a connected graph is the length of the shortest path between the vertices. For example, in both graphs below, the distance between $u$ and $v$ is 4.

$G_1 =$ 

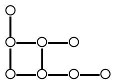
$G_2 =$ 


It is important to note that although $G_2$ does contain a path of length 6 from $u$ to $v$, this does not mean that their distance is 6: The *shortest* path from $u$ to $v$ is still length 4.

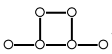
The *diameter* $\mathrm{diam}(G)$ of a graph is the longest distance between any pair of vertices; in symbols,

$$\mathrm{diam}(G) = \max_{u,v \in V} d(u, v).$$

We have $\mathrm{diam}(G_1) = \mathrm{diam}(G_2) = 4$, but $\mathrm{diam}(G_3) = 5$ for the graph below:

$G_3 =$ 

because there is a pair of vertices at distance 5.

▶ a)[1] Find the diameter of $G_4 =$  .

▶ b)[1] Draw a graph with 4 vertices and diameter 1.

▶ c)[1] How many edges are there in a graph with $n$ vertices and diameter 1?

The remaining questions study the diameter of a random graph. Let $G(n, p)$ denote a random graph on $n$ vertices with edge probability $p$. The vertices are numbered $v_1, \ldots, v_n$.

▶ d)[1] Find the probability that a random graph has diameter 1. In symbols, find $p\big(\mathrm{diam}(G(n, p)) = 1\big)$.

In the remaining questions you may fix $p = \frac{1}{2}$ if that makes it easier for you.

▶ e)[1] Find the probability that the distance between $v_1$ and $v_2$ in a random graph is greater than 1. In symbols, find $p\big(d(v_1, v_2) > 1\big)$.

▶ f)[3] Find the probability that the distance between $v_1$ and $v_2$ in a random graph is greater than 2. In symbols, find $p\big(d(v_1, v_2) > 2\big)$. *Hint:* Consider, for each other vertex $v_i$, the probability of a path $v_1 v_i v_2$.

▶ g)[2] Bound the probability that the diameter of a random graph is not 2. In symbols, show that $p\big(\mathrm{diam}(G(n, p)) \neq 2\big)$ is small. Use your answer to the previous question for this.

The conclusion is that a random graph has diameter 2 with high probability. In other words, and more dramatically: almost every graph has diameter 2.

## Exercise 2 (Combinatorics, Sets, and Relations)

Let $S = \{1, \ldots, n\}$. A *partition* of $S$ is a collection of disjoint, nonempty subsets $T_1, \ldots, T_k \subseteq S$ whose union is $T$. In symbols,

$$T_i \neq \emptyset \text{ for } 1 \leq i \leq k; \quad T_i \cap T_j = \emptyset \text{ for } i \neq j; \quad \bigcup_{i=1}^{k} T_i = S.$$

Here is a partition of $S$ for $n = 6$:

$$T_1 = \{1, 4\}, \quad T_2 = \{2, 3, 5\}, \quad T_3 = \{6\}.$$

We don't care about the numbering of the sets $T_i$, so the partition above is the same as the following:

$$T_1 = \{2, 5, 3\}, \quad T_2 = \{6\}, \quad T_3 = \{4, 1\}.$$

Here are the 5 possible partitions for $n = 3$:

$$T_1 = \{1\}, T_2 = \{2\}, T_3 = \{3\}. \qquad T_1 = \{1, 2, 3\}.$$
$$T_1 = \{1\}, T_2 = \{2, 3\}. \qquad T_1 = \{2\}, T_2 = \{1, 3\}. \qquad T_1 = \{3\}, T_2 = \{1, 2\}.$$

▶ a)[1] Find all partitions for $n = 4$.

Given a partition $T_1, \ldots, T_k$ of $S$ define the relation $R\colon S \times S$ where $(x, y) \in R$ if and only if $x$ and $y$ belong to the same subset. In symbols:

$$(x, y) \in R \quad \text{if and only if} \quad \exists i\colon x \in T_i \wedge y \in T_i.$$

▶ b)[1] Is $R$ symmetric? If yes, write 'yes'. If no, give a (small) counterexample.

▶ c)[1] Is $R$ reflexive? If yes, write 'yes'. If no, give a (small) counterexample.

▶ d)[2] Is $R$ transitive? If yes, give a short formal proof. If no, give a (small) counterexample.

Let $B(n)$ denote the number of partitions of $S$ for given $n$ (these are called *Bell* numbers after Eric Temple Bell[1]). For example, we showed $B(3) = 5$ above.

The number $\left\{ {n \atop k} \right\}$ is defined (for $1 \leq k \leq n$) to be the number of ways to partition $\{1, \ldots, n\}$ into exactly $k$ sets (for example, $\left\{ {3 \atop 2} \right\} = 3$).

▶ e)[2] Give a (short) combinatorial argument for

$$B(n) = \sum_{k=1}^{n} \left\{ {n \atop k} \right\}.$$

▶ f)[3] Give an argument for the following recurrence (don't try to solve it!).

$$B(n+1) = \sum_{i=0}^{n} \binom{n}{i} B(i).$$

You can either give a combinatorial argument (recommended), or give an algebraic proof based on the previous question. In that case, you probably want to use

$$\left\{ {n+1 \atop i+1} \right\} = \sum_{k=0}^{n} \binom{n}{k} \left\{ {k \atop i} \right\},$$

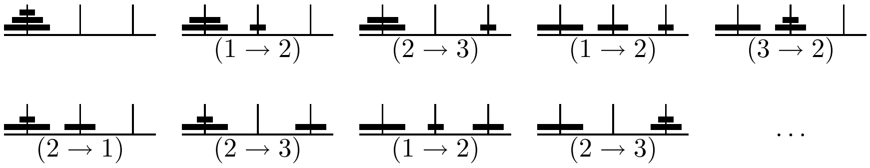
$B(0) = 1$, and $\left\{ {n \atop k} \right\} = 0$ for $k > n$ or $k = 0$ (you don't need to prove these things).

[1]Who wrote science fiction under the pseudonym John Taine

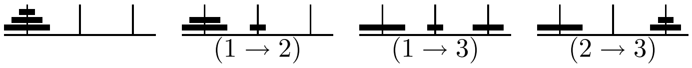
## Exercise 3 (Recurrences)

Consider a variation of the Towers of Hanoi puzzle, which we will call *Tired Hanoi*. In Tired Hanoi we want to move all discs from peg 1 to peg 3, but we forbid moving a disc directly from peg 1 to peg 3 or from peg 3 to peg 1. In other words, all moves must involve peg 2. (Imagine the monks are very tired and cannot move the discs over long distances.)

Consider for example the first eight moves in the case with $n = 3$ discs, all starting on peg 1. (We use the notation $(i \to j)$ for moving the top disc of peg $i$ to the top of peg $j$.)




In the standard version of Towers of Hanoi (where the moves $(1 \to 3)$ and $(3 \to 1)$ are allowed), the same arrangement can be reached with only three moves:




Evidently, Tired Hanoi takes more time than standard Hanoi. Also please note that Tired Hanoi requires the discs to end on peg 3; while standard Hanoi may end on peg 2 or peg 3.

- ▶ a)[1] Finish the example for $n = 3$. Use the $(i \to j)$ notation instead of drawing lots of pretty pictures. State the total number of moves (including the first eight moves from the example).
- ▶ b)[4] Find a recurrence relation for $T_n$, the number of moves needed to solve the Tired Hanoi problem with $n$ discs.
- ▶ c)[3] Solve the recurrence you found in the previous question. If you didn't solve the previous question, solve $D_1 = 3, D_n = 3 + 3D_{n-1}$ instead.
- ▶ d)[2] [Cute but not really hard:] Show that every legal arrangement of the $n$ discs on the 3 pegs is encountered during the solution of Tired Hanoi.