

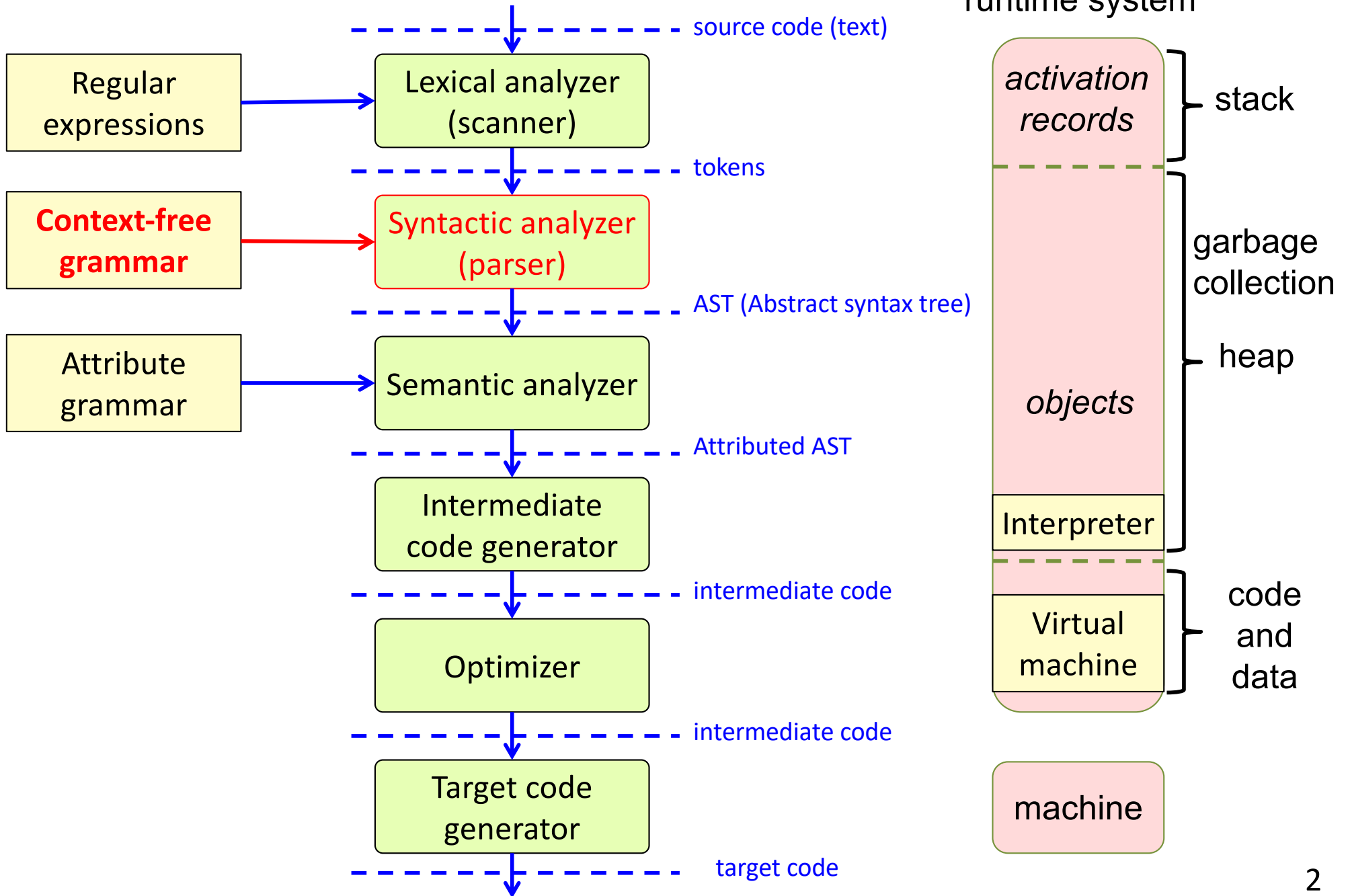
## EDAN65: Compilers, Lecture 04

Grammar equivalence,  
eliminating ambiguities,  
adapting to LL parsing

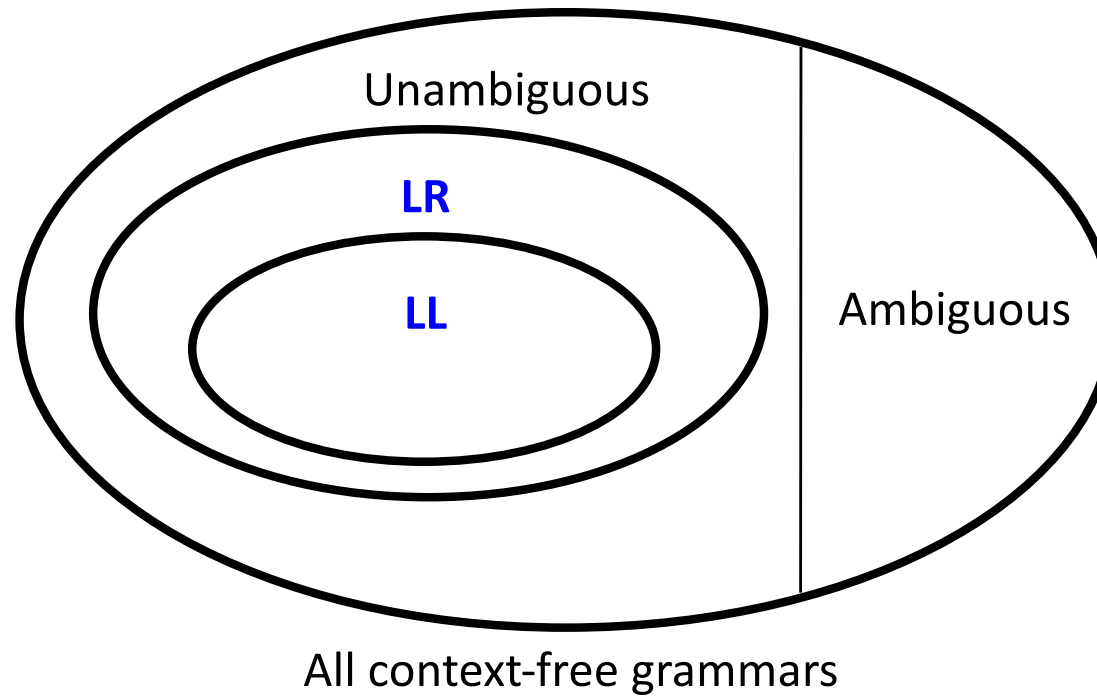
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Revised: 2023-09-04

# This lecture



# Space of context-free grammars



**LL:**

Builds tree top-down  
Simple to understand

**LR:**

Builds tree bottom-up  
More powerful

# Ambiguous grammars

# Recall: the definition of ambiguity

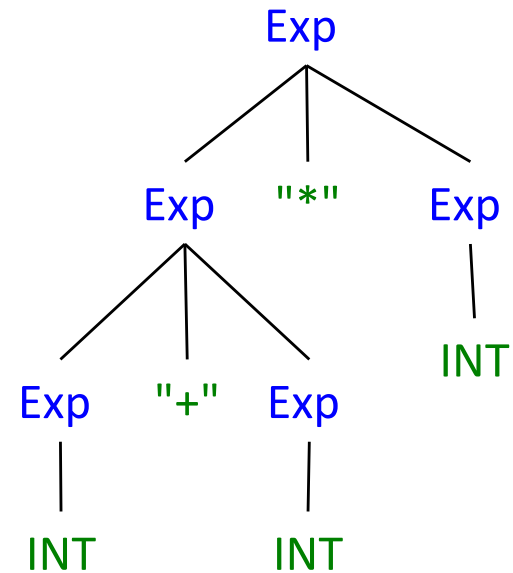
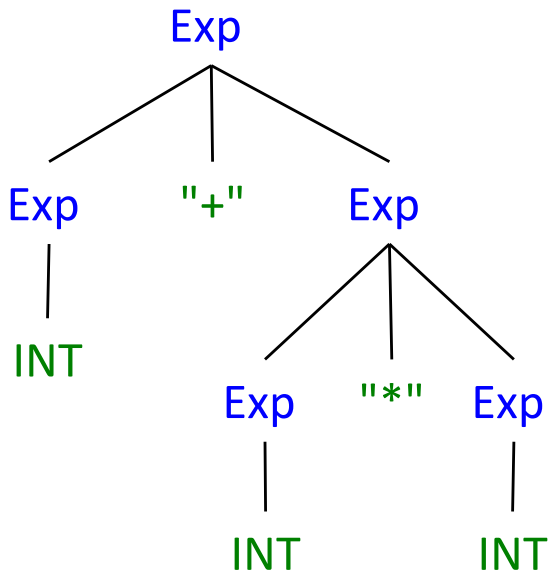
Grammar:

Exp  $\rightarrow$  Exp "+" Exp

Exp  $\rightarrow$  Exp "\*" Exp

Exp  $\rightarrow$  INT

A CFG is *ambiguous* if there is a sentence in the language that can be derived by two (or more) different parse trees.



# Strategies for dealing with ambiguities

# Strategies for dealing with ambiguities

First, decide which parse tree is the desired one.

## Eliminate the ambiguity:

Create an equivalent unambiguous grammar.

Often possible, but there exists grammars for which it cannot be done.

However, the parse tree will be different from the original desired one.

## Alternatively, some parser generators support disambiguation rules:

Use the ambiguous grammar.

Add priority and associativity rules to instruct the parser to select the desired parse tree.

Works for some ambiguities and some parser algorithms.

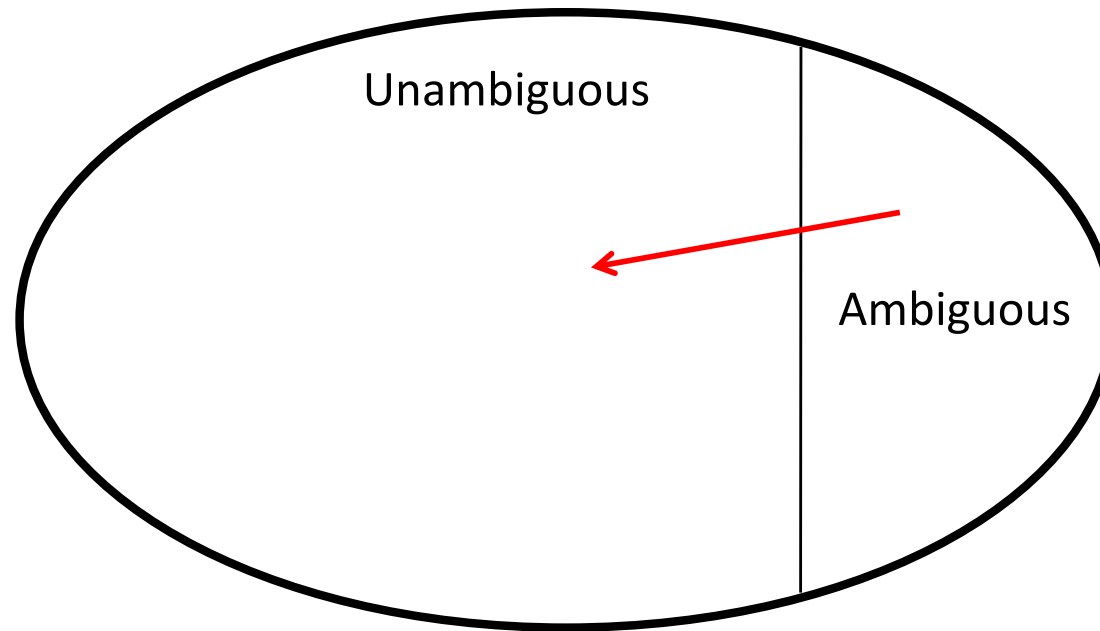
## Alternatively, use general parser:

Constructs all parse trees.

Disambiguate after parsing.

But general parsers are slow (cubic in input length)

# Eliminating ambiguity



Goal: transform an ambiguous grammar to an *equivalent* unambiguous grammar.



# Equivalent grammars

Two grammars,  $G_1$  and  $G_2$ , are *equivalent* if they generate the same language.

I.e., a sentence can be derived from one of the grammars, iff it can be derived also from the other grammar:

$$L(G_1) = L(G_2)$$

# Common kinds of ambiguities

- Binary operators with different priorities:

$a + b * c == d, \dots$

- Associativity of binary operators of the same priority:

$a + b - c + d, \dots$

- Dangling else:

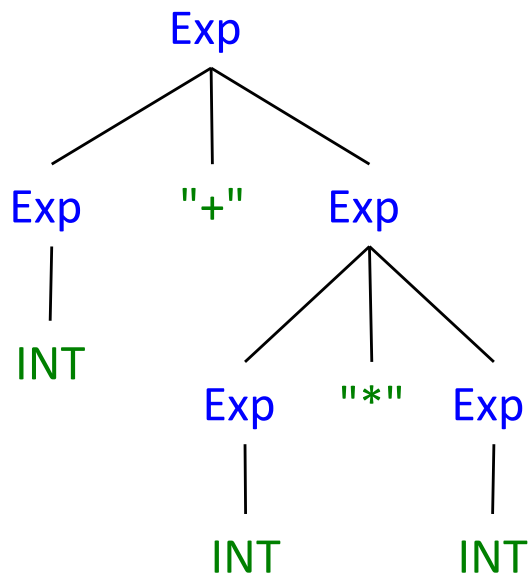
```
if (a)
  if (b) c = d;
  else e = f;
```

# Example ambiguity:

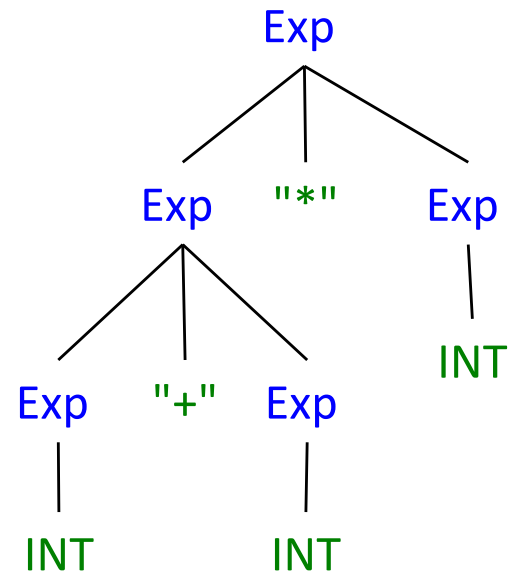
*Priority (also called precedence)*

```
Exp -> Exp "+" Exp
Exp -> Exp "*" Exp
Exp -> INT
```

Two parse trees for INT "+" INT "\*" INT



prio("\*") > prio("+")  
(according to convention)

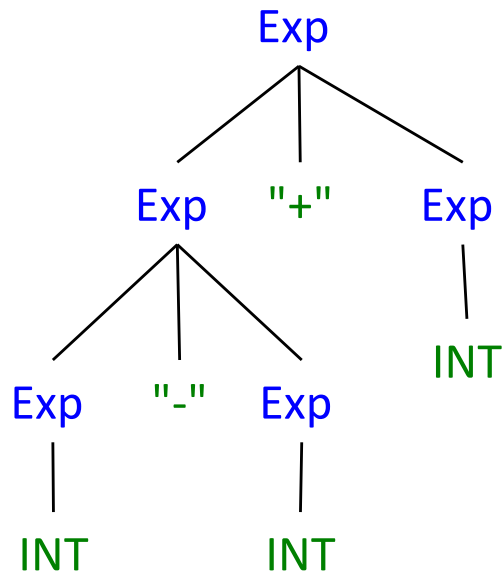


prio("+") > prio("\*")  
(would be unexpected and confusing)

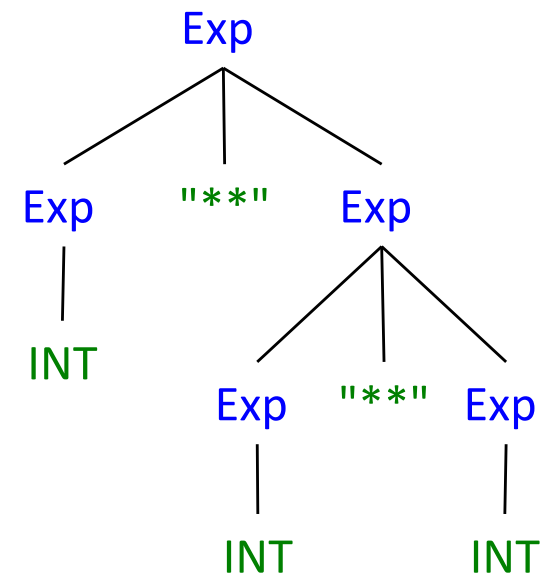
# Example ambiguity: *Associativity*

```
Exp -> Exp "+" Exp
Exp -> Exp "-" Exp
Exp -> Exp "***" Exp
Exp -> INT
```

For operators with the same priority,  
how do several in a sequence associate?



*Left-associative*  
(usual for most operators)



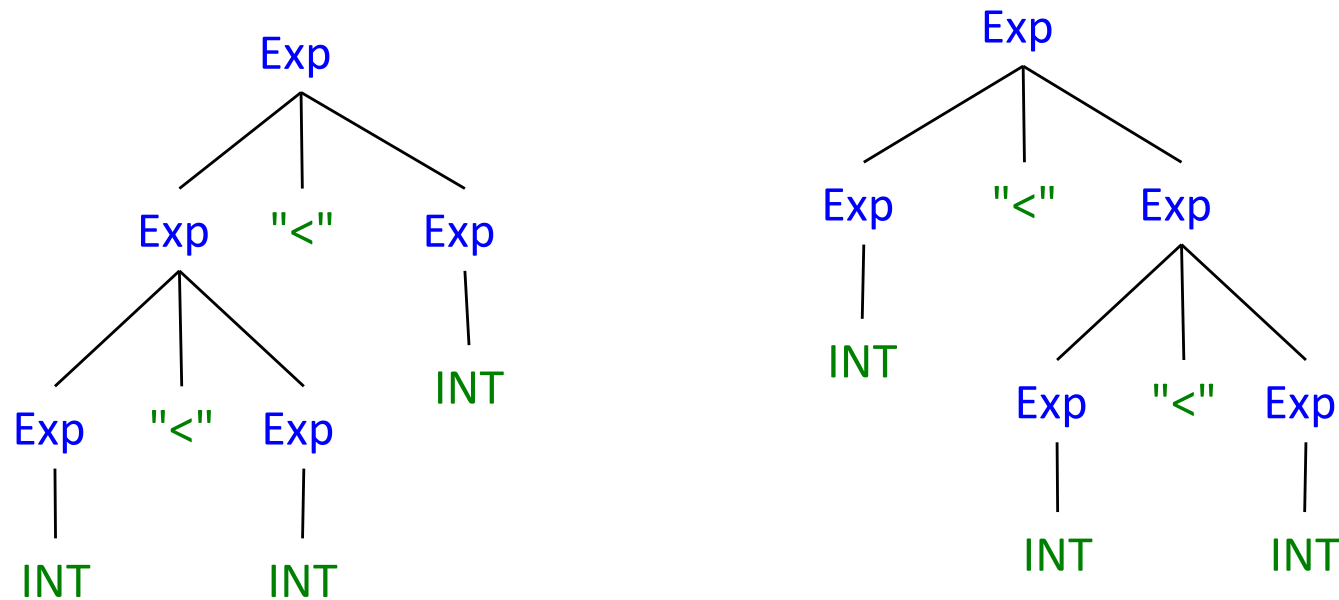
*Right-associative*  
(usual for the power operator)

# Example ambiguity:

*Non-associativity*

Exp  $\rightarrow$  Exp " $<$ " Exp  
Exp  $\rightarrow$  INT

For some operators, it does not make sense to have several in a sequence at all. They are *non-associative*.



We would like to forbid both trees.  
I.e., rule out the sentence from the language.

# Disambiguating expression grammars

How can we change the grammar so that only the desired trees can be derived?

# Disambiguating expression grammars

How can we change the grammar so that only the desired trees can be derived?

**Idea:** Replace some nonterminals with new ones that restrict what can be generated.

**Priority:** Introduce a new nonterminal for each priority level:  
**Term, Factor, Primary, ...**

**Left associativity:** Restrict the right operand so it only can contain expressions of higher priority

**Right associativity:** Restrict the left operand ...

**Non-associativity:** Restrict both operands

# Exercise

Ambiguous grammar:

Expr -> Expr "+" Expr

Expr -> Expr "\*" Expr

Expr -> ID

Expr -> "(" Expr ")"

Equivalent unambiguous grammar:



# Solution

**You will do this in Assignment 2!**

Ambiguous grammar:

```
Expr -> Expr "+" Expr
Expr -> Expr "*" Expr
Expr -> ID
Expr -> "(" Expr ")"
```

Equivalent unambiguous grammar:

```
Expr -> Expr "+" Term
Expr -> Term
Term -> Term "*" Factor
Term -> Factor
Factor -> ID
Factor -> "(" Expr ")"
```

Here, we introduce a new nonterminal, Term, that is more restricted than Expr. That is, from Term, we can not derive any new additions.

For the addition production, we use Term as the right operand, to make sure no new additions will appear to the right. This gives left-associativity.

For the multiplication production, we use Term, and the even more restricted nonterminal Factor to make sure no additions can appear as children (without using parentheses). This gives multiplication higher priority than addition.

# Real-world example: The Java expression grammar

```
Expression -> LambdaExpression | AssignmentExpression
AssignmentExpression -> ConditionalExpression | Assignment
ConditionalExpression -> ...

AdditiveExpression ->
    MultiplicativeExpression
    | AdditiveExpression + MultiplicativeExpression
    | AdditiveExpression - MultiplicativeExpression

MultiplicativeExpression ->
    UnaryExpression
    | MultiplicativeExpression * UnaryExpression
    | ...

UnaryExpression -> ...
...
Primary -> PrimaryNoNewArray | ArrayCreationExpression
PrimaryNoNewArray -> Literal | this | ( Expression ) | FieldAccess ...
```

More than 15 priority levels.

See the Java Language Specification, Java SE 11, Chapter 19, Syntax

<https://docs.oracle.com/javase/specs/jls/se11/html/jls-19.html#jls-19-15>

# The "dangling else" problem

```
S -> "if" "(" E ")" S ["else" S]  
S -> ID "=" E ";"  
E -> ID
```

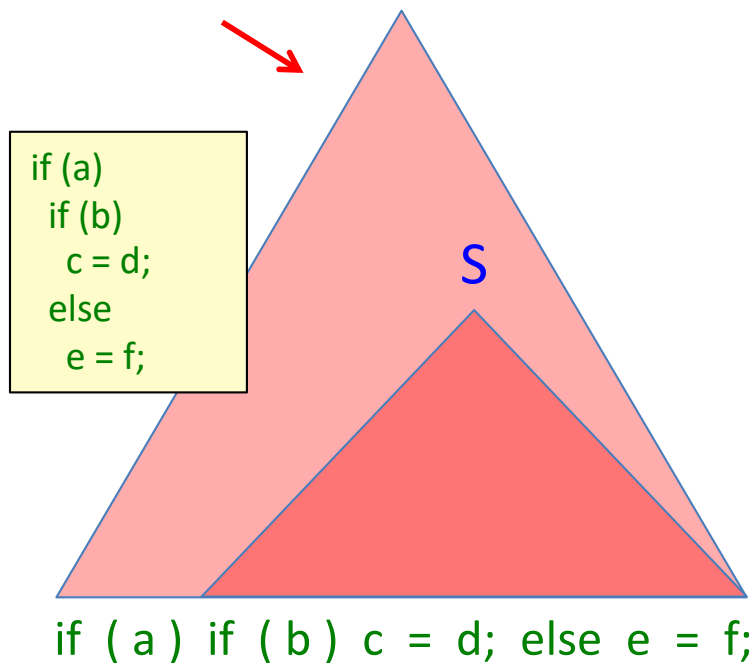
Construct a parse tree for: `if (a) if (b) c = d; else e = f;`

# The "dangling else" problem

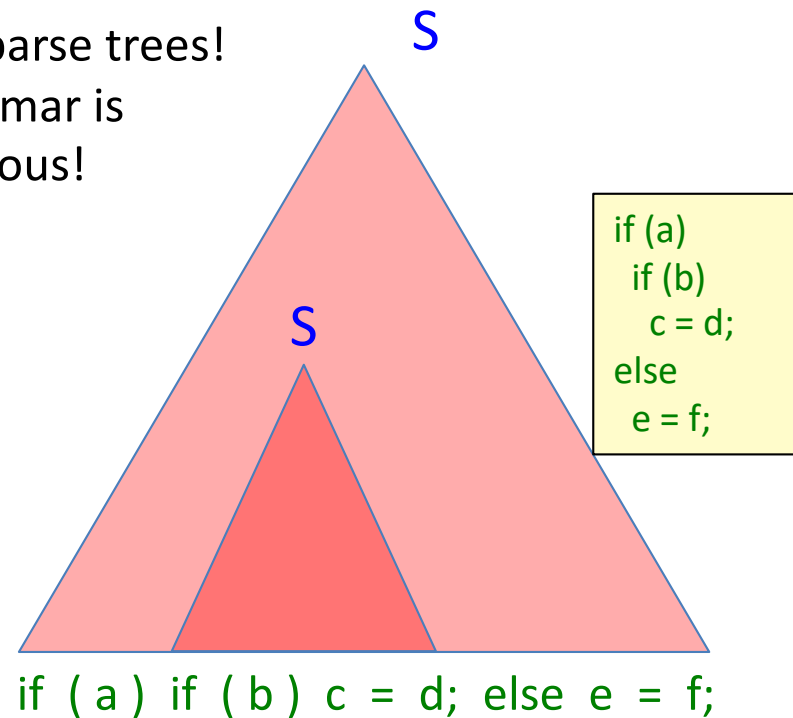
```
S -> "if" "(" E ")" S ["else" S]
S -> ID "=" E ";"
E -> ID
```

Construct a parse tree for: `if (a) if (b) c = d; else e = f;`

The desired tree



Two possible parse trees!  
The grammar is ambiguous!



# Solutions to the "dangling else" problem

## **Rewrite to equivalent unambiguous grammar**

- possible, but results in more complex grammar (several similar rules)

## **Use the ambiguous grammar**

- use "rule priority", the parser can select the correct rule.
- works for the dangling else problem, but not for ambiguous grammars in general
- not all parser generators support it well

## **Change the language?**

- e.g., add a keyword "fi" that closes the "if"-statement
- restrict the "then" part to be a block: "{ ... }". (Recommended for A2)
- only an option if you are designing the language yourself.

The Java Language Specification rewrites the grammar to be unambiguous.  
(See `IfThenStatement` and `IfThenElseStatement`.)

# Rewriting "dangling else"

Ambiguous grammar:

```
S -> "if" E S           // Short if
S -> "if" E S "else" S  // Long if
S -> "while" E "do" S
S -> ID "=" E ";"
S -> "{" S* "}"
```

# Rewriting "dangling else"

Ambiguous grammar:

```
S -> "if" E S           // Short if
S -> "if" E S "else" S  // Long if
S -> "while" E "do" S
S -> ID "=" E ";"
S -> "{" S* "}"
```

Solution idea: Limit S before "else" so that it cannot end with a short if.

Unambiguous grammar:

```
S -> "if" E S
S -> "if" E LimS "else" S
S -> "while" E "do" S
S -> ID "=" E ";"
S -> "{" S* "}"
LimS -> "if" E LimS "else" LimS
LimS -> "while" E "do" LimS
LimS -> ID "=" E ";"
LimS -> "{" S* "}"
```

# Dangling else in JLS

Uses a rewritten grammar using the pattern on the previous slide.

Results in duplication in the grammar.

See <https://docs.oracle.com/javase/specs/jls/se11/html/jls-19.html#jls-19-14>

```
Statement -> IfThenStatement | IfThenElseStatement | WhileStatement | ...
IfThenStatement -> "if" "(" Expression ")" Statement
IfThenElseStatement -> "if" "(" Expression ")" StatementNoShortIf "else" Statement
WhileStatement -> "while" "(" Expression ")" Statement
...

StatementNoShortIf -> IfThenElseStatementNoShortIf | WhileStatementNoShortIf | ...
IfThenElseStatementNoShortIf ->
    "if" "(" Expression ")" StatementNoShortIf "else" StatementNoShortIf
WhileStatementNoShortIf -> "while" "(" Expression ")" StatementNoShortIf
...
```

Other ambiguities are also rewritten using a similar pattern.

(See, e.g., `ClassType` vs `UnannClassType`)



# Disambiguation by requiring block before "else"

Ambiguous grammar:

```
S -> "if" E S  
S -> "if" E S "else" S  
S -> "while" E "do" S  
S -> ID "=" E ";"  
S -> "{" S* "}"
```

# Disambiguation by requiring block before "else"

Ambiguous grammar:

```
S -> "if" E S  
S -> "if" E S "else" S  
S -> "while" E "do" S  
S -> ID "=" E ";"  
S -> "{" S* "}"
```

Unambiguous grammar:

```
S -> "if" E S  
S -> "if" E "{" S* "}" "else" S  
S -> "while" E "do" S  
S -> ID "=" E ";"  
S -> "{" S* "}"
```

Recommendation: Use this approach in A2

# Finding ambiguities in practice

## **You try to run a CFG through an LL or LR parser generator**

- If it is accepted by the parser generator, the grammar is unambiguous
- If not, the grammar *could* be ambiguous, or unambiguous, but outside of the parser generator grammar class. In any case, you need to analyze that particular problem. This can be quite tricky, especially for large grammars. Perhaps you can find an ambiguity, or some other known LL/LR difficulty.

# Transforming to equivalent grammar

EBNF, BNF, Canonical form

# Recall: different notations for CFGs

$A \rightarrow B d e C f$   
 $A \rightarrow g A$

Canonical form

- *sequence* of terminals and nonterminals

$C \rightarrow D a b \mid b E F \mid a C$

BNF (Backus-Naur Form)

- alternative productions ( ... | ... | ... )

$G \rightarrow H^* i \mid (d E)^+ F \mid [d C]$

EBNF (Extended Backus-Naur Form)

- *repetition* (\* and +)
- *optionals* [...]
- *parentheses* (...)

# Writing the grammar in different notations

Canonical form:

```
Expr -> Expr "+" Term
Expr -> Term
Term -> Term "*" Factor
Term -> Factor
Factor -> INT
Factor -> "(" Expr ")"
```

Equivalent BNF (Backus-Naur Form):

Equivalent EBNF (Extended BNF):

# Writing the grammar in different notations

Canonical form:

```
Expr -> Expr "+" Term
Expr -> Term
Term -> Term "*" Factor
Term -> Factor
Factor -> INT
Factor -> "(" Expr ")"
```

Equivalent BNF (Backus-Naur Form):

```
Expr -> Expr "+" Term | Term
Term -> Term "*" Factor | Factor
Factor -> INT | "(" Expr ")"
```

Use *alternatives* instead of several productions per nonterminal.

Equivalent EBNF (Extended BNF):

```
Expr -> Term ("+" Term)*
Term -> Factor ("*" Factor)*
Factor -> INT | "(" Expr ")"
```

Use *repetition* instead of recursion, where possible.

# Translating EBNF to Canonical form

EBNF

Equivalent canonical form

Top level repetition

$X \rightarrow \gamma_1 \gamma_2^* \gamma_3$

Top level alternative

$X \rightarrow \gamma_1 \mid \gamma_2$

Top level parentheses

$X \rightarrow \gamma_1 (\dots) \gamma_2$

Where  $\gamma_k$  is a sequence of terminals and nonterminals



# Translating EBNF to Canonical form

EBNF

Top level repetition

$X \rightarrow \gamma_1 \gamma_2^* \gamma_3$

Equivalent canonical form

$X \rightarrow \gamma_1 N \gamma_3$

$N \rightarrow \varepsilon$

$N \rightarrow \gamma_2 N$

Top level alternative

$X \rightarrow \gamma_1 \mid \gamma_2$

$X \rightarrow \gamma_1$

$X \rightarrow \gamma_2$

Top level parentheses

$X \rightarrow \gamma_1 (\dots) \gamma_2$

$X \rightarrow \gamma_1 N \gamma_2$

$N \rightarrow \dots$

# Exercise:

## Translate from EBNF to Canonical form

EBNF:

$\text{Expr} \rightarrow \text{Term} ( "+" \text{Term} )^*$

Equivalent Canonical Form

# Solution:

## Translate from EBNF to Canonical form

EBNF:

$\text{Expr} \rightarrow \text{Term} ("+" \text{Term})^*$

Equivalent Canonical Form

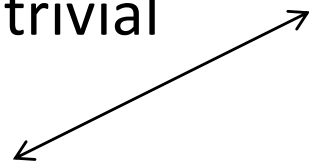
$\text{Expr} \rightarrow \text{Term } N$

$N \rightarrow \varepsilon$

$N \rightarrow "+" \text{Term } N$

# Can we prove that these are equivalent?

trivial



Equivalent Canonical Form

$\text{Expr} \rightarrow \text{Term } N$

$N \rightarrow \varepsilon$

$N \rightarrow "+" \text{Term } N$

EBNF:

$\text{Expr} \rightarrow \text{Term } ("+" \text{Term})^*$

non-trivial



Alternative Equivalent Canonical Form

$\text{Expr} \rightarrow \text{Expr } "+" \text{Term}$

$\text{Expr} \rightarrow \text{Term}$

# Example proof

1. We start with this:

$\text{Expr} \rightarrow \text{Term} ("+" \text{Term})^*$

We would like this:

$\text{Expr} \rightarrow \text{Expr} "+" \text{Term}$

$\text{Expr} \rightarrow \text{Term}$

# Example proof

1. We start with this:

$\text{Expr} \rightarrow \text{Term } "+" \text{ Term}^*$

2. We can move the repetition:

$\text{Expr} \rightarrow (\text{Term } "+")^* \text{Term}$

3. Eliminate the repetition:

$\text{Expr} \rightarrow \text{N Term}$

$\text{N} \rightarrow \epsilon$

$\text{N} \rightarrow \text{N Term } "+"$

4. Replace  $\text{N Term}$  by  $\text{Expr}$  in the third production:

$\text{Expr} \rightarrow \text{N Term}$

$\text{N} \rightarrow \epsilon$

$\text{N} \rightarrow \text{Expr } "+"$

We would like this:

$\text{Expr} \rightarrow \text{Expr } "+" \text{Term}$

$\text{Expr} \rightarrow \text{Term}$

5. Eliminate  $\text{N}$ :

$\text{Expr} \rightarrow \text{Expr } "+" \text{Term}$

$\text{Expr} \rightarrow \text{Term}$

Done!

# Equivalence of grammars

Given two context-free grammars,  $G_1$  and  $G_2$ .  
Are they equivalent?

I.e., is  $L(G_1) = L(G_2)$ ?

# Equivalence of grammars

Given two context-free grammars,  $G_1$  and  $G_2$ .  
Are they equivalent?

I.e., is  $L(G_1) = L(G_2)$ ?

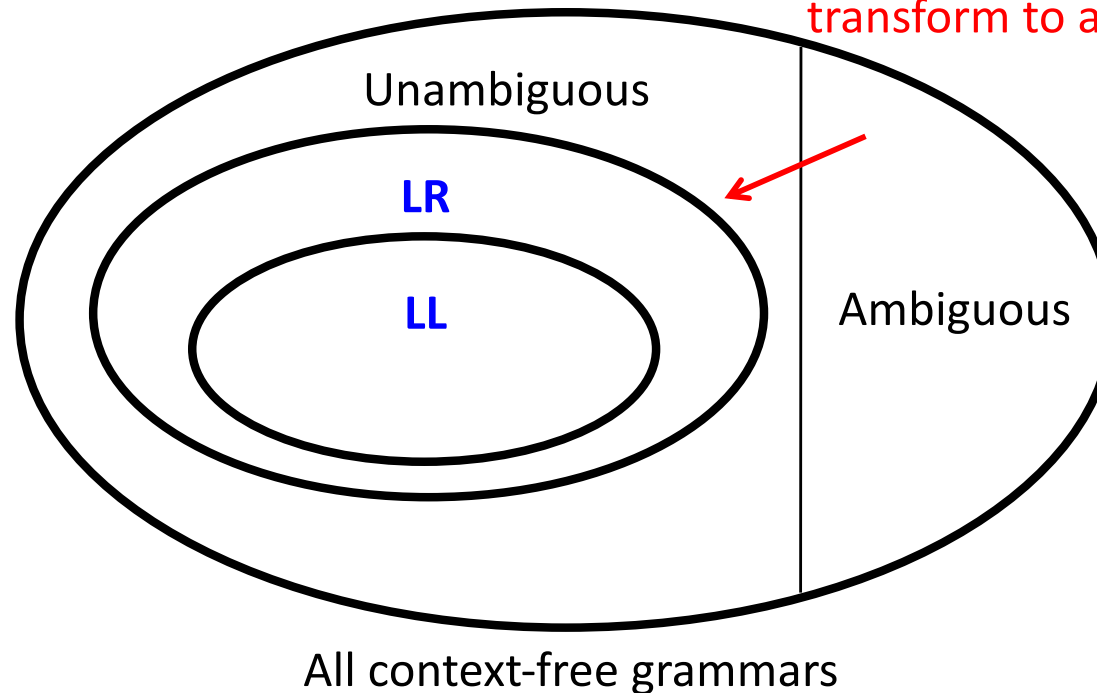
Undecidable problem:  
a general algorithm cannot be constructed.

We need to rely on our ingenuity to find out.  
(In the general case.)



# Space of context-free grammars

After eliminating typical ambiguities, there may still be work to do to transform to an LR or LL grammar.



**LL:**

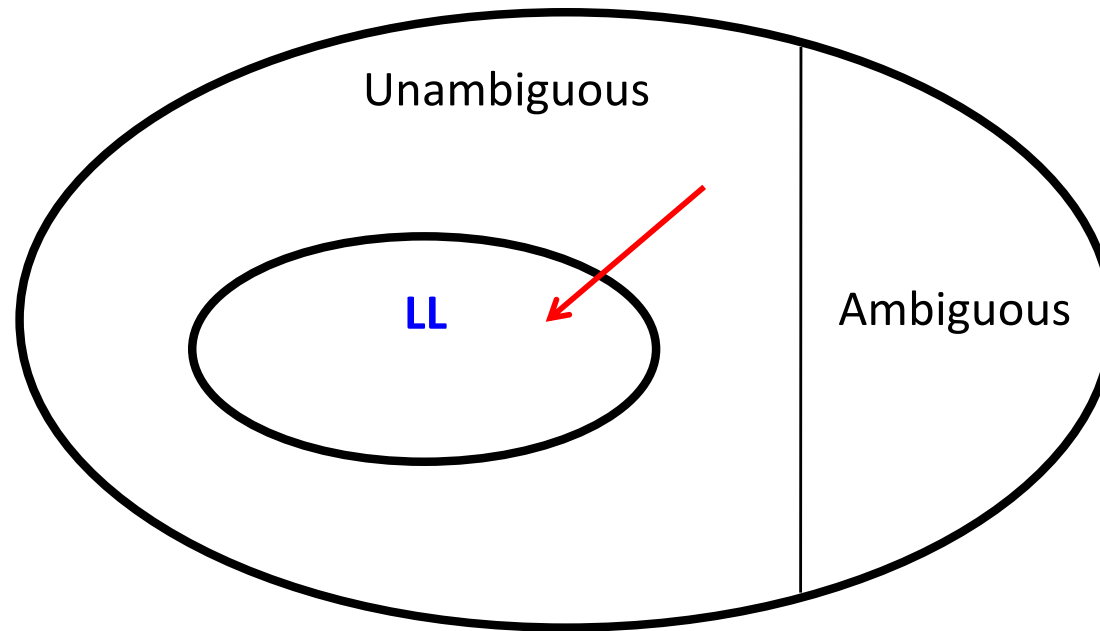
Builds tree top-down  
Simple to understand

**LR:**

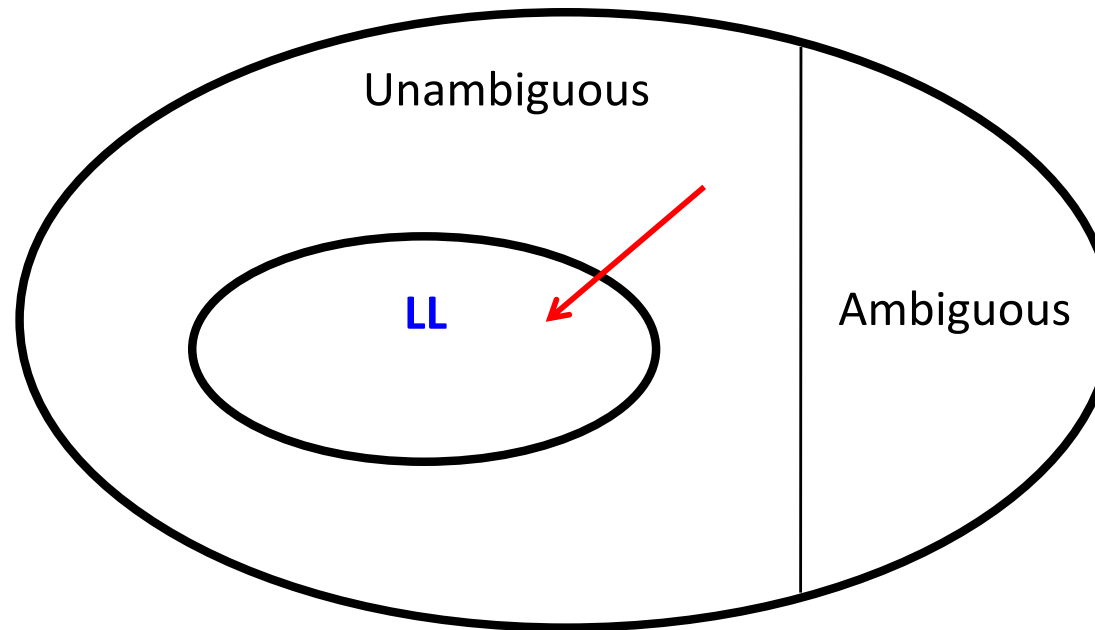
Builds tree bottom-up  
More powerful

# Adapting grammars to LL parsing

# Create equivalent LL grammar



# Create equivalent LL grammar



Typically, need to eliminate Left Recursion and Common Prefixes.  
(But this may not be enough.)

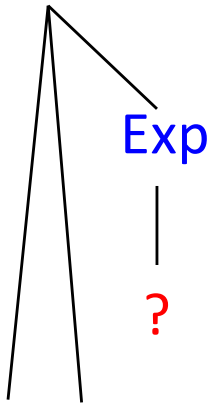
The parse trees will be different from the original desired ones.

Some work needed to build the desired ASTs anyway.

EBNF helps: relatively easy to build the desired AST.

# Recall: LL(1) parsing

Assign



? What node should be built?

ID = ID.ID; ID = ID.ID ( ID );

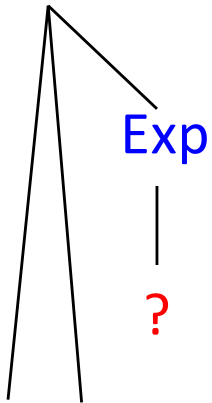


LL(1): decides to build the node after seeing the first token of its subtree.  
The tree is built top down.

```
Assign -> ID = Exp ;  
Exp -> Name Params | Name | ...  
Name -> ID ( . ID )*
```

# Recall: LL(1) parsing

Assign



Exp

?

What node should be built?

ID = ID.ID; ID = ID.ID ( ID );



LL(1): decides to build the node after seeing the first token of its subtree.  
The tree is built top down.

```
Assign -> ID = Exp ;  
Exp -> Name Params | Name | ...  
Name -> ID ( . ID )*
```

Common prefix!  
Cannot be handled by LL(1).  
This grammar is not even LL(k).

# Eliminating the common prefix

Rewrite to an equivalent grammar without the common prefix

Exp -> Name Params | Name

With common prefix - not LL(1)

# Eliminating the common prefix

Rewrite to an equivalent grammar without the common prefix

Exp  $\rightarrow$  Name Params | Name

With common prefix - not LL(1)

Exp  $\rightarrow$  Name OptParams  
OptParams  $\rightarrow$  Params |  $\epsilon$

Without common prefix - LL(1)

Eliminating a common prefix this way is called *left factoring*.



# Exercise

If two productions of *the same* nonterminal can derive a sentence starting in the same way, they share a *common prefix*.

G1:

$A \rightarrow s B$
$A \rightarrow s C$
$B \rightarrow t$
$C \rightarrow u$

G2:

$A \rightarrow B s$
$A \rightarrow B t$
$B \rightarrow u v$

G3:

$A \rightarrow s B$
$B \rightarrow s C$
$B \rightarrow t C$
$C \rightarrow u$

Which grammars have common prefix productions?  
What is the common prefix? Is the grammar LL(1), LL(2), ...?

# Solution

If two productions of a nonterminal can derive a sentence starting in the same way, they share a *common prefix*.

G1: 

$A \rightarrow s B$
$A \rightarrow s C$
$B \rightarrow t$
$C \rightarrow u$

 $A$  has two rules that can derive the prefix  $s$   
G1 is LL(2)

G2: 

$A \rightarrow B s$
$A \rightarrow B t$
$B \rightarrow u v$

 $A$  has two rules that can derive the prefix  $u v$   
G2 is LL(3)

G3: 

$A \rightarrow s B$
$B \rightarrow s C$
$B \rightarrow t C$
$C \rightarrow u$

 This is not a common prefix problem. The two rules that start the same cannot be derived from the same nonterminal.  
G3 is LL(1)

Which grammars have common prefix productions?  
What is the common prefix? Is the grammar LL(1), LL(2), ...?

# The common prefix can be indirect

G1:

```
A -> B
A -> C
A -> D
B -> t s
C -> t v
D -> x
```

G2:

```
A -> B s
A -> B t
B -> B u
B -> v
```

Which grammars have common prefix productions?

What is the common prefix?

Is the grammar LL(1), LL(2), ...?

# The common prefix can be indirect

G1:

```
A -> B
A -> C
A -> D
B -> t s
C -> t v
D -> x
```

A has two rules that can derive the prefix **t**

G1 is LL(2)

G2:

```
A -> B s
A -> B t
B -> B u
B -> v
```

A has two rules that can derive the prefix **v u\***

So, the prefix can become arbitrarily long.

G2 is not LL( $k$ ), no matter what  $k$  we use.

We need to rewrite the grammar, or use another parsing method than LL. (For example, LR has no problem with common prefixes)

Which grammars have common prefix productions?

What is the common prefix?

Is the grammar LL(1), LL(2), ...?

# Eliminating the common prefix

Rewrite to an equivalent grammar without the common prefix

A  $\rightarrow$  B  
A  $\rightarrow$  C  
B  $\rightarrow$  t s  
B  $\rightarrow$  x D  
B  $\rightarrow$  y  
C  $\rightarrow$  t v  
D  $\rightarrow$  B C

Indirect  
common  
prefix

# Eliminating the common prefix

Rewrite to an equivalent grammar without the common prefix

```
A -> B
A -> C
B -> t s
B -> x D
B -> y
C -> t v
D -> B C
```

Indirect  
common  
prefix

First, make the common prefix directly visible:

Substitute all B right-hand sides into the A -> B rule

We can't remove the B rules since B is used in other places.

Similarly for the A -> C rule

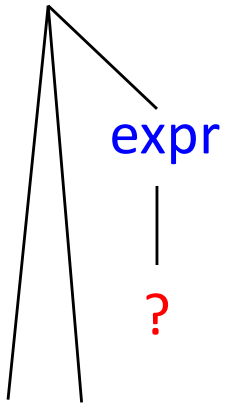
```
A -> t s
A -> x D
A -> y
B -> t s
B -> x D
B -> y
A -> t v
C -> t v
D -> B C
```

Direct  
common  
prefix

Then, eliminate the direct common prefix, as previously.

# Left recursion

assign



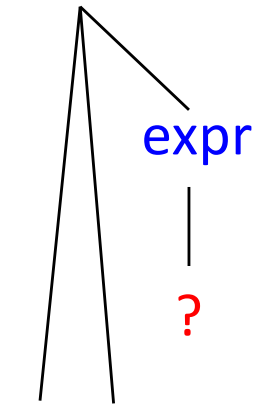
? What node should be built?

ID = ID + ID + ID ;  
↑

```
assign -> ID "=" expr ";"  
expr -> expr "+" term | term  
term -> ID
```

# Left recursion

assign



What node should be built?

ID = ID + ID + ID ;



```
assign -> ID "=" expr ";"  
expr -> expr "+" term | term  
term -> ID
```

The grammar is *left recursive*.

The grammar is not LL(k).

An LL parser would go into endless recursion.

(LR parsers can handle left recursion.)



# Dealing with left recursion in LL parsers

Method 1: Eliminate the left recursion  
(A bit cumbersome)

Left-recursive grammar. Not  
LL(k)

```
E -> E "+" T  
E -> T  
T -> ID
```

# Dealing with left recursion in LL parsers

## Method 1: Eliminate the left recursion (A bit cumbersome)

Left-recursive grammar. Not LL(k)

```
E -> E "+" T  
E -> T  
T -> ID
```

Rewrite to right-recursion! But there is now a common prefix! Still not LL(k).

```
E -> T "+" E  
E -> T  
T -> ID
```

Eliminate the common prefix. The grammar is now LL(1)

```
E -> T E'  
E' -> "+" E  
E' -> ε  
T -> ID
```

With a little work, it is possible to write code that builds a left-recursive AST, even if the parse is right-recursive.

# Dealing with left recursion in LL parsers

## Method 2: Rewrite to EBNF (Easy!)

Left-recursive grammar. Not  
LL(k)

```
E -> E "+" T  
E -> T  
T -> ID
```

# Dealing with left recursion in LL parsers

## Method 2: Rewrite to EBNF (Easy!)

Left-recursive grammar. Not LL(k)

```
E -> E "+" T  
E -> T  
T -> ID
```

Rewrite to EBNF!

```
E -> T ( "+" T ) *  
T -> ID
```

A left-recursive AST can easily be built during the iteration.

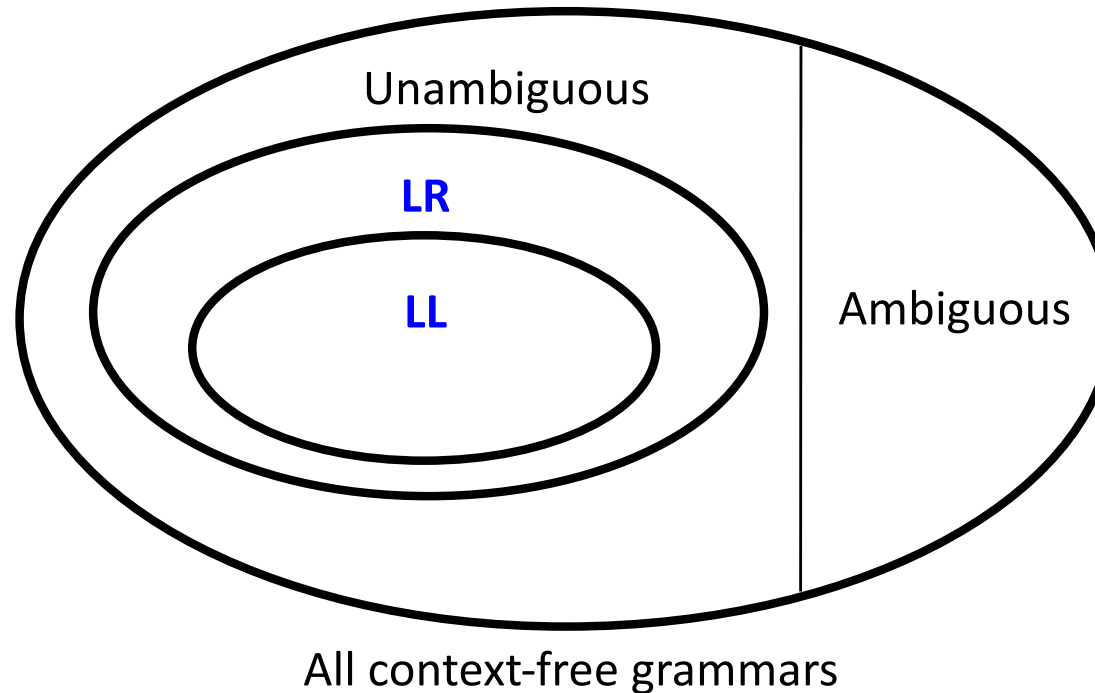
# Advice when using an LL-based parser generator

If the LL parser generator does not accept your grammar, the reason might be

- **Ambiguity** – usually eliminate it. In some cases, rule priority can be used.
- **Left recursion** – can you use EBNF instead? Otherwise, eliminate.
- **Common prefix** – is it limited? You can then use a local lookahead, for example 2. Otherwise, factor out the common prefix.

You might be able to solve the problem, but the grammar might become large and less readable.

# Different parsing algorithms



## **LL:**

Left-to-right scan

Leftmost derivation

Builds tree top-down

Simple to understand

Common prefixes and left recursion  
need to be eliminated

## **LR:**

Left-to-right scan

Rightmost derivation

Builds tree bottom-up

More powerful

Can handle common prefixes and  
left recursion

# LL( $k$ ) vs LR( $k$ )

	LL( $k$ )	LR( $k$ )
Parses input	Left-to-right	
Derivation	Leftmost	Rightmost
Lookahead	$k$ symbols	
Build the tree	top down	bottom up
Select rule	after seeing its first $k$ tokens	after seeing all its tokens, and an additional $k$ tokens
Left recursion	Cannot handle	Can handle!
Unlimited common prefix	Cannot handle	Can handle!
Can resolve some ambiguities by special disambiguation rules	Dangling else	Dangling else, associativity, priority
Error recovery	Difficult	Good algorithms exist
Implement by hand?	Possible. But better to use a generator.	Too complicated. Use a generator.

# Summary questions

- What does it mean for a grammar to be ambiguous?
- What does it mean for two grammars to be equivalent?
- Exemplify some common kinds of ambiguities.
- Exemplify how expression grammars with can be disambiguated.
- What is the "dangling else"-problem, and how can it be solved?
- When should we use canonical form, and when BNF or EBNF?
- Translate an example EBNF grammar to canonical form.
- Can we write an algorithm to check if two grammars are equivalent?
- What is a "common prefix"?
- Exemplify how a common prefix can be eliminated.
- What is "left factoring"?
- What is "left recursion"?
- Exemplify how left recursion can be eliminated in a grammar on canonical form.
- Exemplify how left recursion can be eliminated using EBNF.
- Can  $LL(k)$  parsing algorithms handle common prefixes and left recursion?
- Can  $LR(k)$  parsing algorithms handle common prefixes and left recursion?