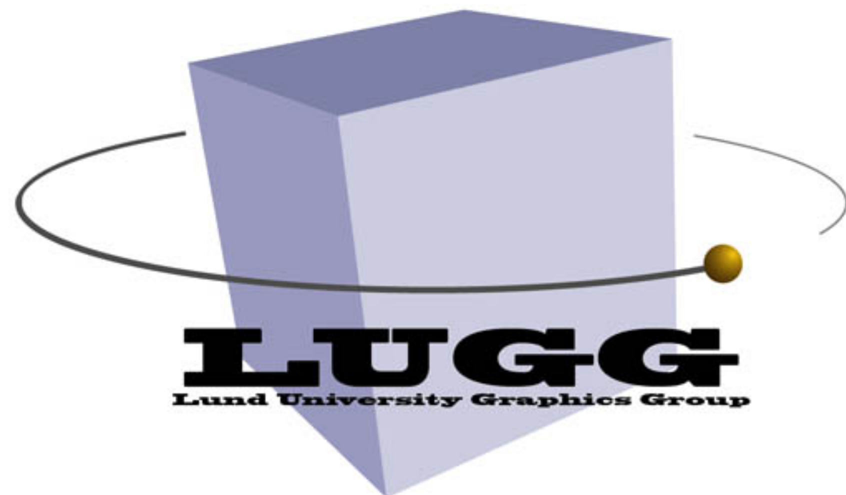
Real-Time Buffer Compression



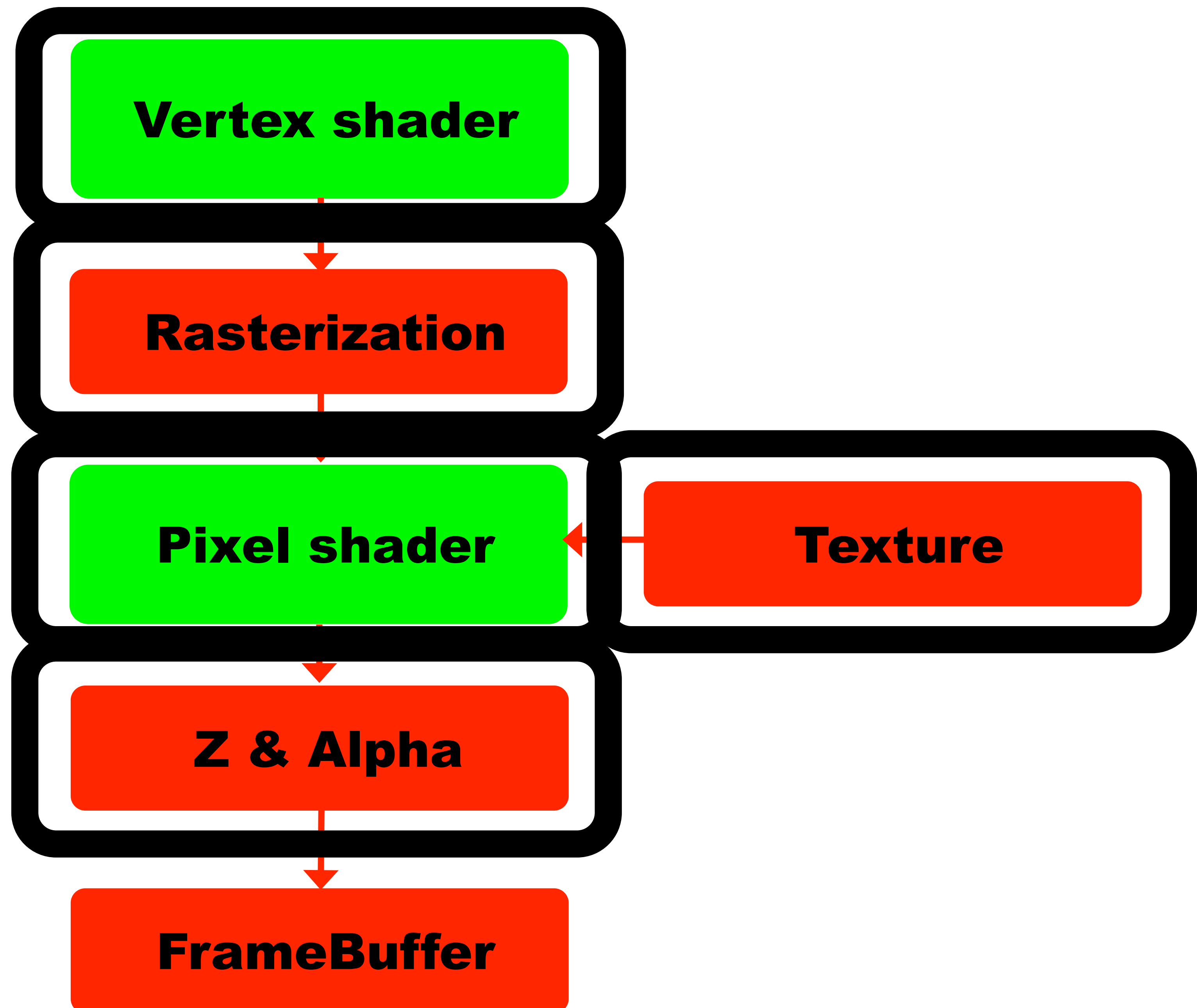
Michael Doggett
Department of Computer Science
Lund university

Project

- 3D graphics project
  - Implement 3D graphics algorithm(s)
  - C++/OpenGL(Lab2)/iOS/android/3D engine
  - Demo, Game
  - **Proposal - Long paragraph by next Thursday**
- More in the next lecture

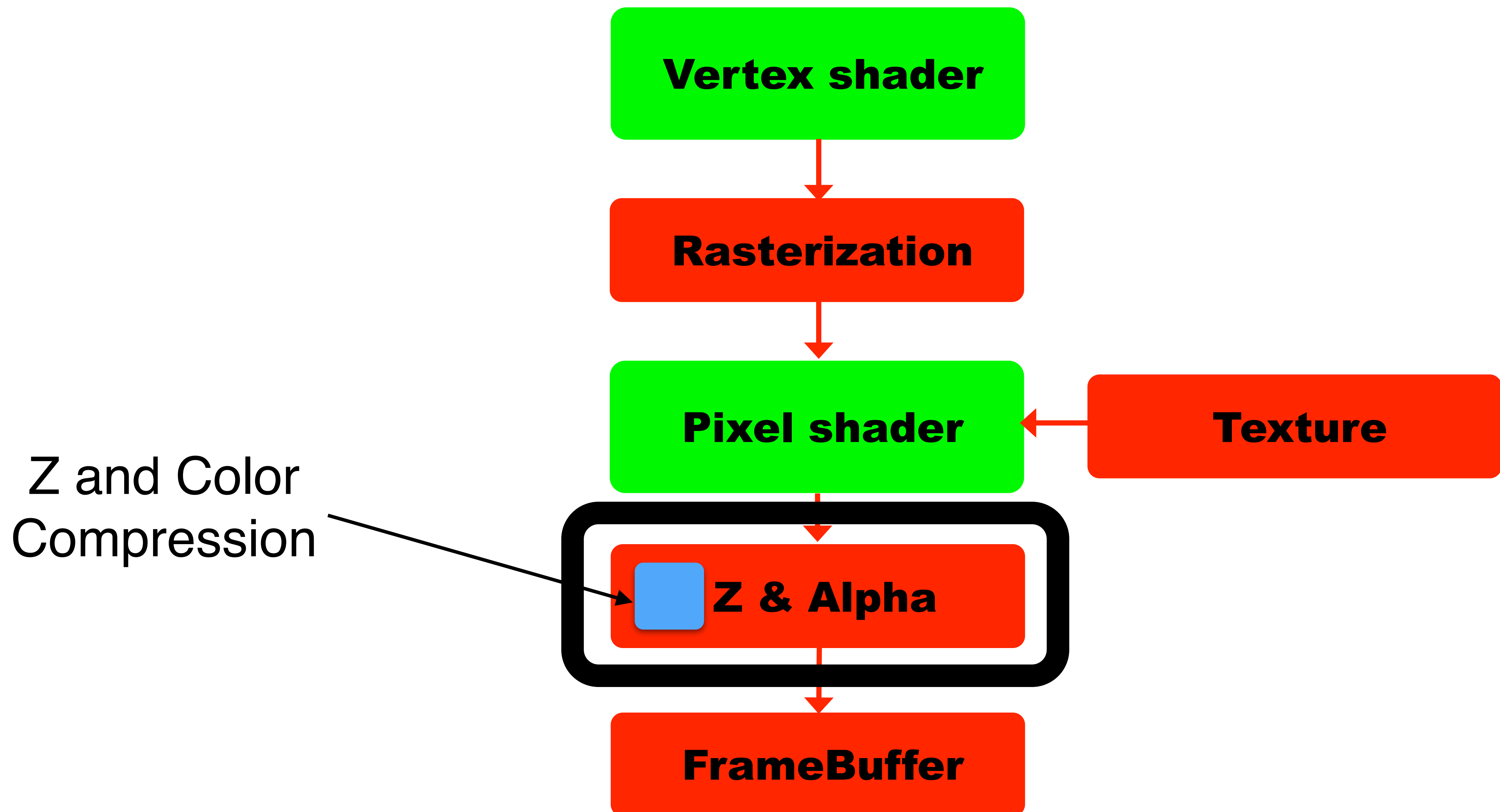
© 2010 Michael Doggett

Stages we have looked at so far



© 2010 Michael Doggett

Today's stages of the Graphics Pipeline



© 2010 Michael Doggett

Z & Alpha Performance

- Recall Memory Bandwidth determines performance
- Both units connect directly to memory
- Computation power of GPUs and CPUs increasing more rapidly than memory bandwidth
- Compression reduces the data before we transfer it

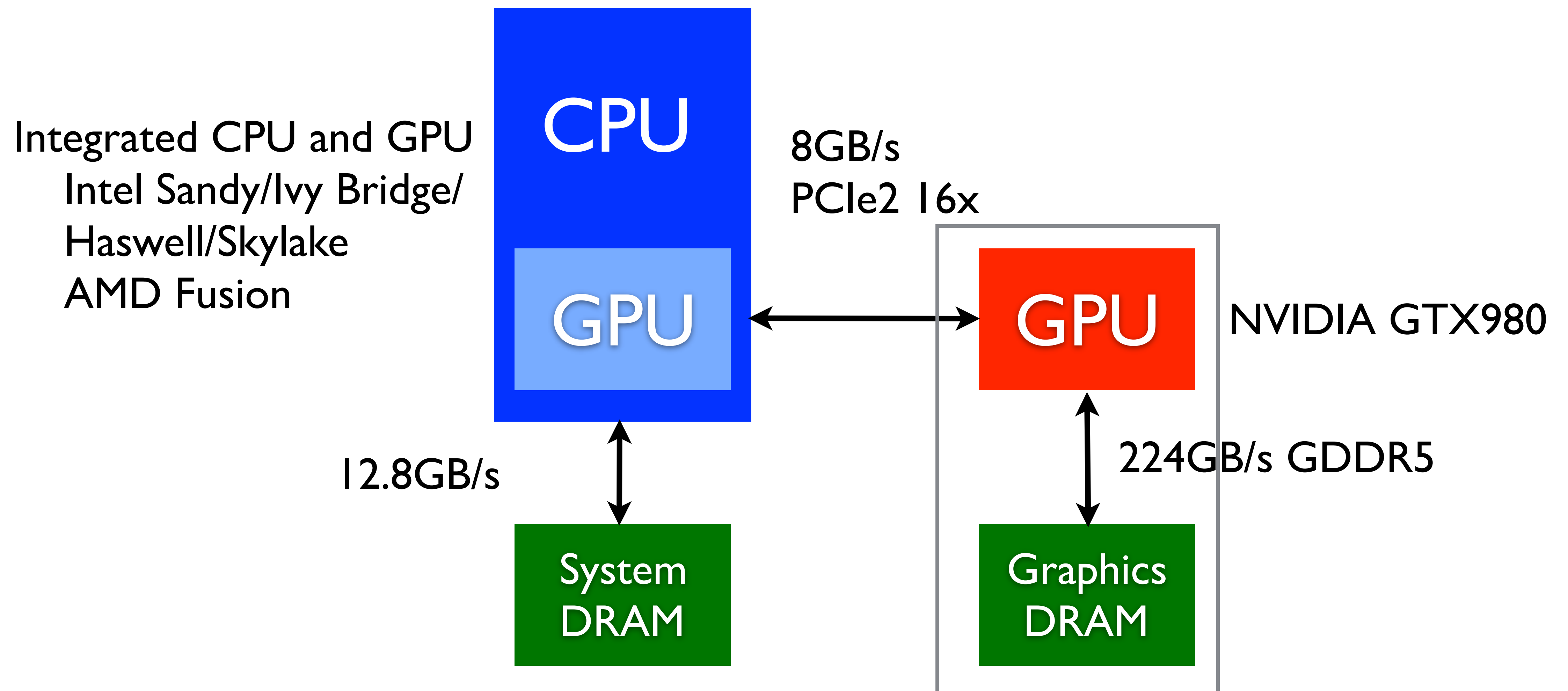
© 2012 Michael Doggett

DRAM overview

- Dynamic Random Access Memory
  - Must have power and be refreshed to maintain it
- Discrete GPU memory - Fast and reasonably priced
- Many different types : SDRAM, VRAM, SGRAM, etc.
- Multiple improvements of data transfer
  - DDR - sends data on both low-to-high and high-to-low
  - QDR, then GDDR (Graphics DDR) versions 2, 3, 4, 5
- HBM - High Bandwidth Memory
  - 3D-stacked DRAM

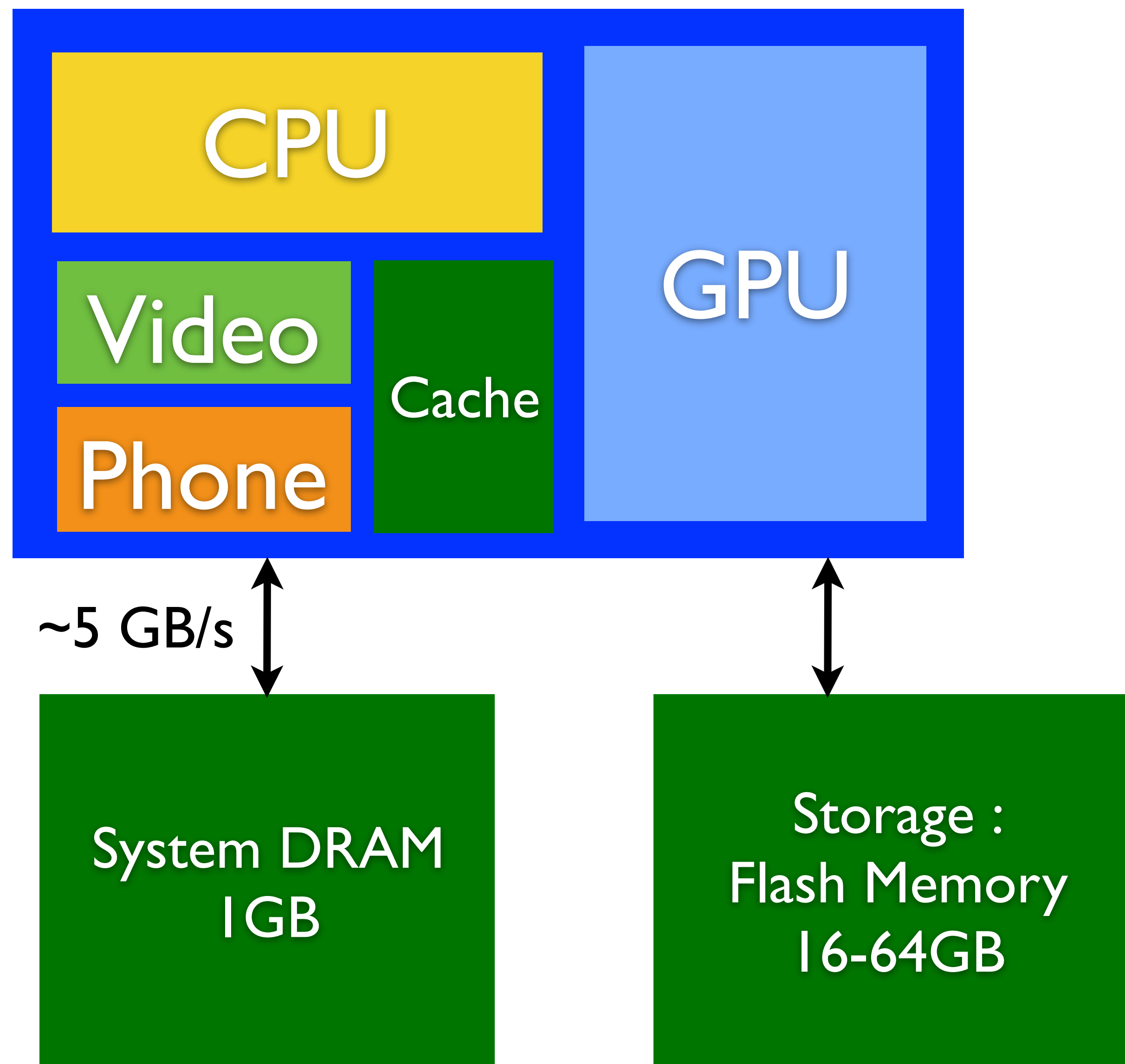
© mmxvi mcd

PC memory architecture



© 2012 Michael Doggett

Mobile Memory Architecture



Based on 2014 iPhone5S

© mmxiii mcd

source http://anandtech.com/show/7335/the-iphone-5s-review

Back to Graphics Hardware algorithms

© 2008 Tomas Akenine-Möller

Why depth buffer?

hardware Thus the only variation of interest here is Newell et al, an order of magnitude less "costly" and the brute-force approach which is already ridiculously expensive.

"A Characterization of Ten Hidden-Surface Algorithms", Ivan Sutherland, Robert Sproul, and Robert Schumacker (ACM Computing Surveys, March 1974)

[Slide courtesy of John Owens]

The "brute-force approach" is depth buffering (aka Z-buffering):
It won over sorting-polygon-methods because memory became ridiculously inexpensive...

© 2008 Tomas Akenine-Möller

Depth buffer bandwidth

- Still could be quite expensive!
- Zmin/Zmax-culling helped (previous lecture)

- Real-Time Buffer Compression can help reduce
  - Depth buffer bandwidth
  - Color buffer bandwidth
  - Other buffers...

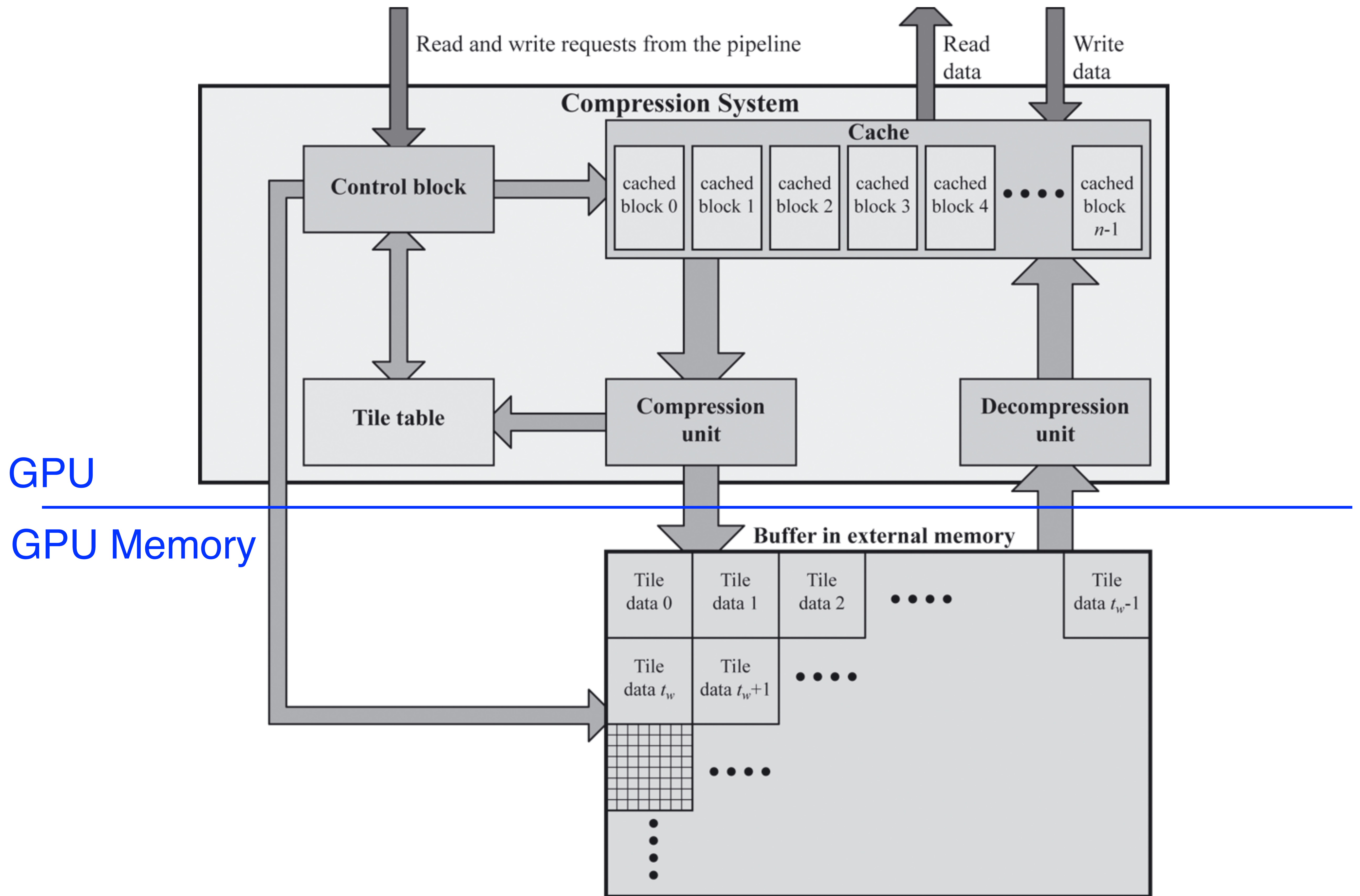
© 2008 Tomas Akenine-Möller

Real-Time Buffer Compression

- Techniques that are or ***may be*** used in GPUs...
- Basic idea:
  - –Lots of coherency (correlation) between pixels
  - –Use that to compress buffer info
  - –Send compressed buffer info over the bus
  - –Special hardware handles compression and decompression on-the-fly
  - –Must be lossless!!

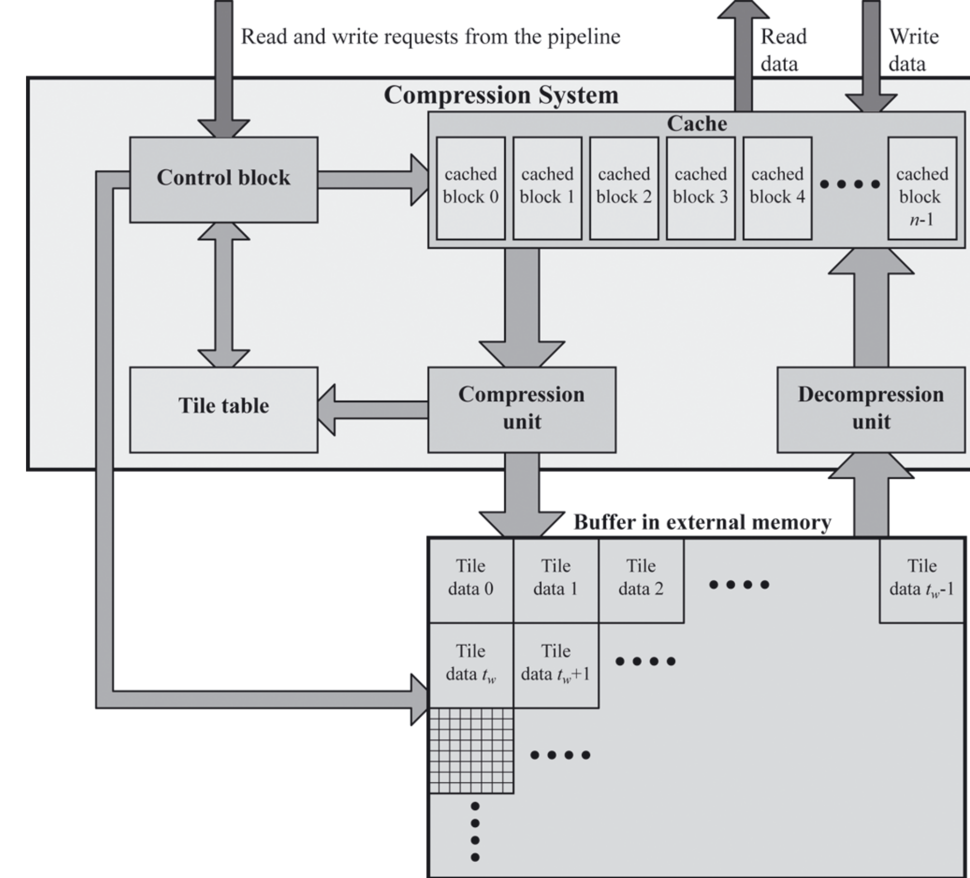
© 2008 Tomas Akenine-Möller

General Compression System



© 2008 Tomas Akenine-Möller

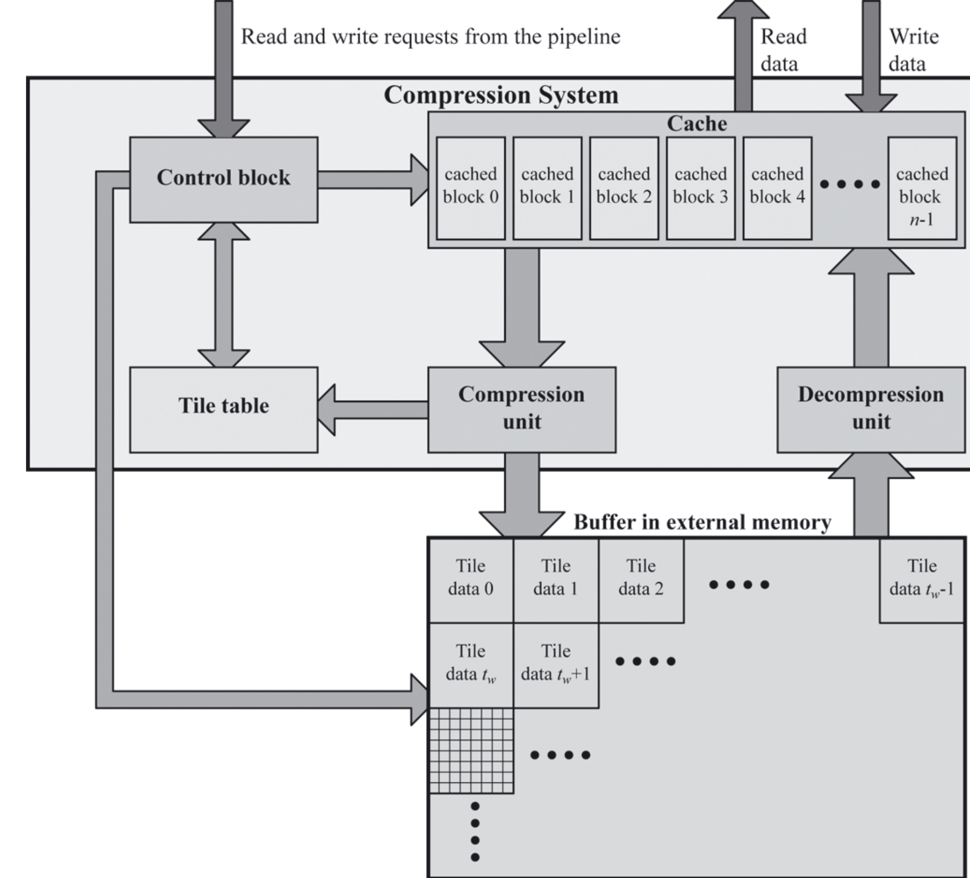
Compression System



- Works on a tile basis
  - –Eg 8x8 pixels at a time
- Cache is important!
  - –Do not want to decompress tile for every fragment that needs access to values in that tile
- Tile table store ”per-tile info”:
  - –E.g., which compression mode is used
  - –Example: 00 is uncompressed, 01 is compressed at 25%, 10 is at 50%, 11 is cleared
  - –Always needs one uncompressed mode as a fallback

© 2008 Tomas Akenine-Möller

Example



- Read request → ctrl block
- Checks cache
  - If there, deliver immediately
  - If not
    - Evict one tile from cache by attempting to compress info, and sending resulting representation, update tile table for that tile
    - Check tile table for requested tile, and
    - Read appropriate amount of bytes
    - Decompress (or send cleared info without reading, or in case of data being uncompressed, no decompression needed)
- Done

© 2008 Tomas Akenine-Möller

Dirty bit

- Each tile in cache has one bit for this
- When new info has been written to a tile in cache, set dirty bit=1
- When a tile in the cache needs to be evicted, check dirty bit
  - If =0, information in external memory is up to date → no need to write back!
  - If =1, attempt to compress, and send to external memory
- Saves a lot when no updates
  - Example: particle systems – do not write depth!

© 2008 Tomas Akenine-Möller

Depth buffer compression

- Hard to get accurate information about this
- Looking at patents we can extract some ideas
- Three techniques:
  - Depth offset compression
  - DPCM compression
  - Layered plane equation compression

© 2008 Tomas Akenine-Möller

Depth buffer compression

- Simplest buffer to compress
  - Highly coherent info (big triangles WRT tile size)
  - Depth is linear in screen space
- Depth cache and depth compression helps Zmax update for Zmax-culling
- Depths, $d(i,j)$ per tile,
  - i is in [0, w-1], j is in [0, h-1]
  - Min depth value is $00...00_b$ (all zeroes, e.g. 24 bits)
  - Max depth value is $11...11_b$ (all ones)
    - i.e., we can use integer math

© 2008 Tomas Akenine-Möller

Depth offset compression (1)

- Identify a set of reference values, r_k,
  - and compress each depth as an offset with respect to one reference value
- Easiest to only use two reference values
  - Use Zmin and Zmax of tile!
  - Rationale: we have two layers
    - One with depths close to Zmin and
    - one with depths close to Zmax

© 2008 Tomas Akenine-Möller

Visually, this means...

- Can encode if all z-values are in the gray regions

© 2008 Tomas Akenine-Möller

Depth offset compression (2)

- Use an offset range of $t=2^p$
- Can use offset, $o(i,j)$, per pixel as:

$$o(i,j) = \begin{cases} d(i,j) - z_{\min}, & \text{if } d(i,j) - z_{\min} < t, \\ z_{\max} - d(i,j), & \text{if } z_{\max} - d(i,j) < t. \end{cases}$$

- If at least one pixel, (i,j), cannot fulfil the above, the tile cannot be compressed!
- Info to store (if compression possible):
  - –Zmin and Zmax
  - –Plus wxh p-bit values

© 2008 Tomas Akenine-Möller

Depth offset compression (3)

- Example with following assumptions:
  - 8x8 pixels per tile
  - $t=2^8$ means 8 bits per offset, *o*
  - 24 bits depth → Zmin and Zmax has 24 bits each

- Storage (uncompressed: 8x8x3=192 bytes):
  - 1 bit per pixel to indicate whether offset to Zmin or Zmax → 8x8x1 bits= 8 bytes
  - Offsets: 8x8x8 bits= 64 bytes
  - Zmin & Zmax: 6 bytes (might be on-chip though)
  - Total: 8+64+6=78 bytes → 100*78/192=41% compression

© 2008 Tomas Akenine-Möller

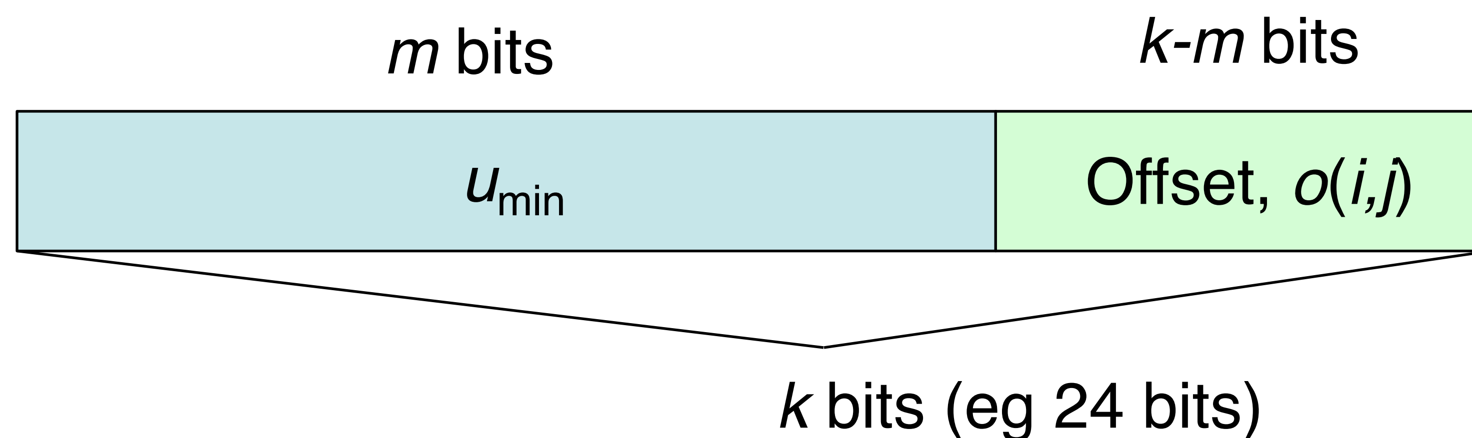
Less expensive implementation

- Only possible to compress if all depths in a tile are in gray regions
- There are some extensions to this that makes the hardware simpler!
- Make the offset computation inexpensive!
  – Currently costs an adder in HW

© 2008 Tomas Akenine-Möller

Inexpensive offsets...

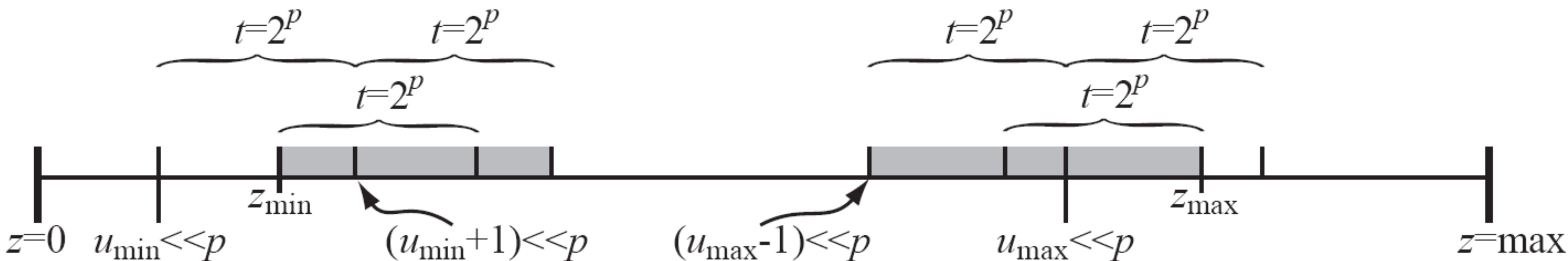
- Instead of storing exactly Zmin and Zmax, store only *m* most significant bits (MSBs)
  - Call these truncated values, u_{min} and u_{max}
- Offset is now simply the *k-m* least significant bits of depth (no add needed)



© 2008 Tomas Akenine-Möller

Disadvantage of cheap offsets, and a solution

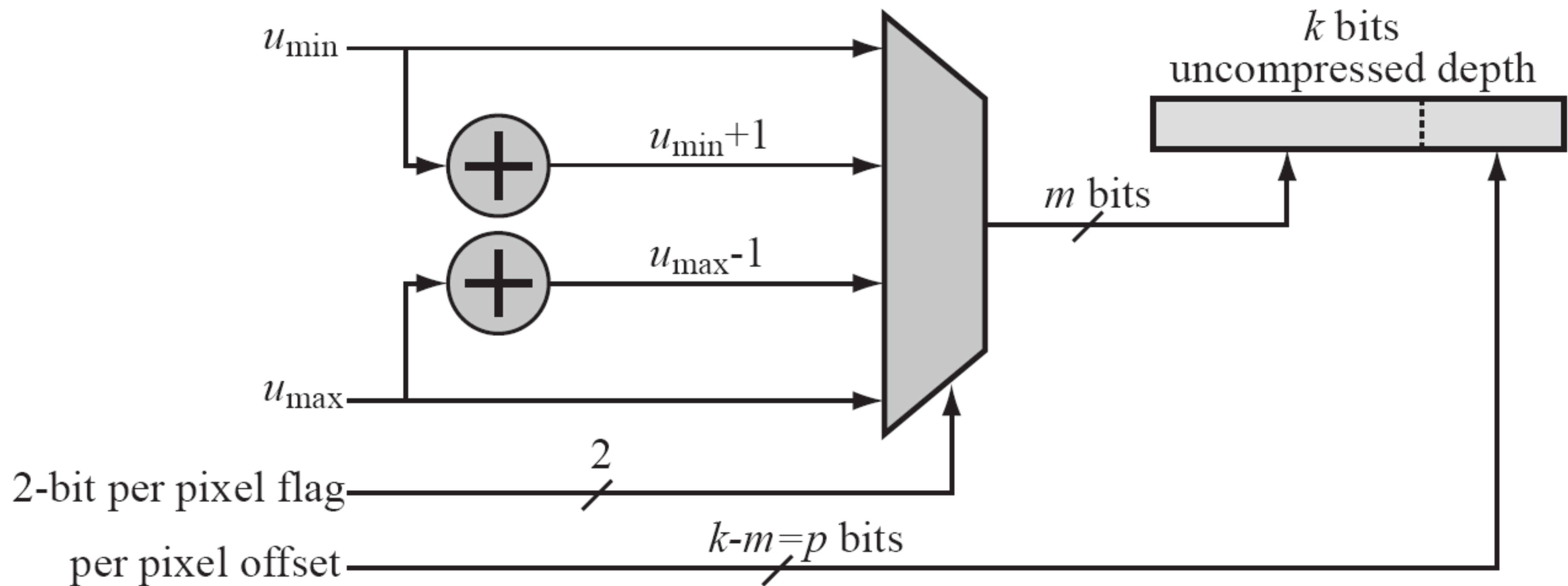
- Only values in dark gray area can be coded → loss of compressibility!
- Simple solution: use one more bit per pixel → four reference values:
  - u_{min}, u_{min}+1, u_{max}-1 and u_{max}



© 2008 Tomas Akenine-Möller

Decompression hardware

- Very simple



© 2008 Tomas Akenine-Möller

Compression ratio with inexpensive variant

- Slightly worse → 44% instead of 41%
- But, range of offsets is larger!
  - Best case: range is twice as large
  - Worst case: range is only one depth value larger
  - Average case: range is about 50% larger!
  - So more tiles can be compressed, but still costs more

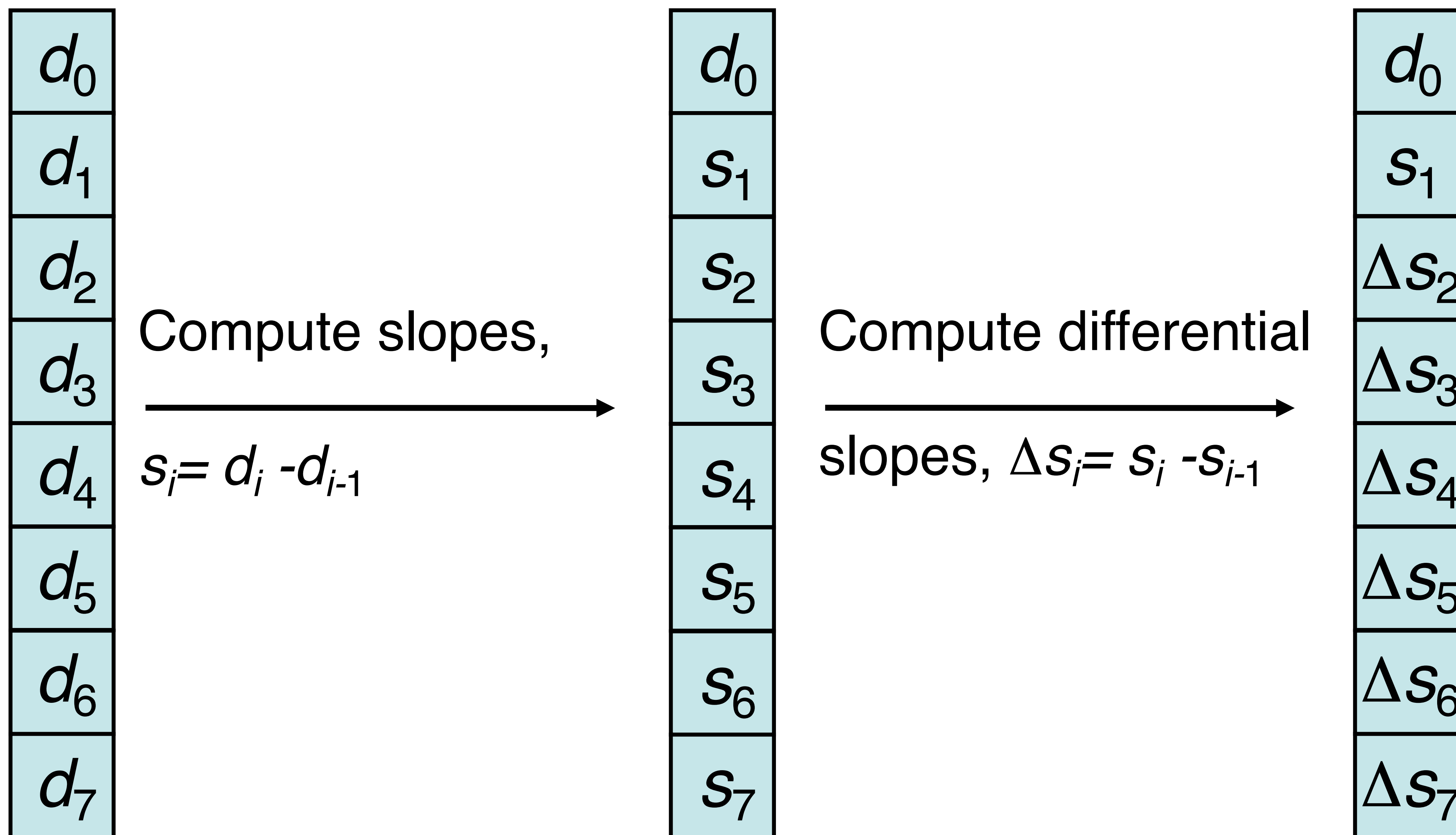
© 2008 Tomas Akenine-Möller

DPCM Compression

- DPCM=differential pulse code modulation
- Basic idea: we usually have linearly varying values in tile
  - Second derivative of linear function is zero!
  - However, we have discretized function, so need discretized ”second derivatives”

© 2008 Tomas Akenine-Möller

DPCM: Focus on one column of depths



- For linear functions, the Δs's will be close to 0
- Reconstruction is simple (next slide)

© 2008 Tomas Akenine-Möller

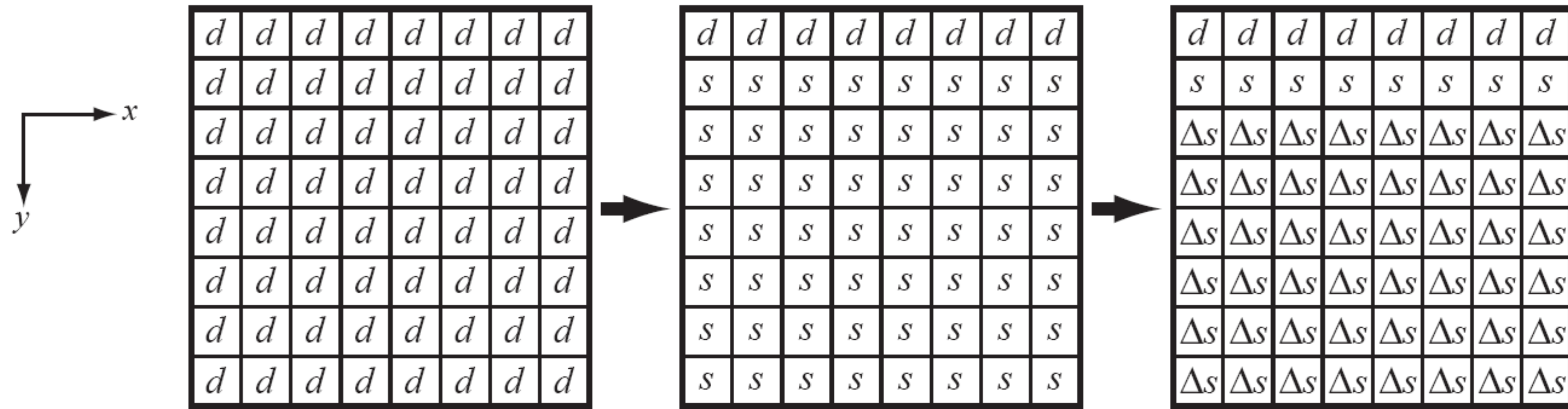
DPCM reconstruction

- From definition, we get: $d_i = d_{i-1} + s_i, \quad i \geq 1$
- Only s_1 is known, but: $s_i = s_{i-1} + \Delta s_i, \quad i \geq 2$

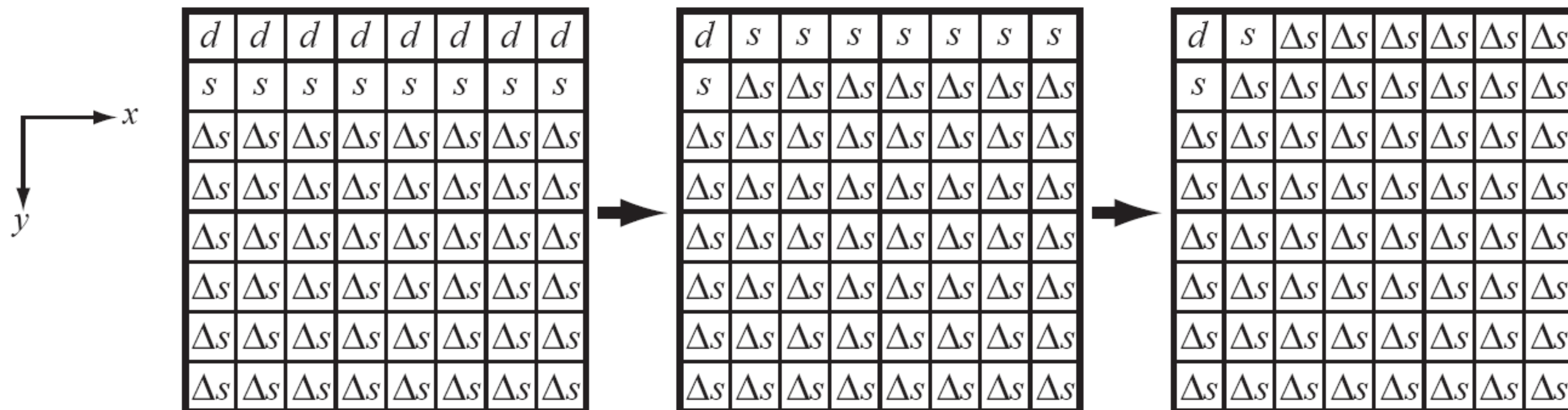
$$
\begin{array}{ll}
d_0, & \text{already known} \\
d_1 = d_0 + s_1, & s_1 \text{ already known} \\
d_2 = d_1 + s_2, & \text{where } s_2 = s_1 + \Delta s_2 \\
d_3 = d_2 + s_3, & \text{where } s_3 = s_2 + \Delta s_3 \\
\ldots \text{and so on} &
\end{array}
$$

© 2008 Tomas Akenine-Möller

Process each column independently



- Not ideal: still many d's and s's
- Process first two rows similarly →



© 2008 Tomas Akenine-Möller

DPCM: How to store this?

- One depth, d, two slopes, s, and 61 Δs
- The Δs are small, in [-1,+1] inside triangle
- Use two extra bits per pixel:
  - 00: add 0
  - 01: add +1
  - 10: add -1
  - 11: use as escape code to handle extraordinary cases...
- Best case compression (no escapes at all):
  - 24 bits + 25 +25 +8x8x2 ~= 25 bytes (13% ratio)
  - If a single triangle covers entire tile
    - Do not need the 11-escape case then though...

© 2008 Tomas Akenine-Möller

DPCM: common case

- Single column:
  - Depths:1,2,3,4,8,10,12,14
  - Slopes: 1,1,1,4,2,2,2
  - Diff slopes: 0,0,3,-2,0,0
- Two escape codes needed per column to change from one plane eq (tri) to another
  - Becomes expensive! 40% compression ratio
- Solution: encode from the top & down and from bottom & up
  - Store also where transition happens
  - Gives about 20% compression ratio!
  - Might be possible to use fewer bits per slope
  - Can only handle two plane equations per tile
  - Still does not use escape

© 2008 Tomas Akenine-Möller

Plane Equation Compression

- Each triangle can be represented as a plane
- For every triangle in a tile store the triangle's plane equation
  - Store one depth in center of tile, and an x-slope (dz/dx), and y-slope (dy/dz) across the tile
- For every pixel in the tile store an index to find the matching plane equation
- Works great for multisample!
- Random access
  - only decompress necessary pixels
- More info
  - [VanHook07] US Patent 7,242,400

© 2010 Michael Doggett

Plane Equation Compression

- Plane 0 : Z_c, x slope, y slope
- Plane 1 : Z_c, x slope, y slope
- Plane 2 : Z_c, x slope, y slope

- Plane Equations
  - 3 x (3 + 2 + 2)Bytes = 21Bytes
- Indexes
  - 64 x 2bits = 16Bytes
- Compressed
  - 37Bytes
- Uncompressed
  - 64 x 3Bytes = 192Bytes
- Compression ratio
  - 19%

1	1	0	0	0	0	0	0
1	1	1	0	0	0	0	0
1	1	1	1	0	0	0	0
1	1	1	1	2	0	0	0
1	1	1	2	2	2	0	0
1	1	2	2	2	2	2	0
1	2	2	2	2	2	2	2
2	2	2	2	2	2	2	2

© 2010 Michael Doggett

Color Buffer Compression

- Could use offset compression for R, G, and B separately (perhaps)
- Could use JPG's non-lossy algorithms
- Can do simple color compression for multi-sample anti-aliasing
- Can compress clear color
- Is generally very difficult due to restrictions
  - Cannot be lossy
  - Must decode very fast for alpha blending

© 2008 Tomas Akenine-Möller

Conclusion

- Compression reduces bandwidth further
  - Several options for depth
    - For more details read Notes chapter 7
  - Harder for color
  - Needs cache
  - Needs fallback for non-compressed mode

© 2008 Tomas Akenine-Möller

Next ...

- Today and Friday
  - Assignment 2 marking in Uranus
- Next lecture
  - Antialiasing
  - Texture Compression
  - Start **project**
- Next week :
  - GPU Architecture
  - Graphics Architecture and OpenCL

© 2010 Michael Doggett