Lecturer is Jonas.Skeppstedt@cs.lth.se with office E:2190
C page http://cs.lth.se/edaa25
Algimp page http://cs.lth.se/edaf15
You will get an account on a POWER8 machine (10 cores, 80 hardware threads)
You can work on other machines if you wish but performance measurements are to be done on this.
You can access it with ssh -Y user@power.cs.lth.se
Contents of the courses

- EDAA25 C Programming: odd lectures
- EDAF15 Algorithm implementation: all lectures but you can skip some from F7, F9, and F11 — I will tell you what.

F1  Introduction to C
F2  Project: linear and integer programming
F3  More C
F4  Instruction set architectures: POWER
F5  Types, conversions, and linkage
F6  Superscalar processors: POWER8
F7  Declarations and expressions
F8  Cache memories
F9  Statements and the C preprocessor
F10 Performance analysis
F11 The C Standard library
F12 Optimizing compilers
Sedgewick and Flajolet in "An Introduction to the Analysis of Algorithms":

The quality of the implementation and properties of compilers, machine architecture, and other major facets of the programming environment have dramatic effects on performance.
You will learn a methodology to maximise algorithm performance on a specific real machine (we will use an IBM POWER8 chip).

To write efficient code, you need competence in:
- Mathematics, algorithms and data structures (not the focus in this course)
- The C programming language and UNIX C programming tools
- Pipelined and superscalar processors
- Cache memories
- What optimizing compilers can do for you — and what you need to fix yourself

EDAA25: you will learn language details on C.
Contents Lecture 1

- The purpose of learning C
- Some simple C programs
Some views of C

- The *other* language for high-performance, FORTRAN, is mainly focussed on numerical computing and not for writing code eg for embedded systems, operating system kernels, or compilers.

- Very often other languages such as Clojure, Rust, Go, Scala, Haskell, Lisp, Prolog, Ada, Java, C++, Mathematica, or Matlab are preferable because they have many convenient features which enable faster program development.

- When performance in terms of memory usage and/or speed is *the* most important aspect, however, the programmer must have complete control over what is happening and then the overhead of many language features can lead to inferior performance.
Your lecturer’s relationship with C

• C is great but not ideal for *everything*. My default language since 1988. Just like Lisp and Prolog, it’s beautiful because it’s powerful *and* has few language features.

• I have written the second ISO validated C99 compiler (EDG was first).

• If I would manage a large software project with several million lines of code, I would still use C.

• I will not try to convince you that C ”is best” because there is no such thing as a best language.
Principles of the C Programming Language

- Trust the programmer
- Don’t prevent the programmer from doing what needs to be done
- Keep the language small and simple
- Provide only one way to do an operation
- Make it fast, even if it is not guaranteed to be portable
- Support international programming

*Update since the C99 version: Don’t trust the programmer.*
Writing a C program

```c
#include <stdio.h>

int main(int argc, char** argv)
{
    printf("hello, world\n");
    return 0;
}
```

- A Java method is called a function in C.
- A C program must have a `main` function.
- A C function must be declared before it is used.
The C preprocessor

- The command `#include <stdio.h>` reads a file with a declaration of `printf`.
- Commands in a C file which start with a hash, `#`, are performed by the C preprocessor before the compiler starts.
- You can run the preprocessor by typing `cpp`.
- The preprocessor can include files and deal with macros, eg `INT_MAX` is the largest number of type `int`.
- Notice that `cpp` knows nothing about C syntax.
In this course we will use the GNU C compiler, called gcc.

To compile one or more C files to make an executable program type:
```
gcc hello.c
```

The command gcc will first run cpp, then the C compiler, and then two more programs called an assembler and a link-editor.

Later in the course you will learn about assembler and the operating system course you can learn about link-editors.

For this course, gcc takes care of the link-editor and tells it to produce an executable file.
Running a C program

- By default the executable file (made by typing gcc hello.c) is called a.out.
- To execute it in Linux (or MacOS X, or another UNIX), type ./a.out.
- You can tell gcc that you want a certain name: gcc hello.c -o hello.
- Now you type ./hello.
Separate compilation

- If you have many big source code files, it is a waste of time to recompile all files every time.
- You can tell gcc to compile a file and produce a so called object file (has nothing to do with object-oriented programming).

```
gcc -c hello.c

gcc hello.o
```

- The above two lines are identical to `gcc hello.c` but useful if you have many files. The second line should then contain all .o files.
Example of I/O: scanf and printf

#include <stdio.h>
int main(int argc, char** argv)
{
    int a;
    float b;
    double c;

    scanf("%d %f %lf\n", &a, &b, &c);
    printf("%lf\n", a + b + c);
}

- %d for int, %f for float, and %lf for double.
- The program will read three numbers from input and print the sum.
In the call to the function `scanf`, we need `&` to tell the compiler that the variables should be modified by the called function.

This does not exist in Java. You cannot ask another method to modify a number passed as a parameter to the method.

Other useful format-specifiers include:
- `%x` for a hexnumber (base 16),
- `%s` for a string,
- `%c` for a char,
Writing to files in C

```c
#include <stdio.h>
int main(int argc, char** argv)
{
    int    a = 1;
    float  b = 2;
    double c = 3;
    FILE*  fp;

    fp = fopen("data.txt", "w"); // open the file for writing.
    fprintf(fp, "%d %f %lf\n", a, b, c);
    fclose(fp);
}
```

- This will create a new file on your hard disk.
#include <stdio.h>
int main(int argc, char** argv)
{
    int a;
    float b;
    double c;
    FILE* fp;

    fp = fopen("data.txt", "r"); // open the file for reading.
    fscanf(fp, "%d %f %lf\n", &a, &b, &c);
    fclose(fp);
}

- Note again the & since fscanf will modify the variables.
Three ways to make arrays in C

```c
#include <stdio.h>
#include <stdlib.h>

int size = 10;
int main(int argc, char** argv)
{  
    int a[10], n, i;
    int* b;
    int c[size];       // called a variable length array.
    sscanf(argv[1], "%d", &n);  // assumes program is run eg as $ ./a.out 10
    b = calloc(n, sizeof(int));  // like to Java’s b = new int[n];
    for (i = 0; i < n; i += 1)
        b[i] = i;  // play around with b as if it was an array
    free(b);
}
```
The `a` and `c` arrays are allocated with other local variables.

Note that `a` and `c` are "real" arrays.

On the other hand, `b` is like an array in Java for which you must allocate memory yourself. Use `new` in Java and `eg calloc` in C.

Java automatically takes care of deallocating the memory of objects.

In C you must do it yourself using `free`.

The variable `b` is not an array — it is a pointer.
int fun(int m, int n)
{
    int a[n];
    int b[m][n];
}

- Before C99 the above was illegal due to m and n are not constants.
- In C99 it is OK to write like that but only for local variables.
- Most C compilers still only support C89 and thus it may be wise to stick to that at least sometimes.
- Variable lengths arrays are only optional in C11.
C has no classes.

C has structs which are Java classes with everything public and no methods.

```c
struct s {
    // this s is a tag.
    int a;
    int b;
} s;  // this s is a variable identifier.
```

- Struct names have a so called tag which is a different namespace than variables and functions: so the above declares a `struct s` which is a type and a variable `s`.

- If we write `Link p` in Java we declare `p` to be a reference but not the object itself whereas `s` above is the real object, or data.
In Java we can declare a List class something like this:

```java
class List {
    List next;   // Next is a reference to another object.
    int a;
    int b;
}
```

- `next` above only holds the address of another object but `next` is not a List object itself. The list does not contain a list.
- Java let's you use pointers conveniently without giving you too much head ache.
- C does not.
We cannot write the following in C:

```c
struct list_t {
    struct list_t next;   // Compilation error!!
    int a;
    int b;
};
```

- It is impossible to allocate a list within the list!
- We really want to declare `next` to simply hold the address of a list object.
- In C this is done as: `struct list_t* next;` which makes `next` a pointer.
The following is correct in C:

```c
struct list_t {
    struct list_t* next;
    int a;
    int b;
};
```

After going into pointers in more detail we will see how to avoid typing `struct list_t` more than twice using `typedef`.
As you all know, your computer has something called memory. It is sufficient to view it as a huge array: `char memory[4294967296];`

It is preferable in the beginning to view it as: `int memory[1073741824];`

Forget about strings for the moment. Now our world consists only of ints.

As you know, a compiler translates a computer program into some kind of language which can be understood by a machine.

That has happened for the software in everybody’s mobile phone.
You will see more details about it later, but the C program which controls your phone is translated to commands which are numbers and can be represented as ints.

These ints are also put in the memory.

We can for instance put the instructions at the beginning of the array.

The instructions will occupy a large number of array elements.

No problem — our array is huge.
int x = 12;
int main()
{
    return x * 2;
}

- We also put the variable x in the memory.
- This program will have a few instructions for reading x from memory, multiplying with two, and returning the result.
- It is a good idea to put x after the instructions: next page
The array element where we have put a variable is called its **address**.

The instructions above are not written as integers but rather as commands to make them more readable.

An instruction is represented in memory as a number however.

It would be too complicated to demand that the hardware should read text such as **MUL** — it is easier is to build hardware if there simply is a number which means multiplication.

<table>
<thead>
<tr>
<th></th>
<th>READ from 3 into R</th>
<th>read the data in x from memory at address 3</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>MUL 2</td>
<td>R = R * 2</td>
</tr>
<tr>
<td>2</td>
<td>RETURN</td>
<td>return R</td>
</tr>
<tr>
<td>3</td>
<td>12</td>
<td>x lives here</td>
</tr>
</tbody>
</table>
When you call a function or method, all the local variables must be stored somewhere.

It is a convention to put them at the end of the memory array.

The local variables of the main function are put at the very end of the array.

When main calls a function, its local variables are put just before main’s.

In general, when a new function starts running, it puts its local variables at the last (highest index) unused memory array elements.

This works like a stack of plates: main is at the bottom and you put newly called functions on the plate at the top.
```c
int main()
{
    int x = 12;
    return f(x);
}

int g(int a)
{
    int b = a+1;
    return g(b+2);
}
```

<table>
<thead>
<tr>
<th>Address</th>
<th>Value</th>
<th>Info</th>
</tr>
</thead>
<tbody>
<tr>
<td>1073741817</td>
<td>15</td>
<td>a in g lives here.</td>
</tr>
<tr>
<td>1073741818</td>
<td></td>
<td>return address from g is here.</td>
</tr>
<tr>
<td>1073741819</td>
<td>13</td>
<td>b in f lives here.</td>
</tr>
<tr>
<td>1073741820</td>
<td>12</td>
<td>a in f lives here.</td>
</tr>
<tr>
<td>1073741821</td>
<td></td>
<td>return address from f is here.</td>
</tr>
<tr>
<td>1073741822</td>
<td>12</td>
<td>x in main lives here.</td>
</tr>
<tr>
<td>1073741823</td>
<td></td>
<td>return address from main is here.</td>
</tr>
</tbody>
</table>

- When a function returns, it deallocates its memory space.
- This is managed by the compiler which uses a register for holding the current free memory index, called the **stack pointer**.
int x = 12;
int *p;
int main()
{
    p = &x;
    *p = 13;
    return x * 2;
}

- A pointer is just a variable and it can hold the address of another variable.
- When `p` points to `x`, writing `*p` accesses `x`. 
## Memory layout

<table>
<thead>
<tr>
<th>instruction/data</th>
<th>Java</th>
<th>comment</th>
</tr>
</thead>
<tbody>
<tr>
<td>0    STORE 6 at 7</td>
<td>MEMORY[7] = 6</td>
<td>&amp;x is put in element 7, ie p</td>
</tr>
<tr>
<td>1    READ from 7 into R</td>
<td>R = MEMORY[7]</td>
<td>read data in p: R=6</td>
</tr>
<tr>
<td>2    STORE 13 at R</td>
<td>MEMORY[R] = 13</td>
<td>*p = 13</td>
</tr>
<tr>
<td>3    READ 6 into R</td>
<td>R = MEMORY[6]</td>
<td>fetch the value of x</td>
</tr>
<tr>
<td>4    MUL 2</td>
<td>R = R * 2</td>
<td>multiply x and R</td>
</tr>
<tr>
<td>5    RETURN</td>
<td>return R</td>
<td></td>
</tr>
<tr>
<td>6    12</td>
<td></td>
<td>x lives here</td>
</tr>
<tr>
<td>7    0</td>
<td></td>
<td>p lives here</td>
</tr>
</tbody>
</table>
In Java, you have used pointers all the time, but they are called object references.

Suppose you have `Link p`, then `p` is a pointer.

In Java, pointers can only point at objects.

The address of some object is, as you might know, the location in memory where that object lives, ie just an integer number.

In Java, `new` returns the address of a newly created object.

In C, `new` is a normal function and is called `malloc`. 
More about pointers

- In C, but not in Java, the programmer can ask for the address of almost anything and thus get a pointer to that object (or function).
- To change the value of a variable in a function, you need to pass the address of the variable as a parameter to the function:

```c
void f(int* a)
{
    *a = 12;
}

void g()
{
    int b;
    f(&b);
}
```
If the type of the variable is a pointer, then you will need two stars:

```c
void f(int** a)
{
    *a = NULL;
}

void g()
{
    int* b;
    f(&b);
}
```
More about pointers

- To return multiple values in Java, you create and return an extra object.
- Option 1 in C: use a plain struct which is allocated on the stack.
- Option 2 in C: Pass additional arguments as pointers (preferable).

```c
struct s f()
{
    struct s a;
    a.x = ...;
    a.y = ...;
    a.u = ...;
    return a;
}

void g(int* x, int* y, int* u)
{
    *x = ...;
    *y = ...;
    *u = ...;
}
```
Arrays vs Pointers

- Despite common belief, arrays and pointers are not equivalent.
- An array declares storage for a number of elements, except when it is a function parameter:

```c
int fun(int a[], int b[12], int c[3][4]);
int fun(int *a, int *b, int (*c)[4]);
int main()
{
    int   x, y[12], z[4];
    fun(&x, y, &z);  // valid.
}
```

- The compiler changes the first `[]` to `*` for array parameters.
- Array parameters are not arrays. They are pointers.
- Doing so avoids copying large arrays in function calls.
C has row-major matrix memory layout

```c
int c[3][4] = { { 1, 2, 3, 4}, { 5, 6 }, { 7 } }; int i, j; for (i = 0; i < 3; i++)    for (j = 0; j < 4; j++)        x += c[i][j];
```

- In a two-dimensional array, one row is layed out in memory at a time, ie row-major.
- Could also be called "rightmost index varies fastest".
int fun(int c[3][4])
{
    printf("%zu %zu\n", sizeof c, sizeof c[0]);
}

- If the output is "8 16", what conclusions can we draw about the size of a pointer and the size of an int?
- Answer: an pointer is eight bytes and an int is four bytes.
- The variable c in the function is simply a pointer: int (*c)[4].
• a[i] is represented as *(a+i) internally in the compiler.

    int main()
    {
        int a[10], *p, i = 3;

        /* the following are equivalent: */

        i[a];
        a[i];
        p = a; p[i]; i[p];
        p = a+i; 0[p]; p[0]; *p;
    }
Variables with static storage duration (globals, static).

Stack variables.

alloca(size_t size); normally adjusts the stack-pointer (no need for free).

Allocate from the heap with malloc/calloc/realloc. Must free explicitly.

Java has (1) for some static data, (2) for scalars and (4) for objects.
Memory allocation is extremely important to master.
Doing it right will lead to a much faster application.
Doing it wrong will certainly introduce painful bugs.
Memory allocation bugs can show up after you think your code works!
However, it is actually quite simple.
As you all know, Java uses garbage collection and this can make life easier but sometimes you need the exact control.
It is possible to have memory leaks in Java as well: if you accidentally keep the last reference to the root of a tree then the entire tree must be preserved.
A variable declared outside all functions is global and is visible from others source files.

Global variables are automatically initialised with zeros.

Explicit initialisers can also specify values other than zero.

Think of global variables as static variables in Java (ie static in a class).

Often it is best to avoid global variables due to:
- Compilers are not good at putting them in registers.
- They sometimes make it more difficult to understand the program.
Precede the type specifier with the **static** storage-class specifier.

For example: `static int a;`

Static variables are initialised with zeros (unless initialisers override that).

A static variable is only visible in the scope where it was declared, i.e., in a file, in a function, or in a particular block in a function.

A function can also be declared static making its visibility is limited to the file.

Unless a variable or function must be "exported", you **should** declare it static!
Stack allocated variables are usually allocated registers by the compiler, hence accessing them becomes more efficient.

Allocating large variables on the stack can cause stack overflow.

Returning and using the address, or contents, of a stack-allocated variable is an error.

Parameters are stack variables and at function calls they are given values, as if by assignment (except for arrays; which really are no arrays as parameters).

Passing a huge struct as a parameter will on some platforms result in copying the struct and on others result in passing the address (and letting the called function do the copying only if needed, to preserve the as if behaviour).
void f()
{
    struct { int a, char b, int c[100]; } s = { 1 };
    if (s.b != 0) abort();
}

- Initialising a part of a stack variable ensures that the rest become initialised to zero (NB: only variables, ie not "padding" between variables in eg a struct).
- Above only s.a is initialised but b and c also becomes initialised, with zero.
- There is typically padding space between b and c and it is not defined whether that is initialised to anything.
To allocate extra memory on the stack, use `alloca`.
No need to free such memory explicitly.
Use only if function returns "sufficiently often".
Problem 1: `alloca` is not standard.
Problem 2: if no memory is available, `NULL` is not returned (as for `malloc/calloc`).
Somewhat bad reputation, but nevertheless used.
Much more efficient than `malloc/calloc`.
To allocate memory for lists, trees, etc, use `malloc` or `calloc` (or `alloca`).

- `malloc` takes one parameter: number of bytes to allocate.
- `calloc` takes two parameters: number of items and size of each item.
- `calloc` returns memory initialised to zero; `malloc` leaves it undefined.
- Potential problem with using `malloc` (avoided with `calloc`):

```c
typedef struct { int a, c, d; } info_t;
info_t* info = malloc(sizeof *info);
info->a = ...; info->c = ...; info->d = ...;

Then add b to `info_t` but forget to initialise it.
```
Allocating memory with `malloc/calloc` (or `new` in Java) takes time.

Try to avoid allocation eg by reusing data of old objects, using a freelist:

- When allocating an object, see if there is something on the freelist first.
- When deallocating an object, put it first in the freelist instead of returning the memory to `free`.
- Sometimes the freelist grows too much should be deallocated.

There are numerous implementations of `malloc/calloc/free` and you might want to try different ones, or write your own.

Using data already freed is illegal —actually, it is illegal to even look at the address! (although nothing happens on most systems).
Use the `sizeof` operator when requesting memory.

The `sizeof` operator either takes a type or an expression as operand:
- `int* p; ..... p = calloc(n, sizeof(int));` (type in parenthesis)
- `int* p; ..... p = calloc(n, sizeof *p);`

The latter is better because if somebody changes the type, the compiler will automatically calculate the correct value for `sizeof`, otherwise if the type is changed from `int` to `long long` and the operand of `sizeof` is not also changed to `long long` someone (eg you) will suffer from a painful bug at some point in the future.