



EDAF40/EDAN40: Functional Programming Introduction

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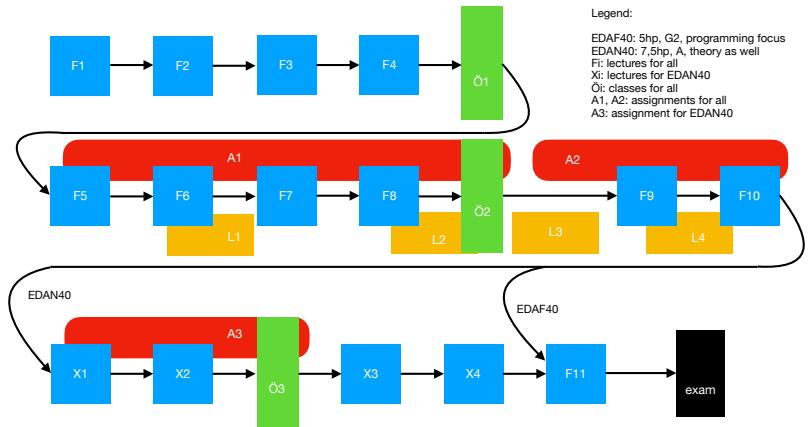


Administrativia

- Standard notification: 140/200h total compared with 20/28h with lecturers + 14/6 with TAs.
- Language-learning period in the beginning (syntax, basics).
- Two/Three not too tough programming assignments. (15, 10, 6 hrs)
- Kursombud (course representative) must be chosen. Today!
- Programming assignments verified by you, then machine and then teaching assistants (Christian Söderberg and Sven Gestegård Robertz, possibly more).
- Any problems (deadlines?) – please discuss **IN ADVANCE** with me!
- Slides based a lot on Lennart Andersson and Lennart Ohlsson's material. Thank you.



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Textbooks

- ① Graham Hutton, *Programming in Haskell*, 2nd ed., Cambridge University Press, 2016, ISBN 978-1316626221
- ② Bryan O'Sullivan, Don Stewart, and John Goerzen, *Real World Haskell*, O'Reilly Media, 2008, ISBN 0-596-51498-0
- ③ Miran Lipováča, *Learn You a Haskell for Great Good!*, No Starch Press, 2011, ISBN 1-59327-283-9
- ④ Paul Chiusano and Rúnar Bjarnason, *Functional Programming in Scala*. Manning Publications, 2014, ISBN: 9781617290657.
- ⑤ Simon Thompson, *Haskell - The Craft of Functional Programming*, 3rd edition, Addison-Wesley 2011, ISBN 0-201-88295-7

Software



- Glasgow Haskell Compiler, or `ghc`
- Interpreter is called `ghci`
- Currently in its version 7.10.3. (@ `login.student.lth.se`), or higher
- *.student.lth.se all run this version (please report issues)
- consider installing haskell-stack environment on your machine (<http://haskellstack.org>)

What is functional programming?



“Functional programming is so called because a program consists entirely of functions. [...] These functions are much like ordinary mathematical functions [...] defined by ordinary equations.”

(John Hughes)

Suggestions



- Read the assignment completely before you begin coding;
- Read the assignment text **after** the official announcement date;
- Complain to me or to a course student representative, if something does not work or is unclear;
- Check the course web;
- Do not mail `fp@cs.lth.se` unless you are filing in a **working** solution to an assignment;
- Do not mail `edan40@cs.lth.se` if you want to contact a human;
- Plan your time!
- Use our time (JM Mo 15.30-16.30, CS ..., SGR ...)!

A function



Let A and B be arbitrary sets.

Any subset of $A \times B$ will be called a *relation* from A to B .

A relation $R \subset A \times B$ is a *function* if and only if

$$\forall x \in A \forall y_1, y_2 \in B ((x, y_1) \in R \wedge (x, y_2) \in R) \rightarrow (y_1 = y_2)$$

A function



Our domain and range here: natural numbers

$$\begin{aligned} f_0 &= 1 \\ f_n &= n * f_{(n-1)} \end{aligned}$$

A function



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mathematical induction vs. computational recursion vs.
mathematical recursion

Equals for equals



If

$$\begin{aligned} f_0 &= 1 \\ f_n &= n * f_{(n-1)} \end{aligned}$$

then what is f_3 ?

Equals for equals



If

$$\begin{aligned} f_0 &= 1 \\ f_n &= n * f_{(n-1)} \end{aligned}$$

then what is f_3 ?

$$\begin{aligned} f_3 &= 3 * f_2 \\ &= 3 * 2 * f_1 \\ &= 6 * 1 * f_0 \\ &= 6 * 1 \\ &= 6 \end{aligned}$$

called also *rewrite semantics*

Imperative programming



Think like a computer:

```
public int f(int x) {
    int y = 1;
    for (int i=1; i<=x; i++) {
        y = y*i;
    }
    return y;
}
```

Then

$f(3) = y = y*i = ????$

The basic principle



NO ASSIGNMENTS!

not exactly, but the meaning is:

NO SIDE EFFECTS!

The basic principle



NO ASSIGNMENTS!

The problem with side effects



Example:

```
public int f(int x) {
    int t1 = g(x) + g(x);
    int t2 = 2*g(x);
    return t1-t2;
}
```

The problem with side effects



Example:

```
public int f(int x) {
    int t1 = g(x) + g(x);
    int t2 = 2*g(x);
    return t1-t2;
}
```

Then **of course**

$$f(x) = t1 - t2 = g(x) + g(x) - 2*g(x) = 0$$

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Example:

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public int f(int x) {
    int t1 = g(x) + g(x);
    int t2 = 2*g(x);
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}
```

Then **of course**

$$f(x) = t1 - t2 = g(x) + g(x) - 2*g(x) = 0$$

But suppose:

```
public int g(int x) {
    int y = input.nextInt();
    return y; }
```

The concept of a variable



Is a variable the name of a

memory cell

or the name of an

expression?

The core of functional programming



Functional programming

=

ordinary programming – assignments / side effects

It provides good support for

- higher order functions
- infinite data structures
- lazy evaluation

Recursion: The sum of a list



```
sum1 [] = 0
sum1 (x:xs) = x + (sum1 xs)
```

Note1: *recursion* is intimately connected to *computability*.

Note2: (x:xs) - a very important idiom in FP/Haskell.

Higher order functions



```
sum1 [] = 0
sum1 (x:xs) = x + (sum1 xs)
```

```
ackumulate f i [] = i
ackumulate f i (x:xs) = f x (ackumulate f i xs)
```

Higher order functions



```
sum1 [] = 0
sum1 (x:xs) = x + (sum1 xs)
```

```
ackumulate f i [] = i
ackumulate f i (x:xs) = f x (ackumulate f i xs)
```

```
sum2 = ackumulate (+) 0
```

Higher order functions



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sum1 [] = 0
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```
ackumulate f i [] = i
ackumulate f i (x:xs) = f x (ackumulate f i xs)
```

```
sum2 = ackumulate (+) 0
```

```
product2 = ackumulate (*) 1
anyTrue2 = ackumulate (||) False
allTrue2 = ackumulate (&&) True
```

Infinite lists



Primes computed with Eratosthenes sieve:

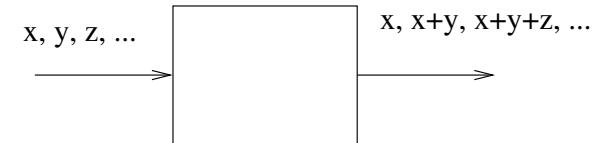
```
primes = sieve [2..]
  where
    sieve (n:ns) =
      n : sieve [ x | x <- ns, x `mod` n > 0 ]
```

Is this programming? Or just math?

Data flow programming



The running sums of a list of numbers:

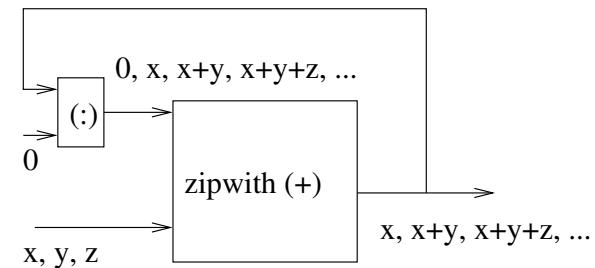


Running sums



```
runningSums xs = theSolution
  where
    theSolution = zipWith (+) xs (0:theSolution)
```

Data flow programming



Exact approximations



The Taylor series of the exponential function:

$$e^x = \sum_{i=0}^{\infty} \frac{x^i}{i!}$$

Exact approximations



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for example like a list of approximations:

```
eExp x = runningSums [ (x^i)/(fac i) | i <- [0..] ]
```