EDAP15: Program Analysis

DATA FLOW ANALYSIS 1

Christoph Reichenbach

Welcome back!

Some Administrativa:

- Labs:
  - CodeProber study with Anton:
    - Participation is 100% optional, but greatly appreciated
    - If you participate:
      Please make sure to check in auto-genered log files!
  - Ideally: done with exercise 0, working on exercise 1
  - No new exercise this week (exercise 1 is the biggest one)
  - Optional polls on how much time exercises took in Moodle
- Lectures:
  - Guest lecturer on Wednesday: **Alexandru Dura**
  - Guest lecture on 14 February:
    Patrik Åberg, **Magnus Templing** from Ericsson

Questions?

Getting More out of Type Inference (1/3)

- Recall our typing rules from last lecture:

$$\frac{}{\texttt{true} : \textsc{Bool}} \qquad \frac{}{\texttt{false} : \textsc{Bool}}$$

$$\frac{e_1 : \textsc{Bool} \quad e_2 : \tau \quad e_3 : \tau}{\texttt{if}\ e_1\ \texttt{then}\ e_2\ \texttt{else}\ e_3 : \tau}$$

- Could we make them more precise so we can e.g. tell that:
  - `if false then false else true` evaluates to `true`
  - `if true then false else false` evaluates to `false`

Getting More out of Type Inference (2/3)

Replacing $\textsc{Bool}$ by *more precise types* $\textsc{True}$ and $\textsc{False}$:

$$\frac{}{\texttt{true} : \textsc{True}} \qquad \frac{}{\texttt{false} : \textsc{False}}$$

$$\frac{e_1 : \textsc{True} \quad e_2 : \tau}{\texttt{if } e_1 \texttt{ then } e_2 \texttt{ else } e_3 : \tau}\ (\textit{if-true})$$

$$\frac{e_1 : \textsc{False} \quad e_3 : \tau}{\texttt{if } e_1 \texttt{ then } e_2 \texttt{ else } e_3 : \tau}\ (\textit{if-false})$$

We can now infer:

$$\frac{\frac{}{\texttt{false} : \textsc{False}} \quad \frac{}{\texttt{true} : \textsc{True}}}{\texttt{if false then false else true} : \textsc{True}}\ (\textit{if-false})$$

Getting More out of Type Inference (3/3)

- Consider:

```
fun(x) = if x then false else true
```

- Can't know x in general $\Rightarrow$ we must allow both:
  - x : True
  - x : False

$\Rightarrow$ We don't have a principal type any more
(without adding nontrivial extra structure [Dolan & Mycroft, 2017])

- How would this work for `int`?

Towards Abstract Interpretation

- Consider the following language:

$$
\begin{array}{rcl}
e & ::= & \mathsf{zero} \\
 & | & \mathsf{one} \\
 & | & \langle e \rangle \mathbf{+} \langle e \rangle \\
 & | & \mathsf{neg}\ \langle e \rangle
\end{array}
$$

- **Property of Interest**:
  Does a given program $\varphi \in e$ compute a number ≥ 0?
- We will use a different theoretical framework now:
  Abstract Interpratation
- Similar in many ways, but more suitable for the "subtyping"-like behaviour we just saw

Abstract Domains

- *Abstract Interpretation*:
  - Map all values to a simpler *abstract domain*

- Example: set *abstract domain* $\mathcal{D}$:

$$
\begin{array}{lll}
\mathcal{D} = & \{ & \\
& & D^0, \quad \text{Program computes 0} \\
& & D^+, \quad \text{Program computes a positive value} \\
& & D^- \quad\;\; \text{Program computes a negative value} \\
& \} &
\end{array}
$$

Patrick Cousot & Radhia Cousot, "Abstract Interpretation", published in Principles of Programming Languages, 1977

Correspondence: Concrete and Abstract

$$\underbrace{\cdots \quad -3 \quad -2 \quad -1}_{} \quad 0 \quad \underbrace{1 \quad 2 \quad 3 \quad \cdots}_{} \qquad\qquad D^{-} \qquad D^{0} \qquad D^{+}$$

Abstract Domains

- *Abstract Interpretation*:
  - Map all values to a simpler *abstract domain*

- Example: set *abstract domain* $\mathcal{D}$:

$$\mathcal{D} = \{ \begin{array}{ll} D^0, & \text{Program computes 0} \\ D^+, & \text{Program computes a positive value} \\ D^- & \text{Program computes a negative value} \end{array} \}$$

- Notation: $\boxed{\varphi \leadsto^D a}$, where $a \in \mathcal{D}$

Patrick Cousot & Radhia Cousot, “Abstract Interpretation”, published in Principles of Programming Languages, 1977

Abstract Interpretation

$$
\begin{aligned}
\ominus\, D^0 &= D^0 \\
\ominus\, D^+ &= D^- \\
\ominus\, D^- &= D^+ \\
\ominus\, D^? &= D^?
\end{aligned}
$$

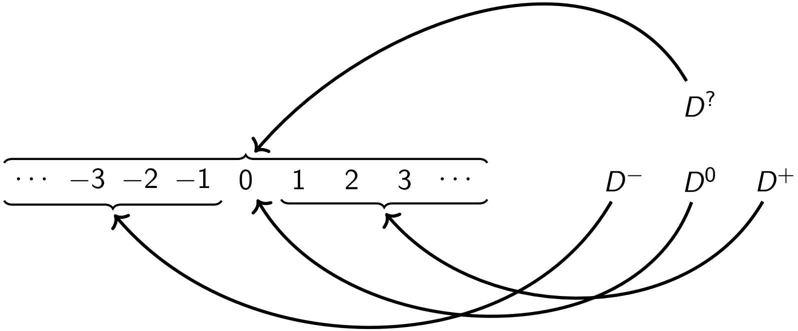
$$
\begin{aligned}
e \;::=\;& \mathsf{zero} \\
\mid\;& \mathsf{one} \\
\mid\;& \langle e\rangle + \langle e\rangle \\
\mid\;& \mathsf{neg}\ \langle e\rangle
\end{aligned}
$$

$$
a_1 \oplus a_2 = \left\{
\begin{array}{c||c|c|c|}
 & D^+ & D^0 & D^- \\
\hline\hline
D^+ & D^+ & D^+ & D^? \\
D^0 & D^+ & D^0 & D^- \\
D^- & D^? & D^- & D^-
\end{array}
\right.
\qquad
D^? \oplus a = D^? = a \oplus D^?
$$

$$
\mathsf{zero} \leadsto^D D^0 \qquad \mathsf{one} \leadsto^D D^+
$$

$$
\frac{x \leadsto^D a}{\mathsf{neg}\ x \leadsto^D \ominus a}
\qquad
\frac{x \leadsto^D a_1 \quad y \leadsto^D a_2}{x + y \leadsto^D a_1 \oplus a_2}
$$

Correspondence: Concrete and Abstract



Abstract Domains

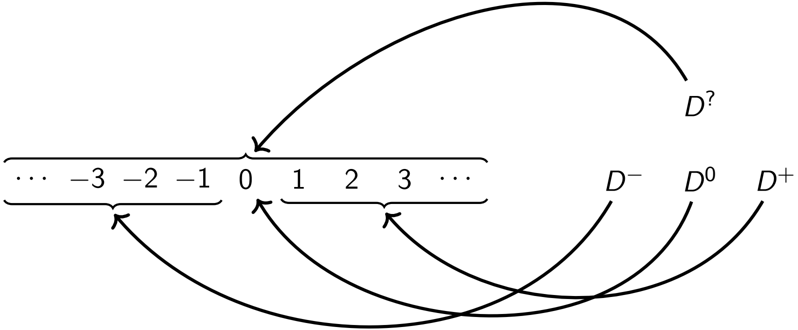
- *Abstract Interpretation*:
  - Map all values to a simpler *abstract domain*
  - For each operation, build an *abstract operations*
- Example: set *abstract domain* $\mathcal{D}$:

$$\mathcal{D} = \{ \quad D^0, \; \text{Program computes 0} \quad D^+, \; \text{Program computes a positive value} \quad D^-, \; \text{Program computes a negative value} \quad D^? \; \text{Program computes any value} \}$$

- Notation: $\boxed{\varphi \leadsto^D a}$, where $a \in \mathcal{D}$

Patrick Cousot & Radhia Cousot, "Abstract Interpretation", published in Principles of Programming Languages, 1977

Correspondence: Concrete and Abstract



Also:

- $\ominus$ *"is compatible with"* neg
- $\oplus$ *"is compatible with"* $+$

Will return later to examine connections between elements in $\mathcal{D}$

Summary

- **Abstract Interpretation** maps concrete values to "abstract value" in **Abstract Domain**
  - Mapping extends from just values to include expressions, statements, whole programs

$$\leadsto^{D}: \mathit{Program} \to \mathcal{D}$$

  - Map operations to *compatible* operations on abstract domain
  - Many design options for abstract domain:
    - Challenge: precision vs. decidability
  - Theoretical foundation/generalisation of other analysis theories
- Unlike type inference, $\leadsto^{D}$ is a *function*
- Other common notation (instead of $\varphi \leadsto^{D} a$):

$$[\![\varphi]\!]^{\mathcal{D}} = a$$

Lecture Overview

Foundations

Static Analysis

Dynamic Analysis

Properties

Control Flow

01 Foundations

02 Constructing Program Analyses in JastAdd

03 04 Types

05 06 07 **Data Flow**

08 09 Memory

05 Intraprocedural

10 Interprocedural

11 Indirect

12 Instrumentation

13 Analysis

14 Review

Teal

Teal-0	Imperative and Procedural
Teal-1	Minor extensions to Teal-0
Teal-2	User-Defined Data Types
Teal-3	Dynamic Dispatch, Inheritance

- Small enough for homework exercises
- Big enough to exhibit real challenges
- "Nonsensical" operations in Teal trigger dynamic *failure*s:
  - `null` dereference:

Teal

```
var a := null;
print(a[0]);
```

  - Array-out-of-bounds access:

Teal

```
var a := [7]; // List with one element, 7
print(a[-1]);
```

A New Analysis Challenge

Teal

```
var x := [0, 0];
print(x);    // A
if z {
  x[0] := 2; // B
  x := null;
}
x[0] := 1;   // C
```

- Analyse: Can there be a *failure* at `B` or `C`?
- Must distinguish between `x` at `A` vs. `x` at `B` and `C`
- Need to model program flow: **Flow-Sensitive Analysis**
  - Type inference is *not Flow-Sensitive*
  - Abstract interpretation can be Flow-Sensitive

Need analysis that can represent *data flow* through program

Evaluation Order

Teal-0

```
fun p(a) = { print(a); return 1; }
fun main() = {
    p(p(0) + p(1));
}
```

Teal-0 with explicit order

```
fun main() = {
    var tmp1   := p(0);
    var tmp2   := p(1);
    var tmp3   := tmp1 + tmp2;
    var tmp4   := p(tmp3);
}
```

- Evaluation order specified in language definition

Every analysis must remember the evaluation order rules!

Evaluation Order: Other Languages

- Complex subexpressions / evaluation order:

Java / C / C++

```
// Many challenging constructions:
a[i++] = b[i > 10 ?  i-- :  i++] + c[f(i++, --i)];
```

- Beware: exact evaluation order is *undefined* in C and C++!
- Short-Circuit Evaluation:

Java (similar in C / C++)

```
int[] a2 = some_array;
bool v = (a == null)
         || ((a2 = a)[0] == 0);
```

- The assignment `a2 = a` is executed while computing `v` …but *only* if `a == null` is **not true**!

Violates most coding styles, but allowed by language!

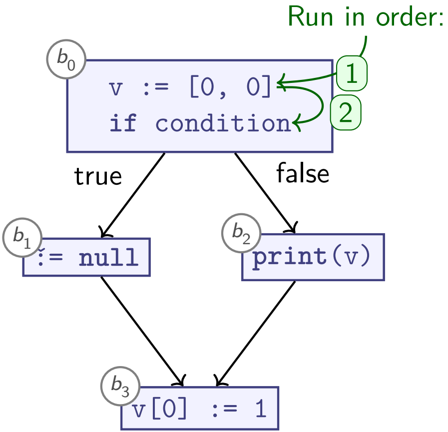
Summary

- Understanding differences before/after variable updates requires **Flow-Sensitive Analysis**
- Type inference is *not* flow sensitive
- "Flow" is complicated, influenced by:
  - Expression evaluation order
  - Short-circuit evaluation
  - Statement execution order

Control-Flow Graphs (CFGs)

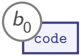
Teal

```
var v := [0, 0];
if condition {
  v := null;
} else {
  print(v);
}
v[0] := 1;
```

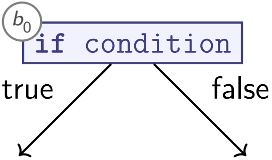


Control Flow Graphs encode statement execution order

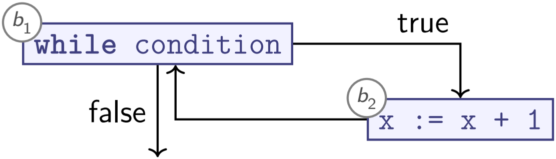
Control-Flow-Graphs

- Encode statement order by *nodes* 

and edges →
- *Multiple* outgoing edges (branches): Add label:



- Uniform representation for control statements:



Summary

Control-Flow Graph (CFG):

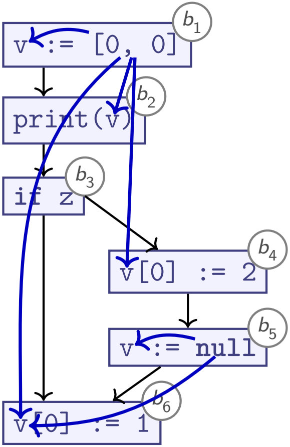
- Motivation:
  - Universal representation of control flow
  - Computed once before running analyses
  - Flow-sensitive analyses can utilise CFG
- Idea:
  - Represent control flow as **Blocks** and **Control-Flow Edges**
  - Edges represent control flow, **labelled** to identify conditionals

Control Flow

Understanding data flow requires understanding control flow:

Teal

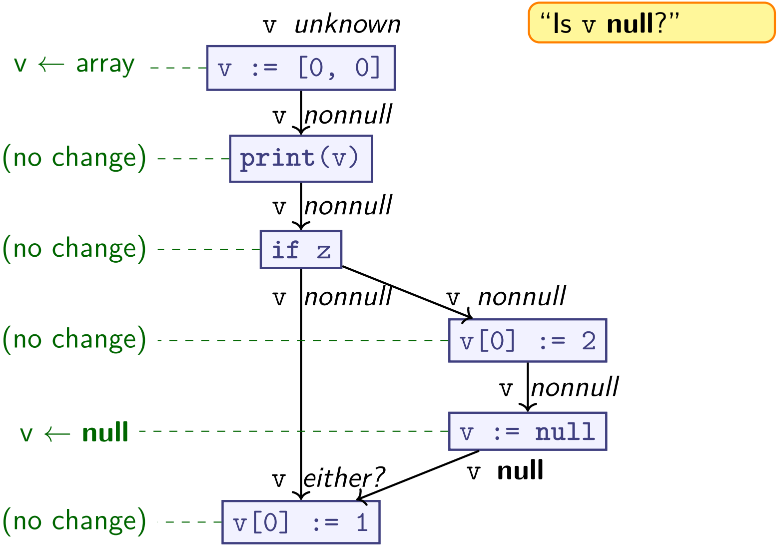
```
var v := [0, 0];
print(v);
if z {
  v[0] := 2;
  v := null;
}
v[0] := 1;
```



⟶ Control flow

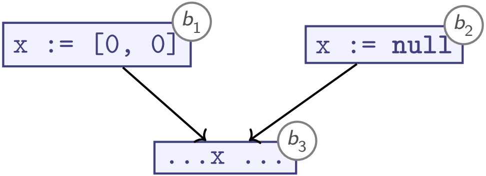
⟶ Data flow

Intuition behind Data Flow Analysis



Knowledge about data “flows” through CFG

What does "either?" mean?



- Should analysis report `x` as **null** or as **nonnull**?
  - New category: **either**
  - "Can I safely dereference without a check?"
    $\Rightarrow$ better assume **null**
  - "Is this guaranteed to be null?"
    $\Rightarrow$ better assume **nonnull**
- We might not need extra **either** category, depending on what properties we are looking for

“May” vs “Must” Analysis

- “**May**” analysis: we *cannot rule out* property
  - “either?” becomes **true**
  - Avoids *False Negatives*
- “**Must**” analysis: we *can guarantee* property
  - “either?” becomes **false**
  - Avoids *False Positives*

Another Analysis

Teal

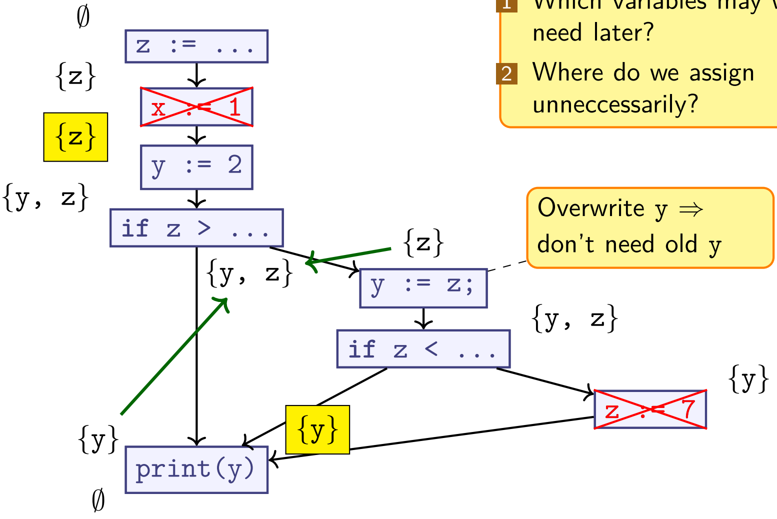
```
z := ...
x := 1;
y := 2;
if z > ... {
  y := z
  if z < ... {
    z := 7
  }
}
print(y);
```

- Which assignments are unnecessary?

$\Rightarrow$ Possible oversights / bugs
(*Live Variables Analysis*)

Unnecessary Assignments: Intuition

1. Which variables may we need later?
2. Where do we assign unneccessarily?

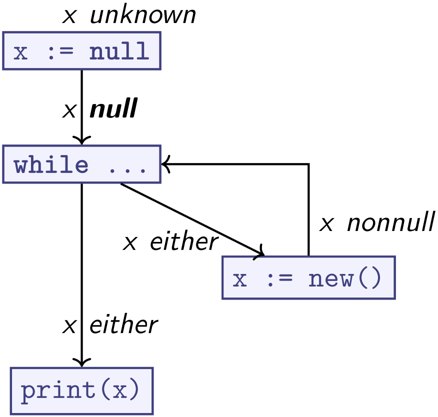


Analysis effective: found useless assignments to z and x

Observations

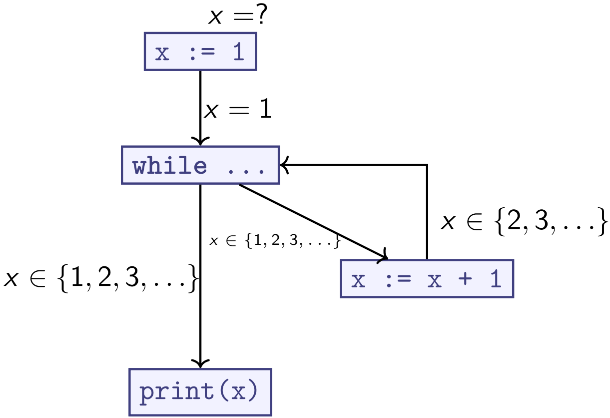
1. Data Flow analysis can be run *forward* or *backward*
2. May have to *join* results from multiple sources
3. Some analyses may need multiple "passes" (steps)

What about Loops? (1/2)



- Analysis: *Null Pointer Dereference*
- May need to analyse each node/edge more than once
- Stop when we're not learning anything new any more

What about Loops? (2/2)



- Analysis: *Reaching Values* / *Reaching Definitions* / *Copy Propagation*

We need to bound repetitions!

Summary: Data-Flow Analysis (Introduction)

- Data flow depends on *control flow*
- Data flow analysis examines how variables or other program state change across control-flow edges
- May have to join multiple results
- When joining "yes" and "no", must decide:
  - "**May**" analysis: optimistically report what is possible
  - "**Must**" analysis: conservatively report what is guaranteed
  - Alternative: introduce value for "don't know"
- Can run *forward* or *backward* relative to control flow edges
- Handling loops is nontrivial

Summary: Some Analyses

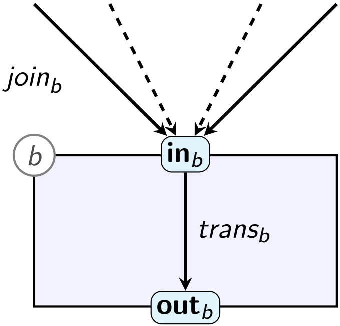
1. **Reaching Values / Reaching Definitions** (Also "Copy Propagation"):
   - *What values might our variables contain?*
   - Forward analysis
   - Most common as a *Must* analysis, where either:
     - 'v has constant value c', or
     - 'v might not have constant value'
   - We will also use it as *May* analysis
2. **Live Variables**
   - *Which variables might still be read later in the program?*
   - Backward analysis
   - *May* analysis
3. **Unnecessary Assignments** (also "Dead Assignments"):
   - Refinement of *Live Variables* analysis
   - Flags assignments on variables that are not live

Engineering Data Flow Algorithms

1. General Algorithm
   - Keep updating until nothing changes
   - JastAdd: *Circular Attributes*
2. Termination
   - Assumption: Operate on Control Flow Graph
   - Theory: Ensure termination
3. (Correctness)

Data Flow Analysis on CFGs

- $\mathbf{in}_b$: knowledge at entrance of basic block b
- $\mathbf{out}_b$: knowledge at exit of basic block b
- $join_b$: combines all $\mathbf{out}_{b_i}$ for all basic blocks b_i that flow into b "**Join Function**"
- $trans_b$: updates $\mathbf{out}_b$ from $\mathbf{in}_b$ "**Transfer Function**"



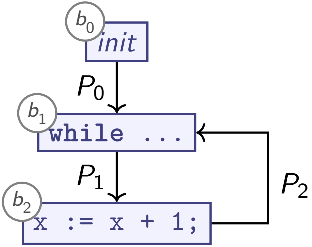
Characterising Data Flow Analyses

Characteristics:

- *Forward* or *backward* analysis
- L: Set of "abstract values" that represent our knowledge about the program
- $trans_b : L \to L$
- $join_b : L \times L \to L$

$$\boxed{\textbf{Require properties of } L\textbf{, } trans_b\textbf{, } join_b \textbf{ to ensure termination}}$$

Limiting Iteration



- Does the following ever stop changing:

$$\begin{array}{rcl} \mathbf{in}_{b_1} & = & join_{b_1}(P_0, P_2) \\ \mathbf{in}_{b_2} & = & trans_{b_1}(\mathbf{in}_{b_1}) \\ P_2 & = & trans_{b_2}(\mathbf{in}_{b_2}) \end{array}$$

- Intuition: we keep generalising information
  - *Growth limit*: bound amount of generalisation
  - Make sure $join_b$, $trans_b$ never throw information away

Eventually, either nothing changes or we hit growth limit

Ordering Knowledge

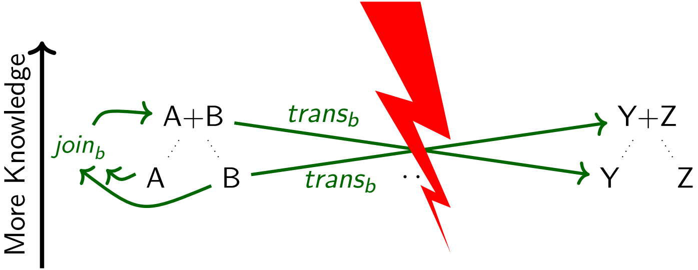
$$A \sqsubseteq B$$

$$\begin{tikzcd} B \arrow[d, dash] \\ A \end{tikzcd}$$

- B describes at least as much knowledge as A
- Either:
  - $A = B$ (i.e., $A \sqsubseteq B \sqsubseteq A$), or
  - B has strictly more knowledge than A

Intuition: Knowing Less, Knowing More

Structure of L:

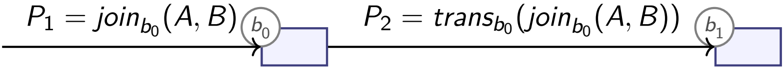


- $join_b$ must not lose knowledge
  - $A \sqsubseteq join_b(A, B)$
  - $B \sqsubseteq join_b(A, B)$
- $trans_b$ must be *monotonic* over amount of knowledge:

$$x \sqsubseteq y \implies trans_b(x) \sqsubseteq trans_b(y)$$

- Introduce bound: $\top$ means 'too much information'

Aggregating Knowledge



- Interplay between $trans_b$ and $join_b$ helps preserve knowledge
- $A \sqsubseteq join_b(A, B)$:
  As we add knowledge, P_1 either:
  - Stays the same
  - Increases knowledge
- Monotonicity of $trans_b$: If P_1 goes up, then P_2 either:
  - Stays the same
  - Increases knowledge

$\Rightarrow$ At each node, we either stay equal or grow

Now we must only prevent *infinite* growth. . .

Ascending Chains

$$\begin{array}{l} a_k = a_{k+1} = \dots \\ \vdots \\ a_3 \\ | \\ a_2 \\ | \\ a_1 \\ | \\ a_0 \end{array}$$

- A (possibly infinite) sequence $a_0, a_1, a_2, \dots$ is an *ascending chain* iff:

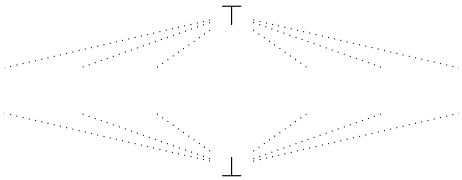
$$a_i \sqsubseteq a_{i+1} \text{ (for all } i \geq 0\text{)}$$

- *Ascending Chain Condition*:
  - For *every* ascending chain $a_0, a_1, a_2, \dots$ in abstract domain L:
  - there exists $k \geq 0$ such that:

$$a_k = a_{k+n} \text{ for any } n \geq 0$$

ACC is formalisation of growth limit

Top and Bottom



- *Convention*: We introduce two distinguished elements:
  - **Top**: $\top$: $A \sqsubseteq \top$ for all A
  - **Bottom**: $\bot$: $\bot \sqsubseteq A$ for all A
- Since $A \sqsubseteq join_b(A, B)$ and $B \sqsubseteq join_b(A, B)$:
  - $join_b(\top, A) = \top = join_b(A, \top)$
  - $\bot \sqsubseteq A \sqsubseteq join_b(\bot, A)$
    - In practice, it is safe and simple to set: $join_b(\bot, A) = A = join_b(A, \bot)$
- Intuition:
  - $\top$: means 'contradictory / too much information'
  - $\bot$: means 'no information known yet'

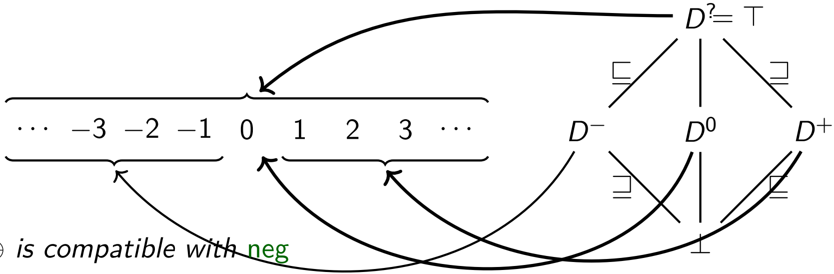
Summary

- Designing a *Forward* or *backward* analysis:
- Pick **Abstract Domain** L
  - Must be **partially ordered** with $(\sqsubseteq) \subseteq L \times L$:
    $A \sqsubseteq B$ iff B 'knows' at least as much as A
  - Unique top element $\top$
  - Unique bottom element $\bot$
- $trans_b : L \to L$
  - Must be *monotonic*:

$$x \sqsubseteq y \implies trans_b(x) \sqsubseteq trans_b(y)$$

- $join_b : L \times L \to L$ must produce an *upper bound* for its parameters:
  - $A \sqsubseteq join_b(A, B)$
  - $B \sqsubseteq join_b(A, B)$
- Satisfy **Ascending Chain Condition** to ensure termination
  - Easiest solution: make L finite

Abstract Domains Revisited



$\ominus$ *is compatible with* neg

$$
\begin{array}{lcl}
\ominus \bot & = & \bot \\
\ominus D^0 & = & D^0 \\
\ominus D^+ & = & D^- \\
\ominus D^- & = & D^+ \\
\ominus D^? & = & D^?
\end{array}
$$

$\ominus$ is monotonic (and $\oplus$ extended with $\bot$ is, too)

Summary

- We can extend $\{D^+, D^-, D^0, D^?\}$ by adding $\bot$

$$L_D = \{D^+, D^-, D^0, D^?, \bot\}$$

- $\bot$ representing "not known" – not needed for our example analysis ($\rightsquigarrow^D$), but would be needed if we had variables / control flow in that language
- L_D is finite, so the DCC holds trivially
- Our *Transfer Functions* $\ominus, \oplus$ are monotonic
  - Concretely, $\oplus$ is "pointwise monotonic", meaning: if $d \in L_D$ is constant, then
    - $x \mapsto d \oplus x$ is monotonic
    - $x \mapsto x \oplus d$ is monotonic

Outlook

- We will continue on Dataflow Analysis
- Next lecture held by **Alexandru Dura**

`http://cs.lth.se/EDAP15`