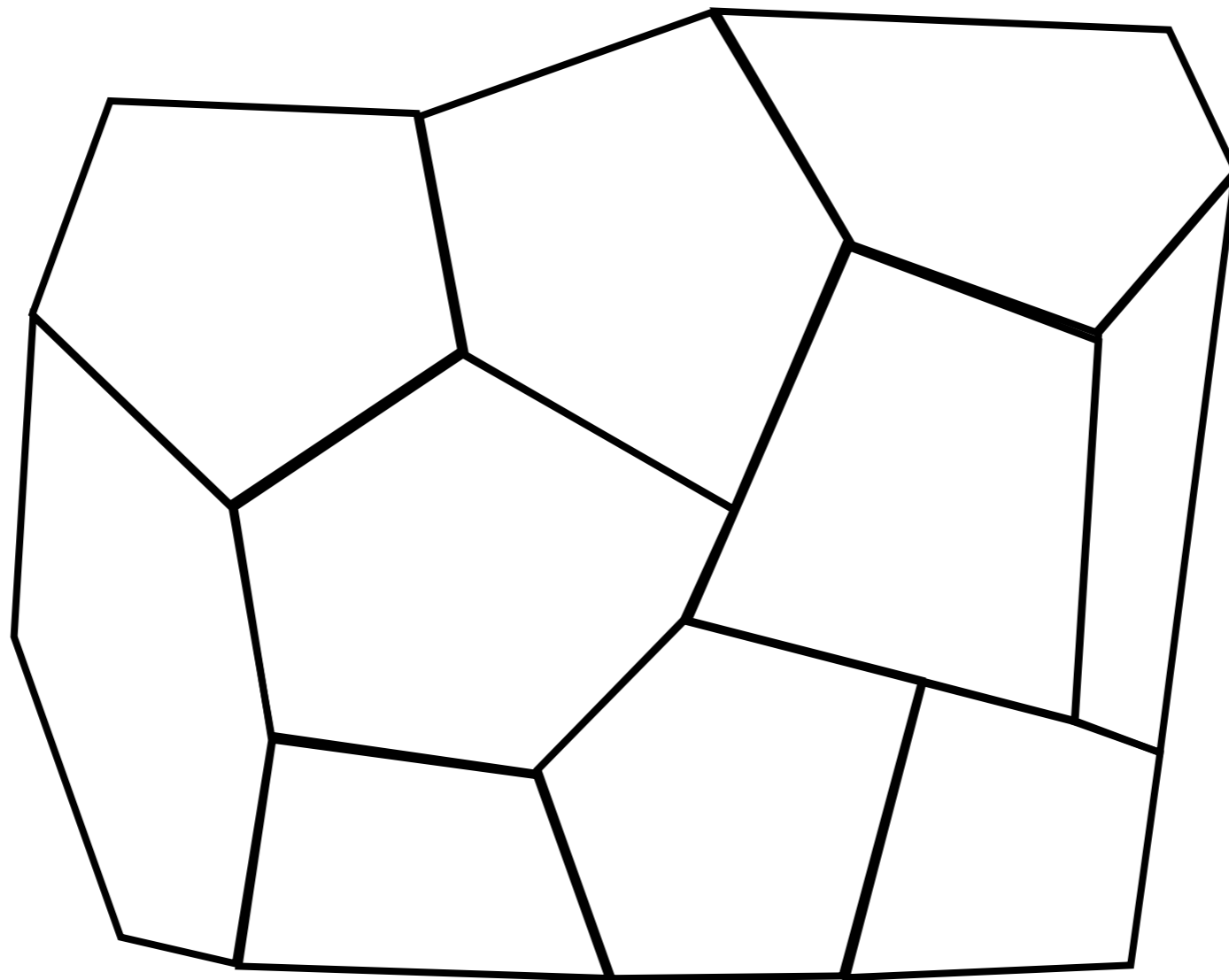


Graph Coloring

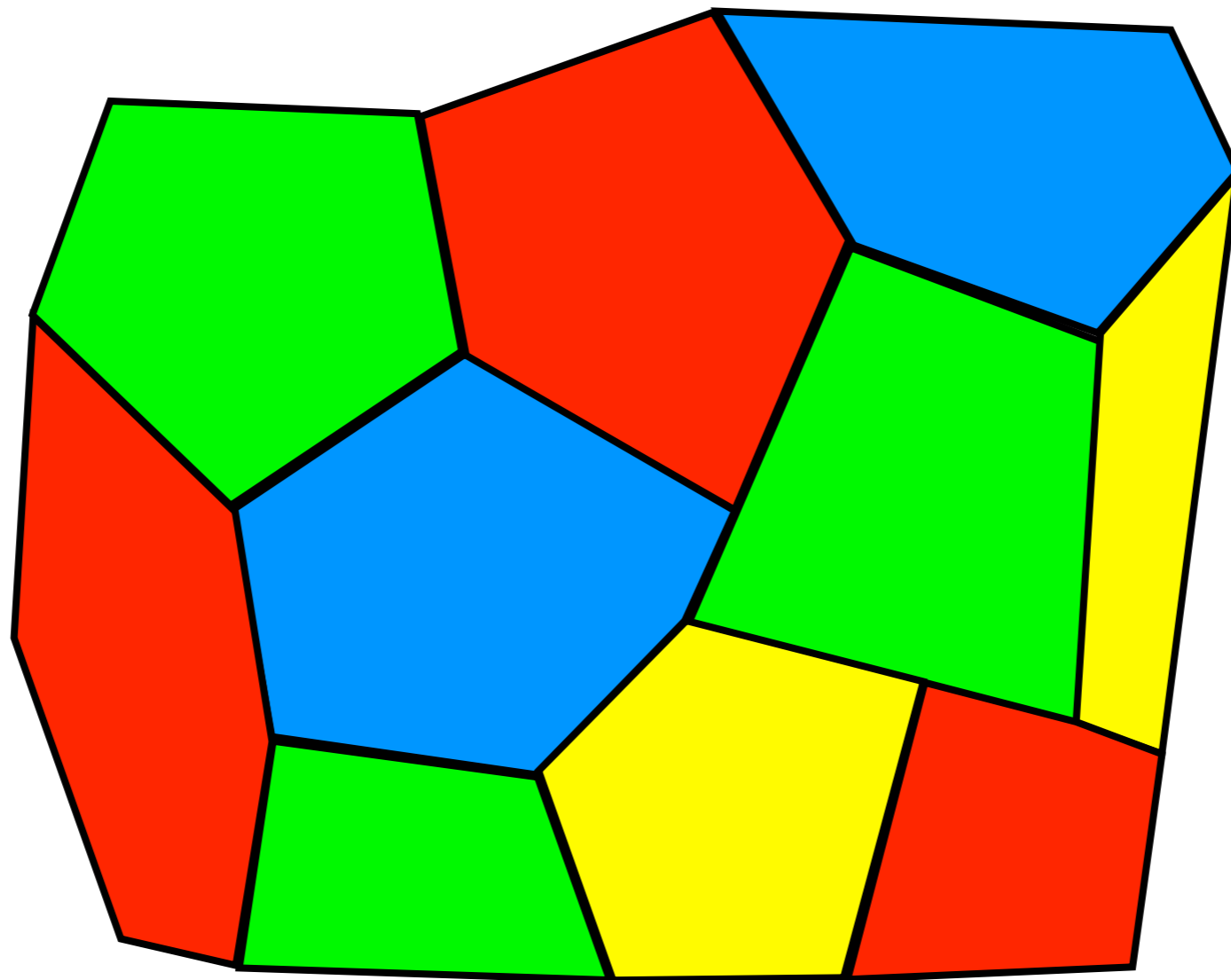
Andreas Björklund

Coloring a Map

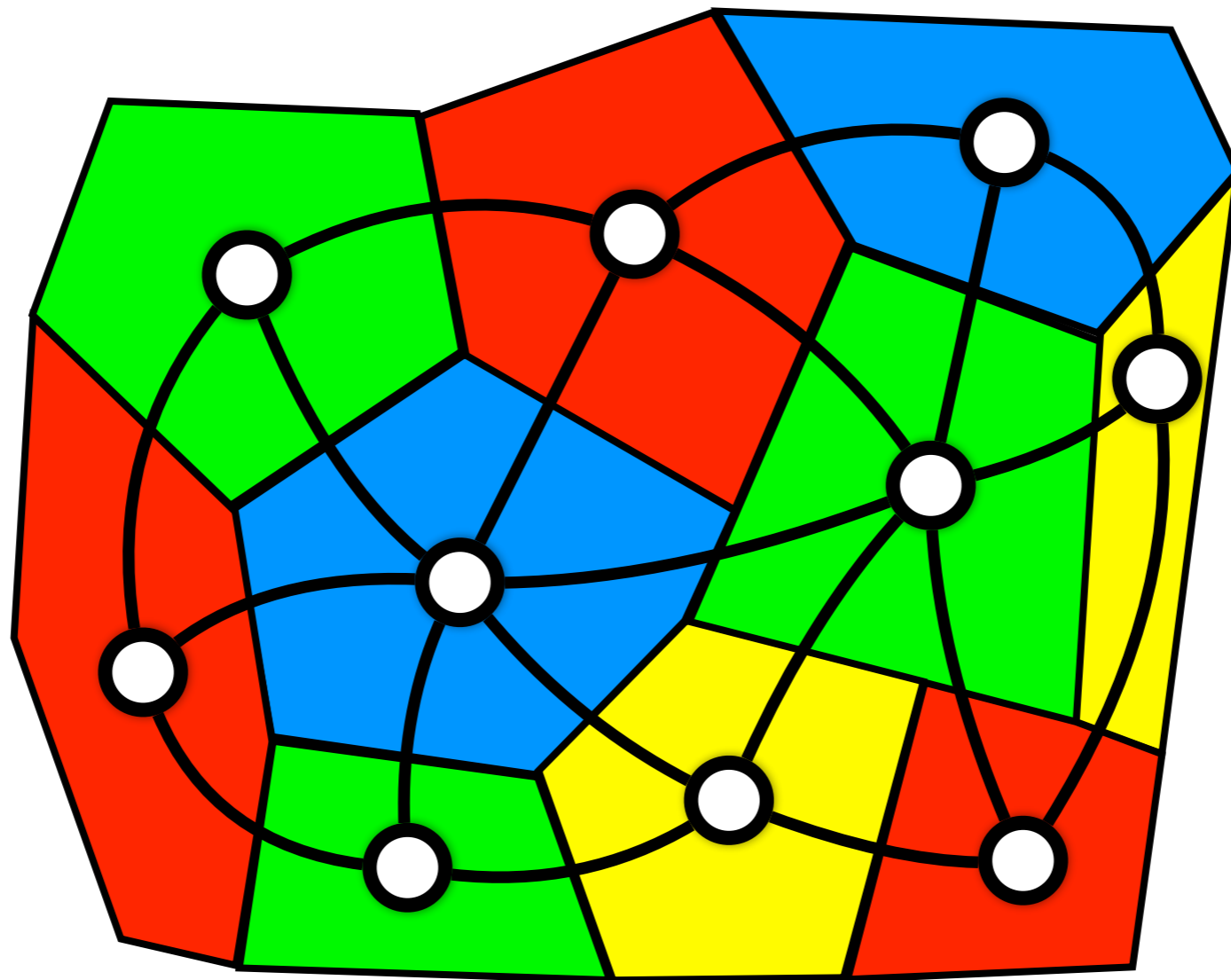
[Francis Guthrie 1852]



Coloring a Map

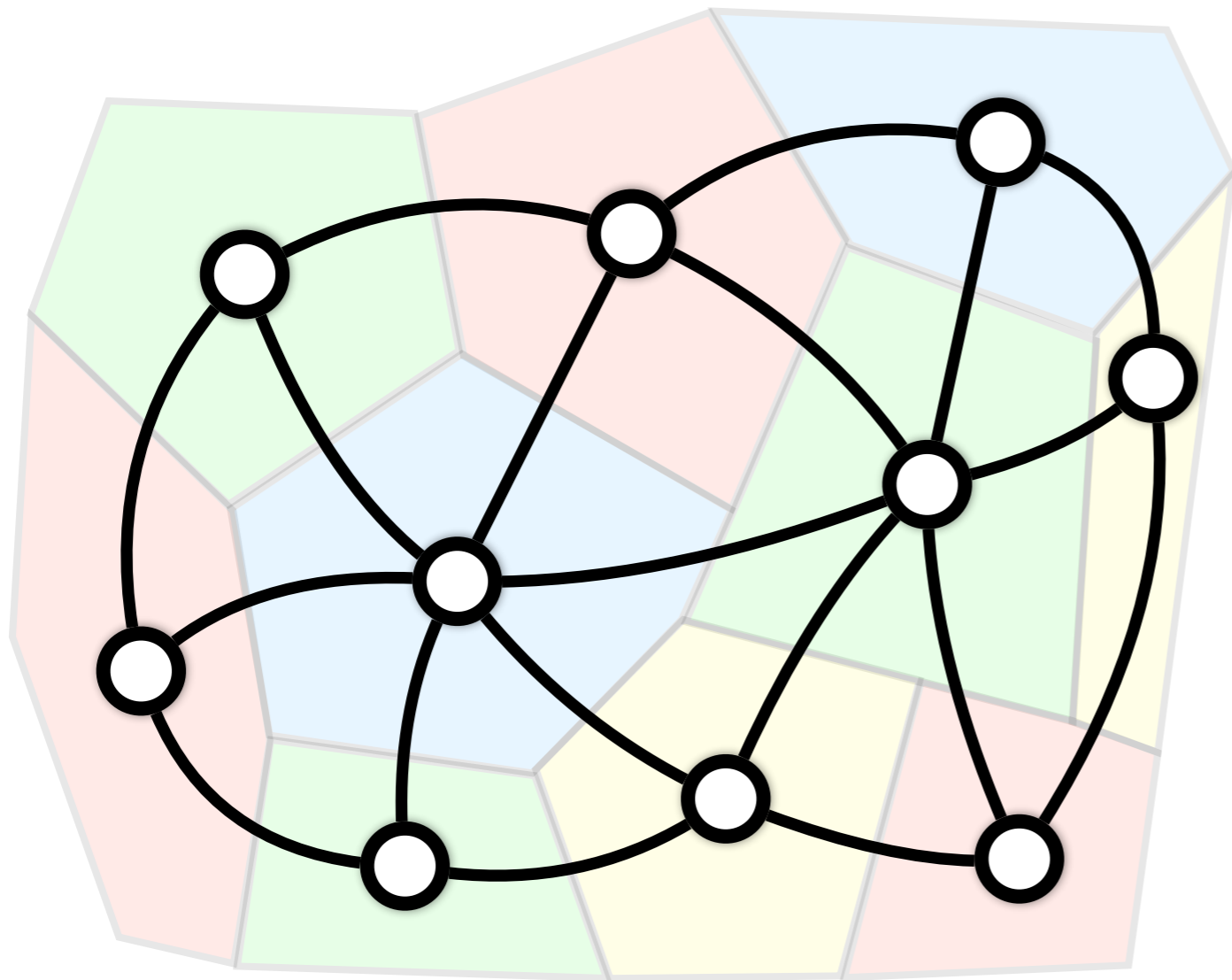


Graph Representation



Compact Description

$$G=(V,E)$$

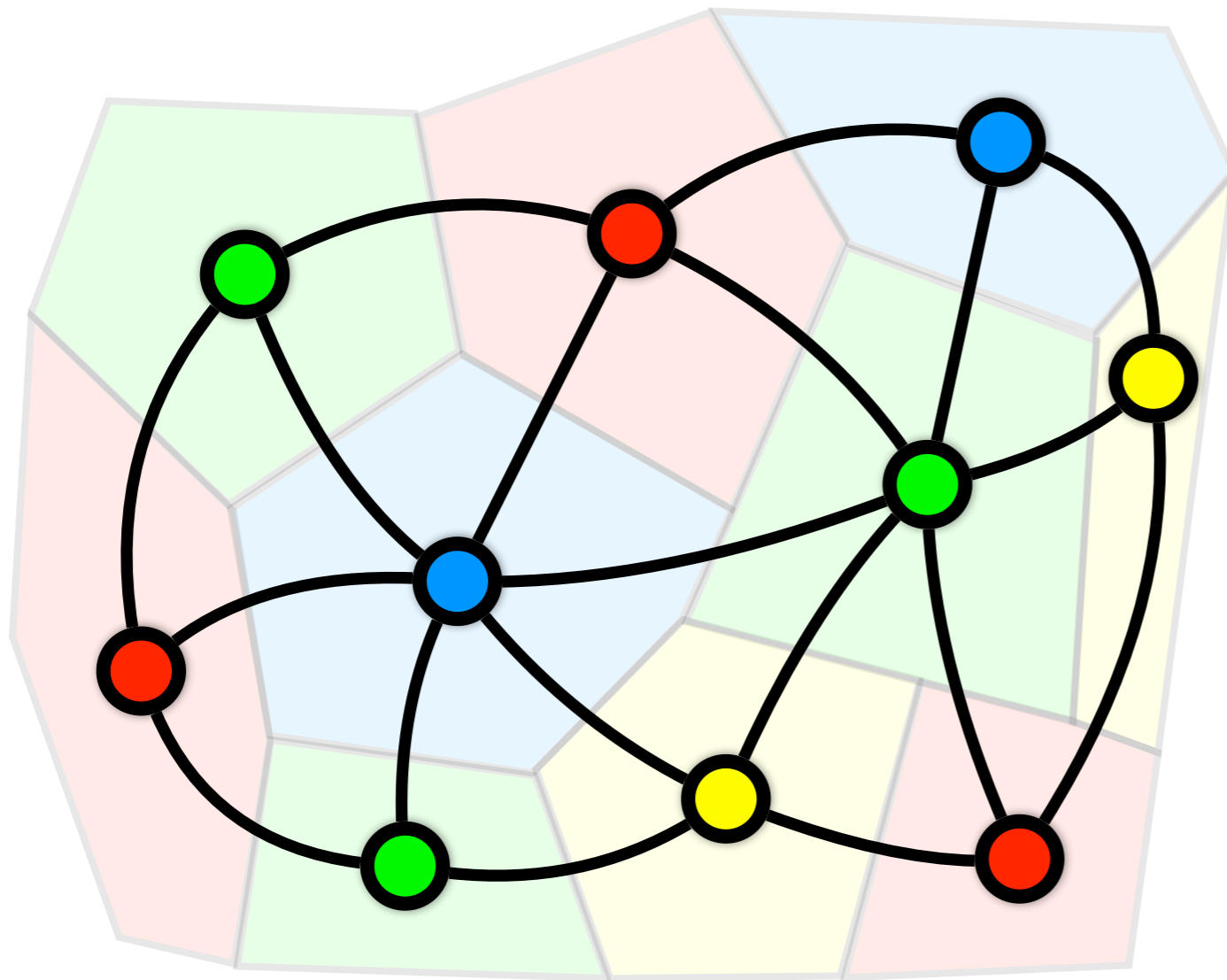


V =Set of vertices

E =Set of edges,
pairs of vertices

Vertex Coloring

$$G=(V,E)$$



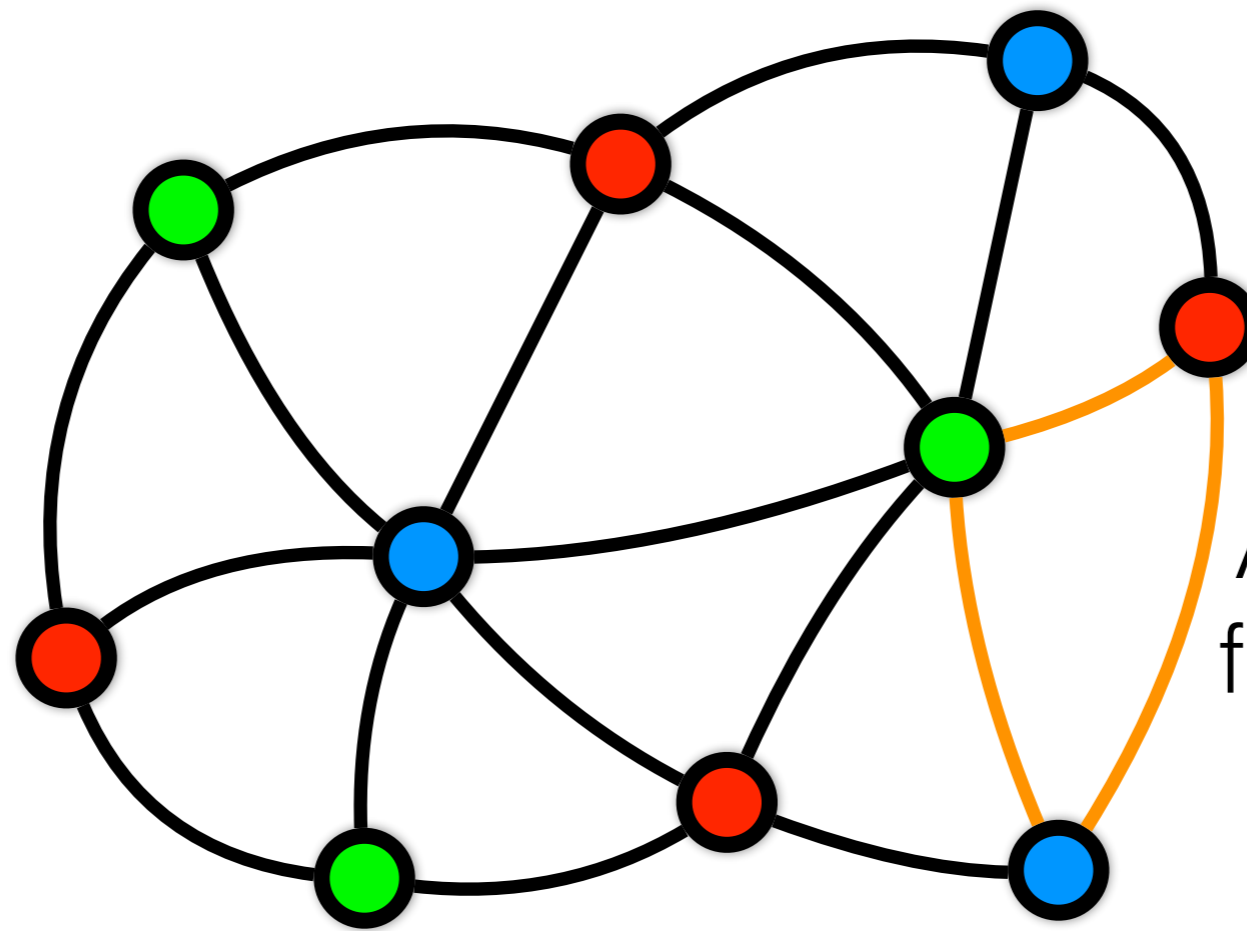
V =Set of vertices

E =Set of edges,
pairs of vertices

Chromatic Number

- Given a graph $G=(V,E)$, the chromatic number $\chi(G)$ is the smallest number of colors needed to color the vertices V so that for every edge e in E , the endpoints have different colors.

The Chromatic Number is Three



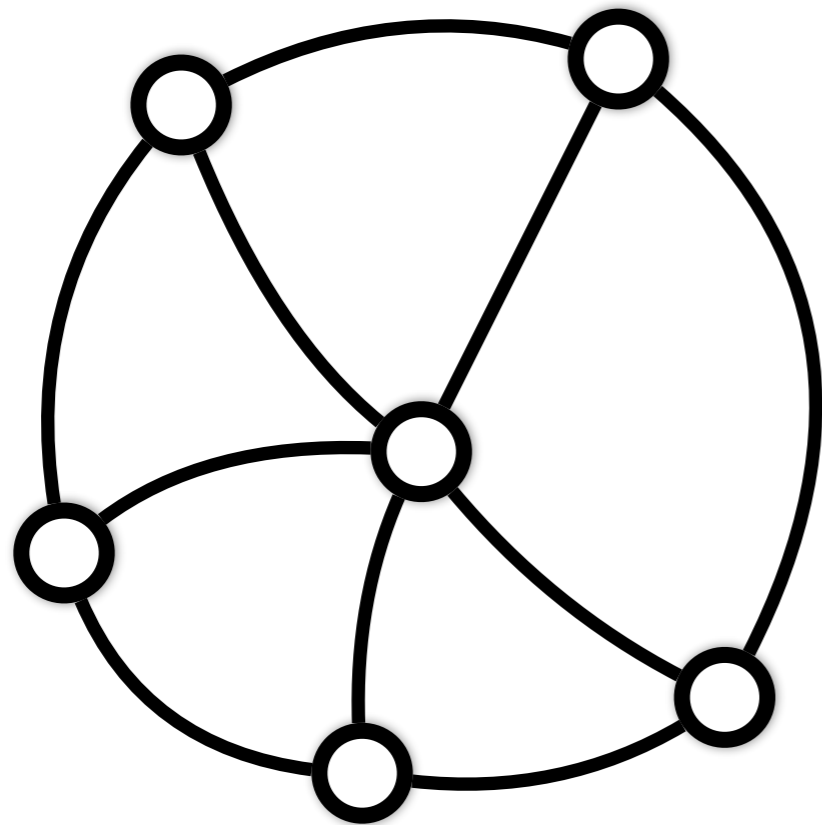
An odd cycle
forces at least
three colors

Planar Graphs

- A graph is planar if it can be drawn in the plane without any edges crossing each other.
- [Appel&Haken 1976] The chromatic number of a planar graph is at most four.
- Proof through computer assisted case analysis, it gives a (complicated) polynomial time algorithm.

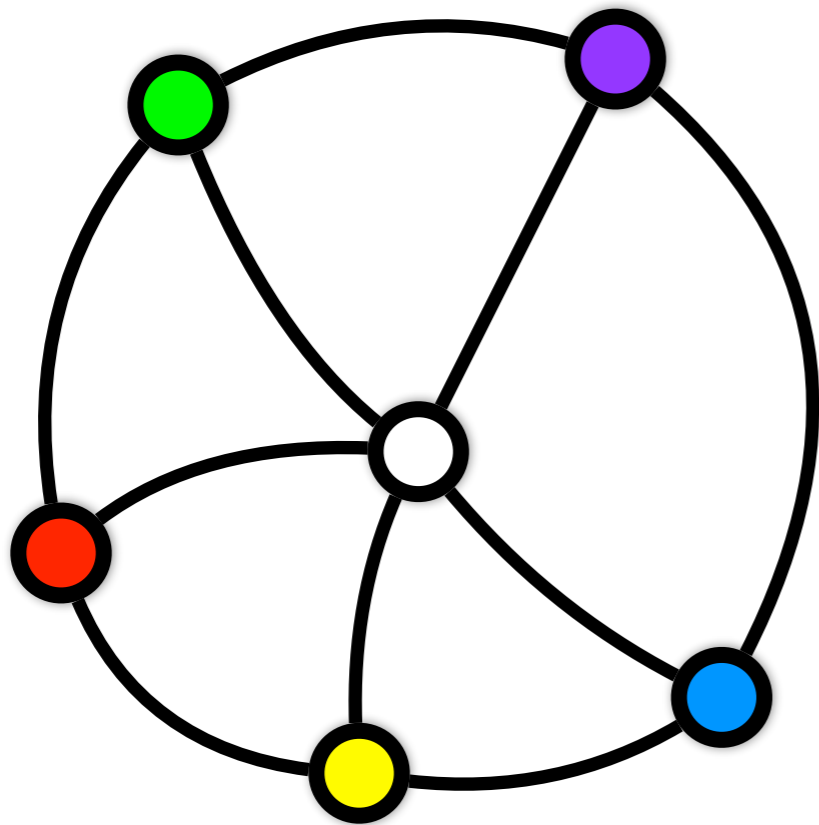
Simple Argument $\chi(G) < 6$

[Heawood 1890]



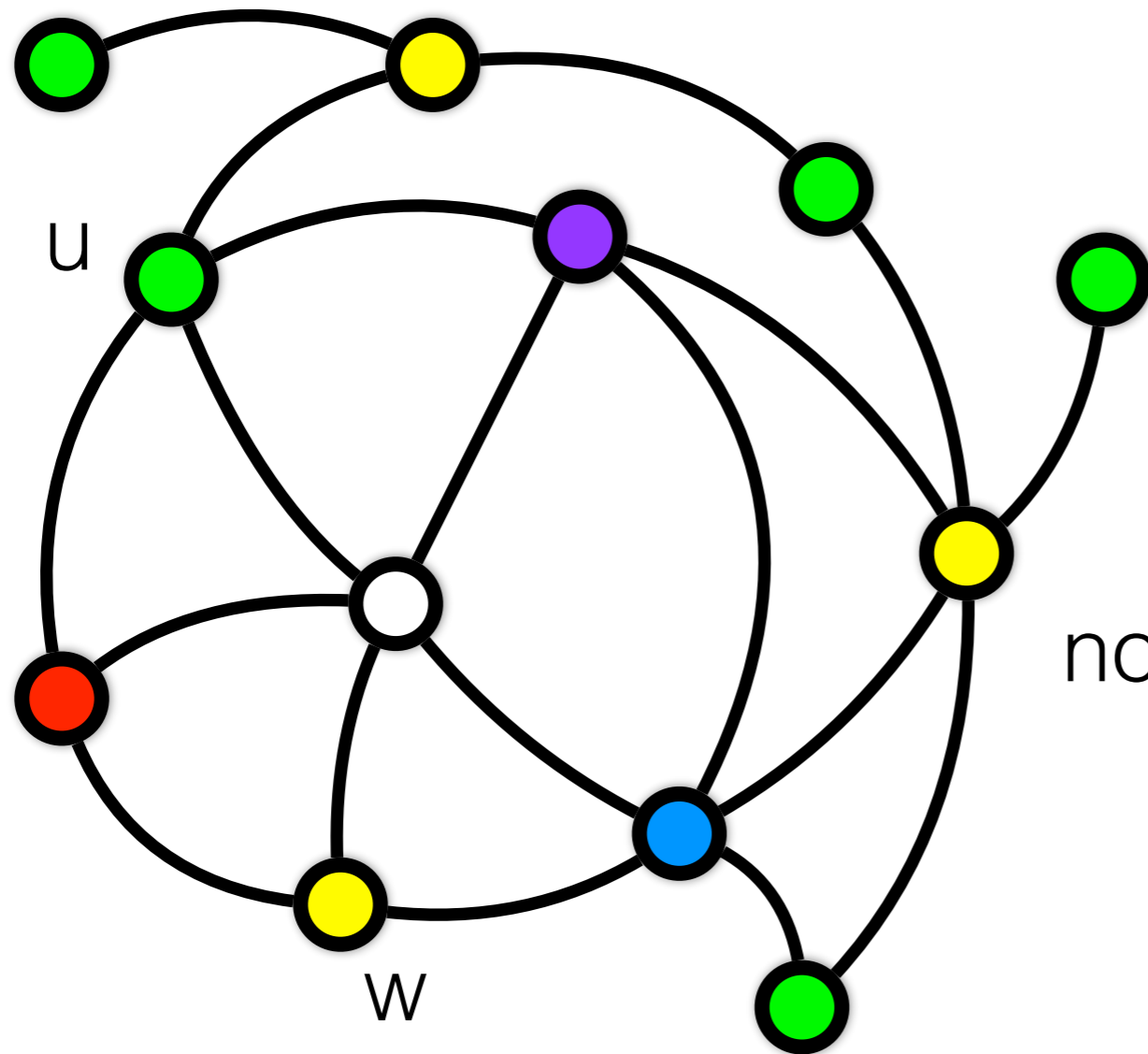
There must be a vertex
of degree at most 5.

Simple Argument $\chi(G) < 6$



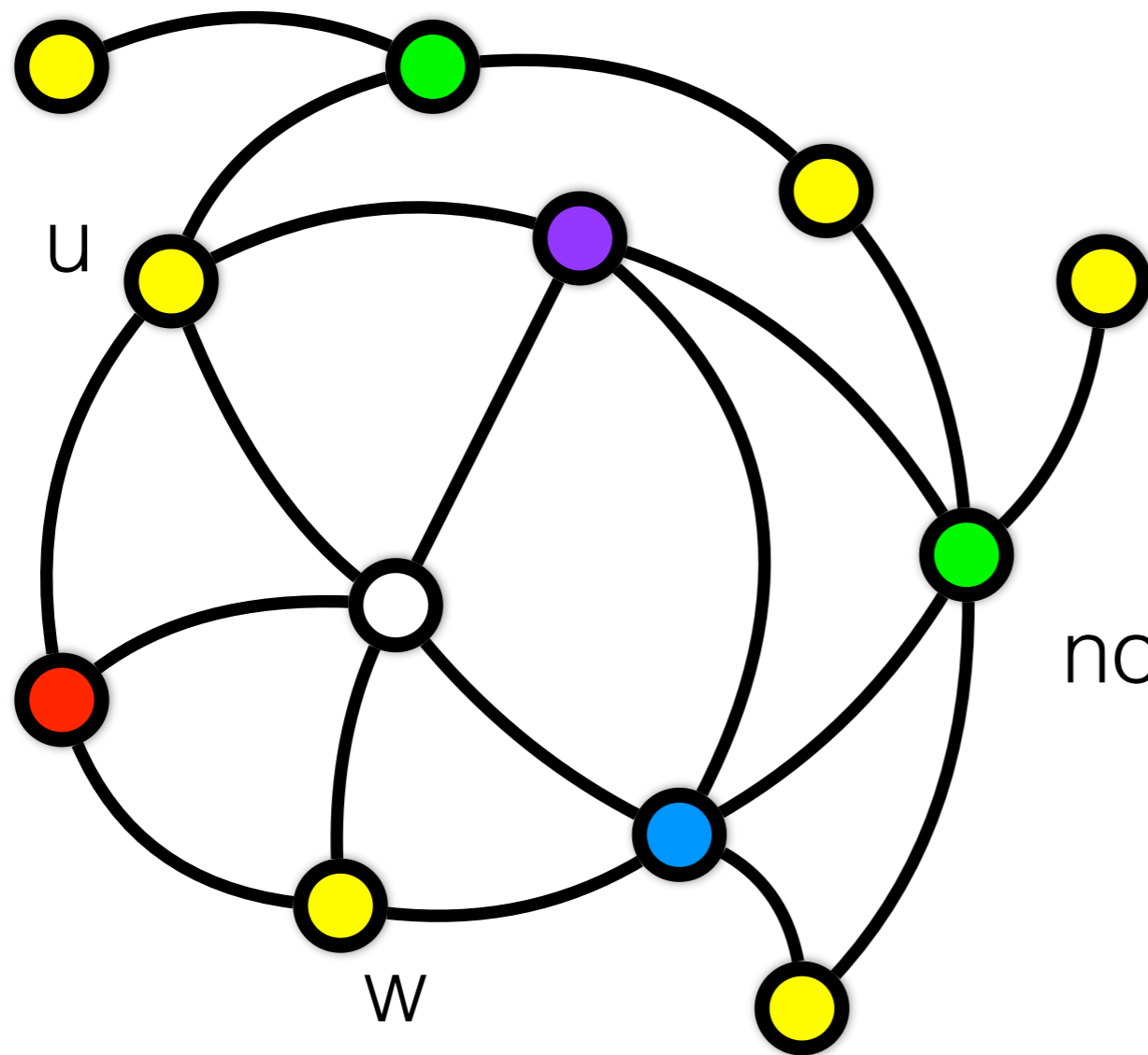
Assume its neighbors are colored in 5 colors.

Simple Argument $\chi(G) < 6$



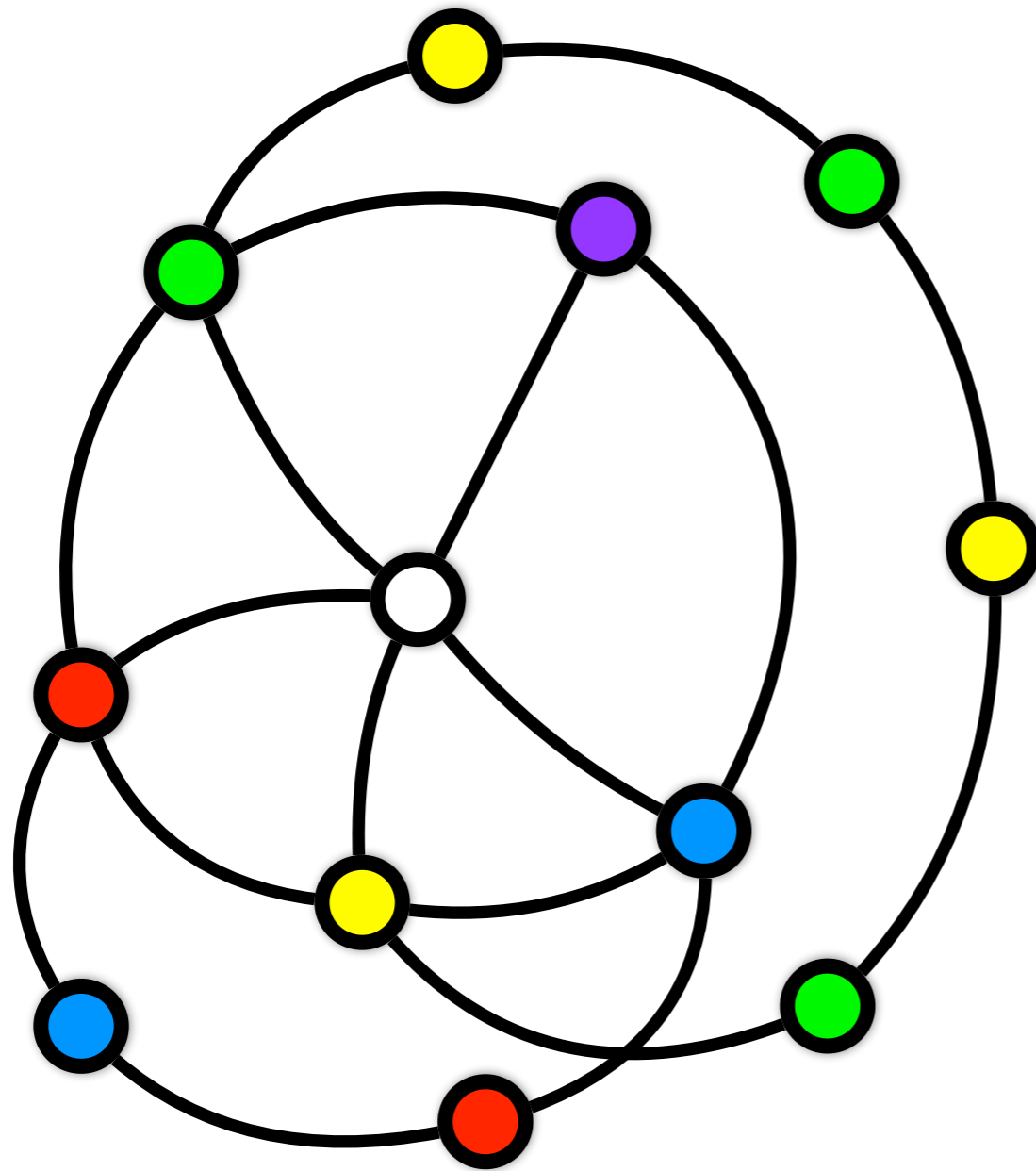
If one neighbor u is in alternating colored component that is not connected to neighbor w of the other color, then Recoloring is possible!

Simple Argument $\chi(G) < 6$



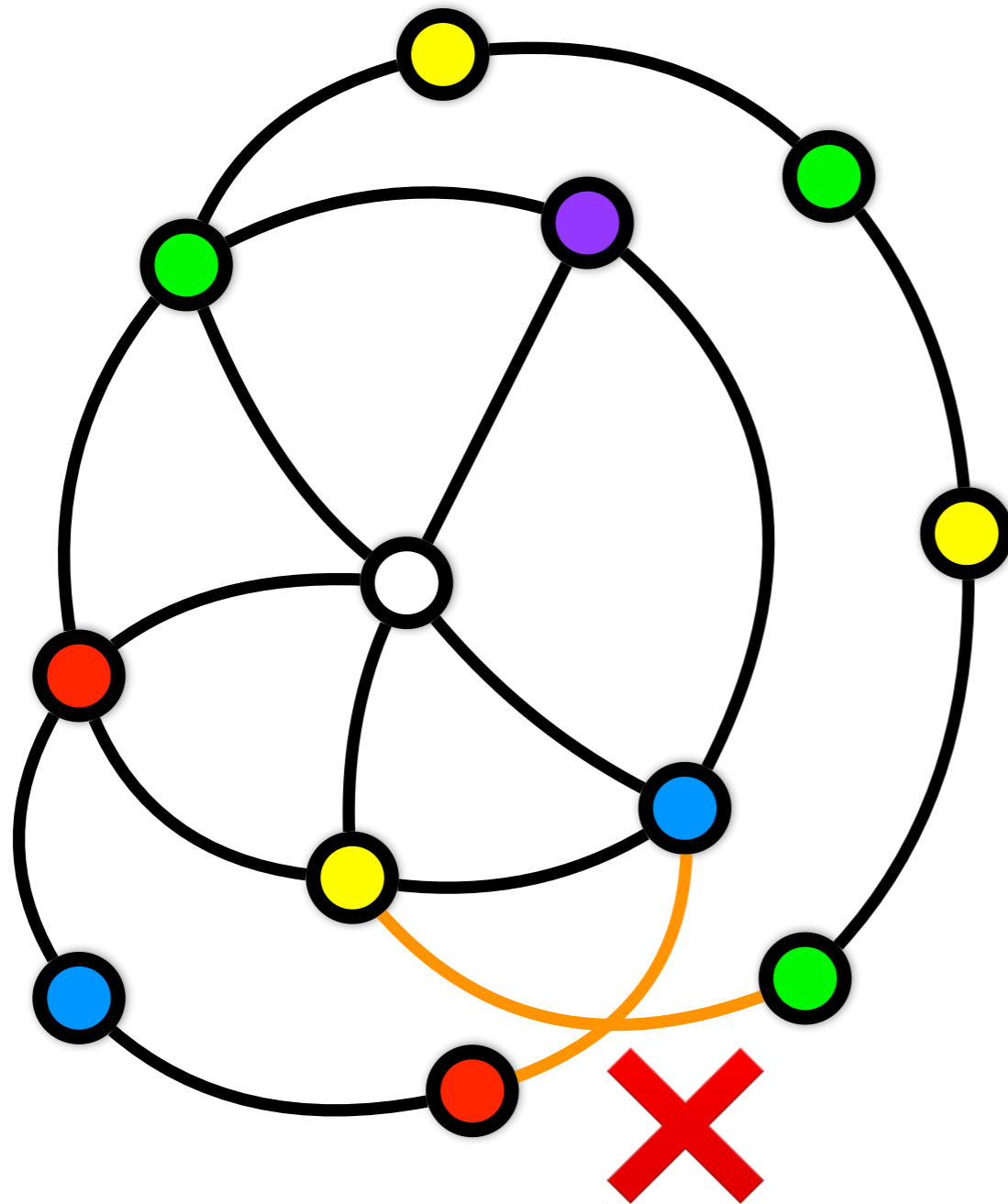
If one neighbor u is in alternating colored component that is not connected to neighbor w of the other color, then Recoloring is possible!

Simple Argument $\chi(G) < 6$



Must be alternately colored paths between neighbors.

Simple Argument $\chi(G) < 6$

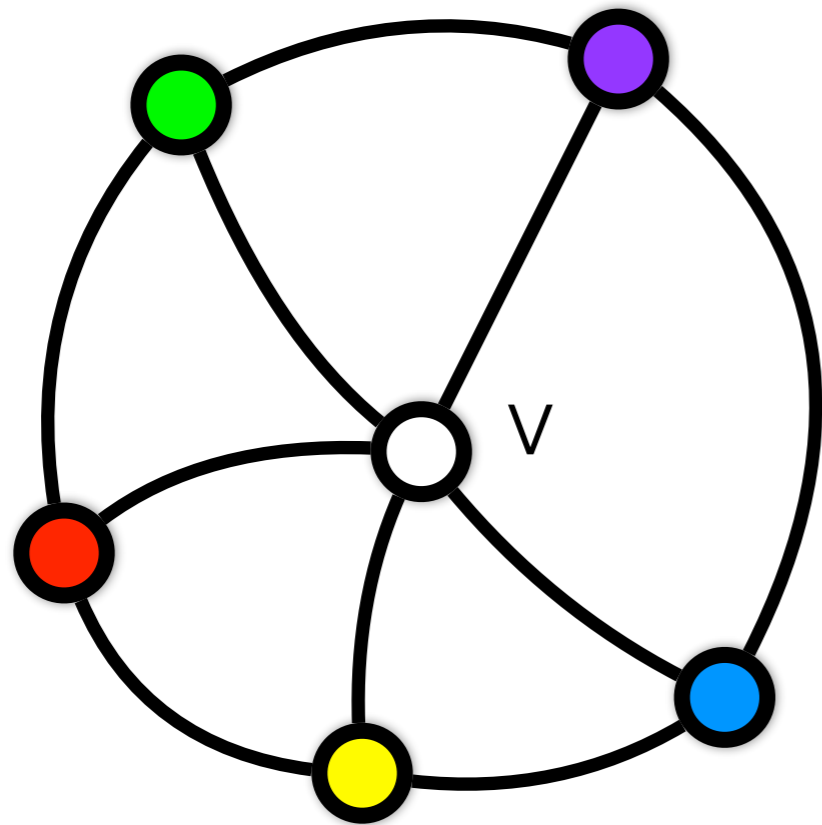


Paths must cross each other
 \Rightarrow contradiction!

Polynomial Time Algorithm for Five Coloring Planar Graphs

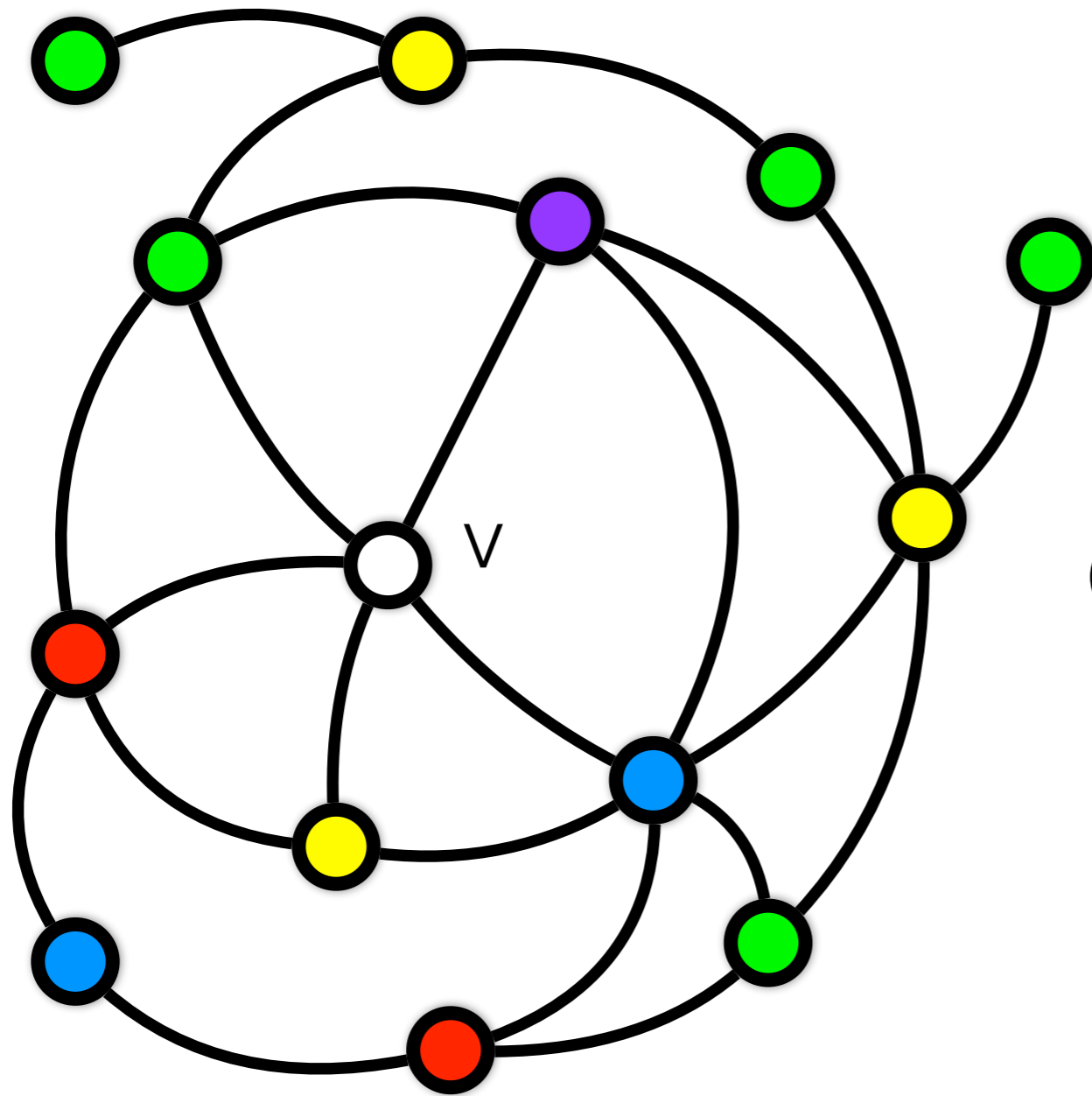
- [**Preprocessing**] Remove the vertex of smallest degree and put it on a stack. Repeat until all vertices are on the stack.
- [**Coloring**] As long as not all vertices have been colored, pop the stack and color the current vertex with the first available valid color given the already coloured vertices.
 - If no valid color is available, it must be possible to recolor a colored neighbor by Heawood's argument.
 - Find a vertex that can be recolored, do that and continue.

Recoloring a Neighbor



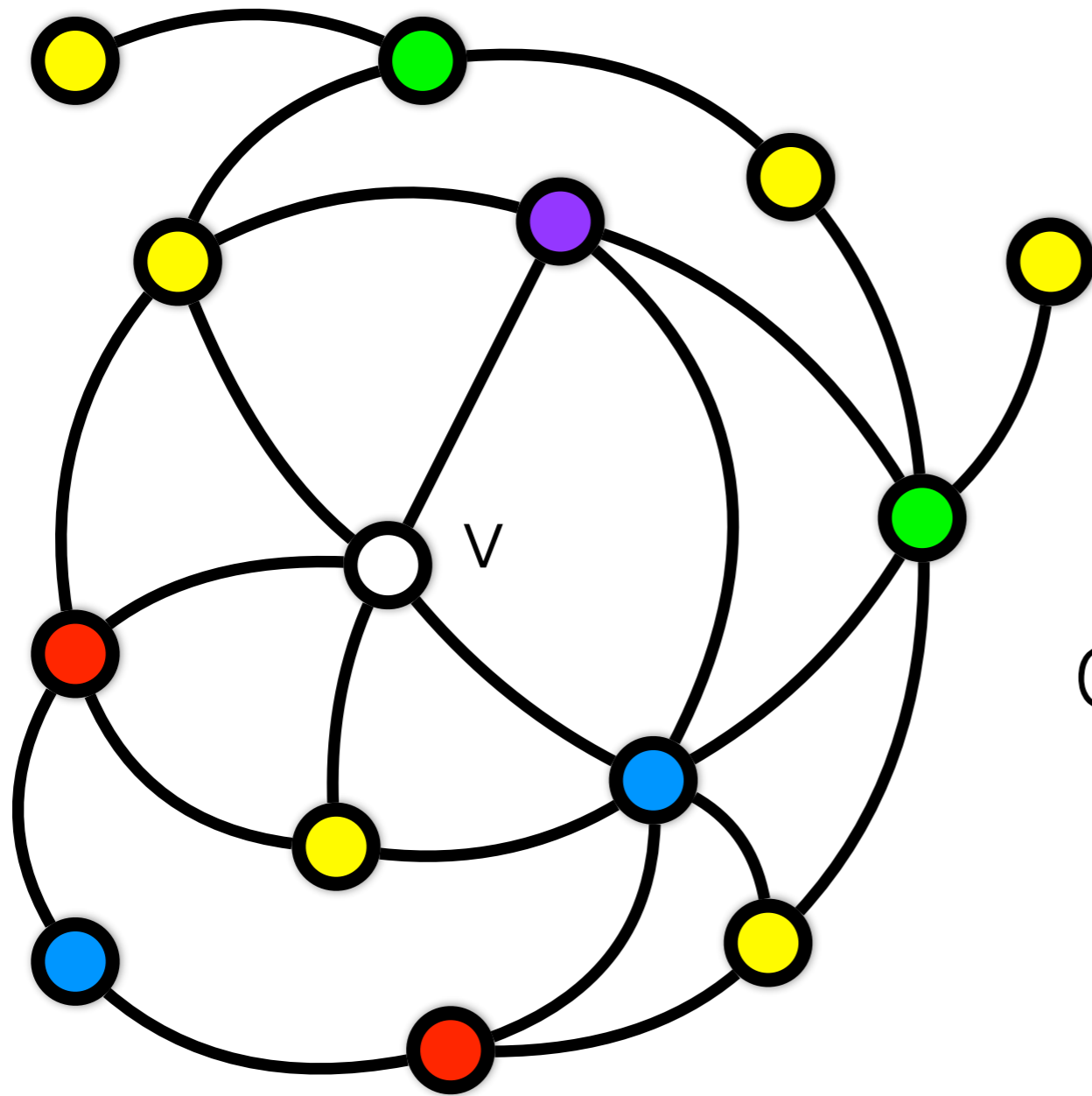
None of the five colors
are valid for v .

Recoloring a Neighbor



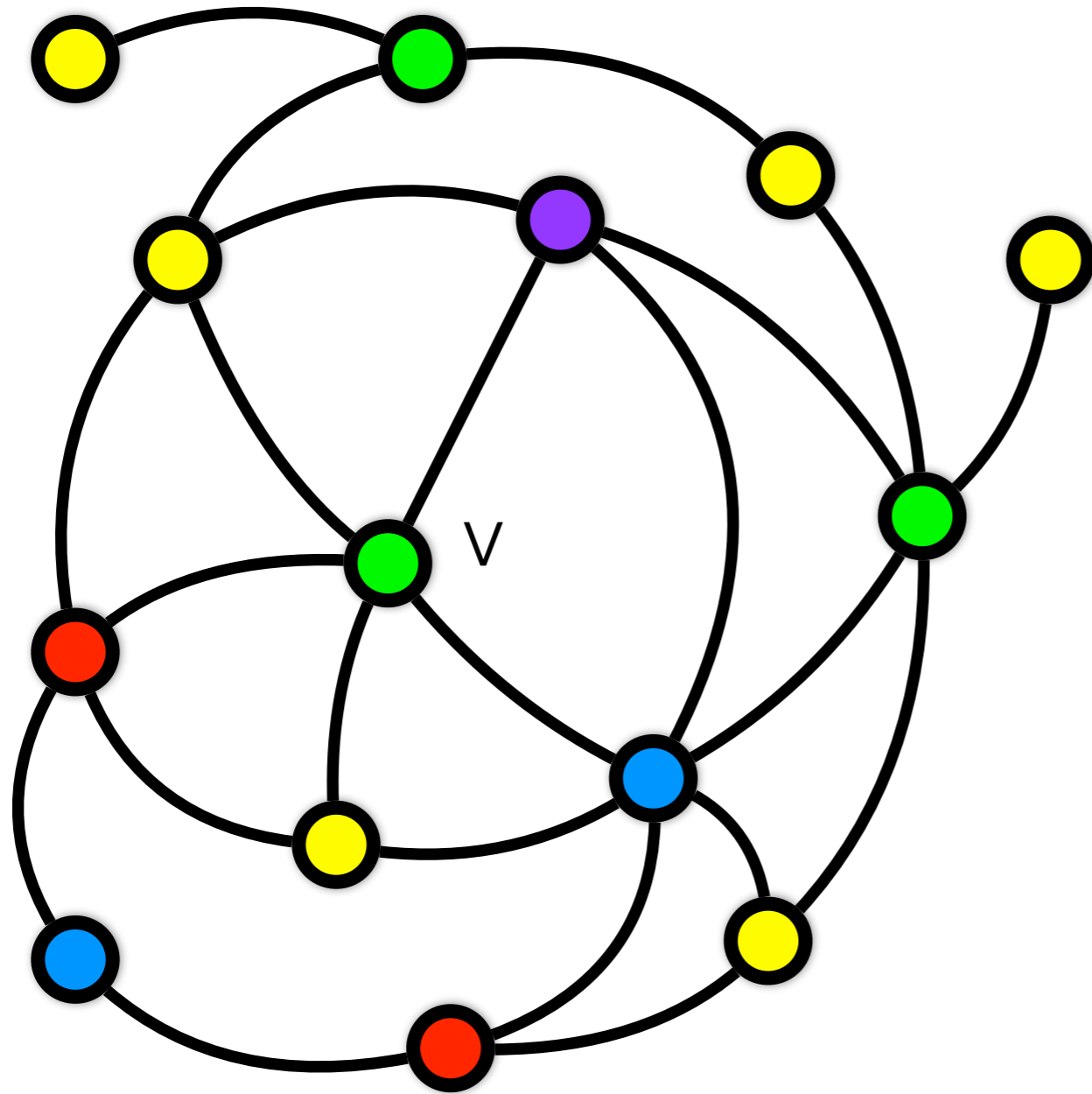
Green-Yellow component
only incident with one
neighbor of v .

Recoloring a Neighbor



Swap colors on
Green-Yellow component.

Recoloring a Neighbor



Green color is now
valid for v.

Hardness of Coloring

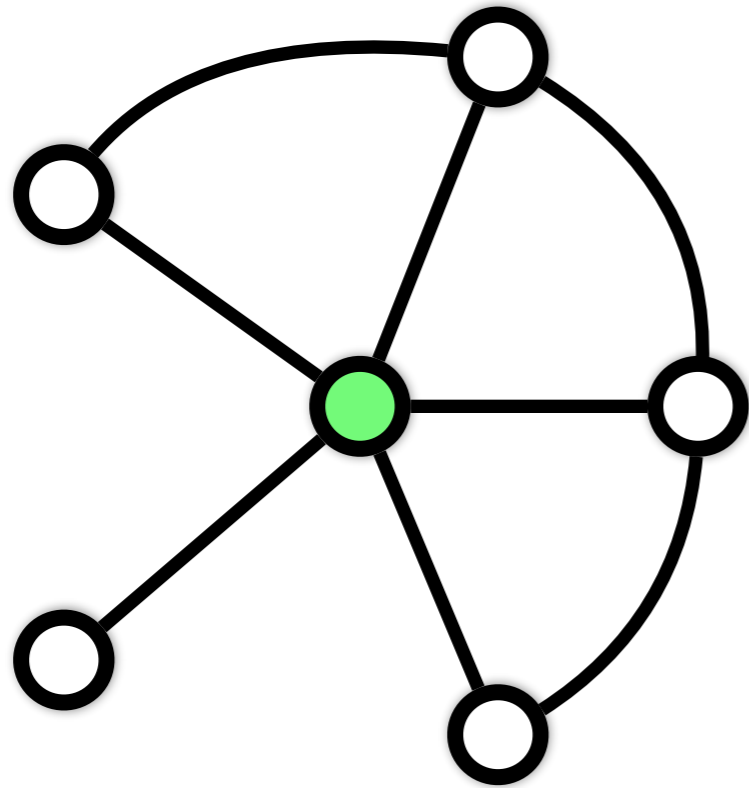
- Computing the chromatic number in a general graph is NP-complete. In fact deciding if $\chi(G)=3$ in a planar graph is NP-complete.
- The best polynomial time algorithms for coloring a 3-colorable graph uses more than $n^{1/6}$ colors!

Coloring a 3-colorable Graph with $O(n^{1/2})$ Colors in Polynomial Time

[Wigderson 1982]

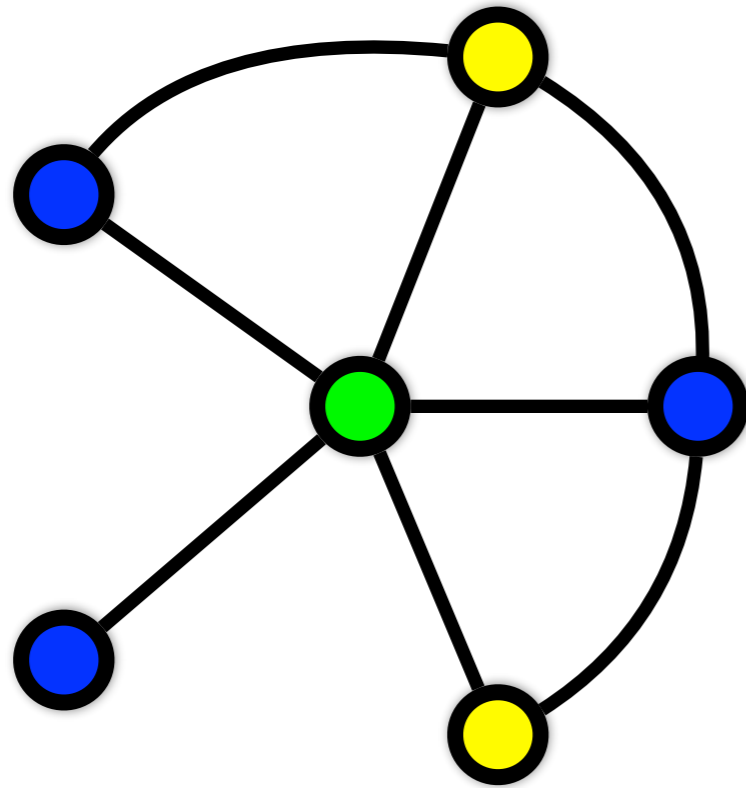
Note: Graphs need not be planar,
(they are just easier to visualise).

Coloring a 3-colorable Graph with $O(n^{1/2})$ Colors



If there exists vertex of degree more than $n^{1/2}$, color it and its neighbours with three fresh colors.

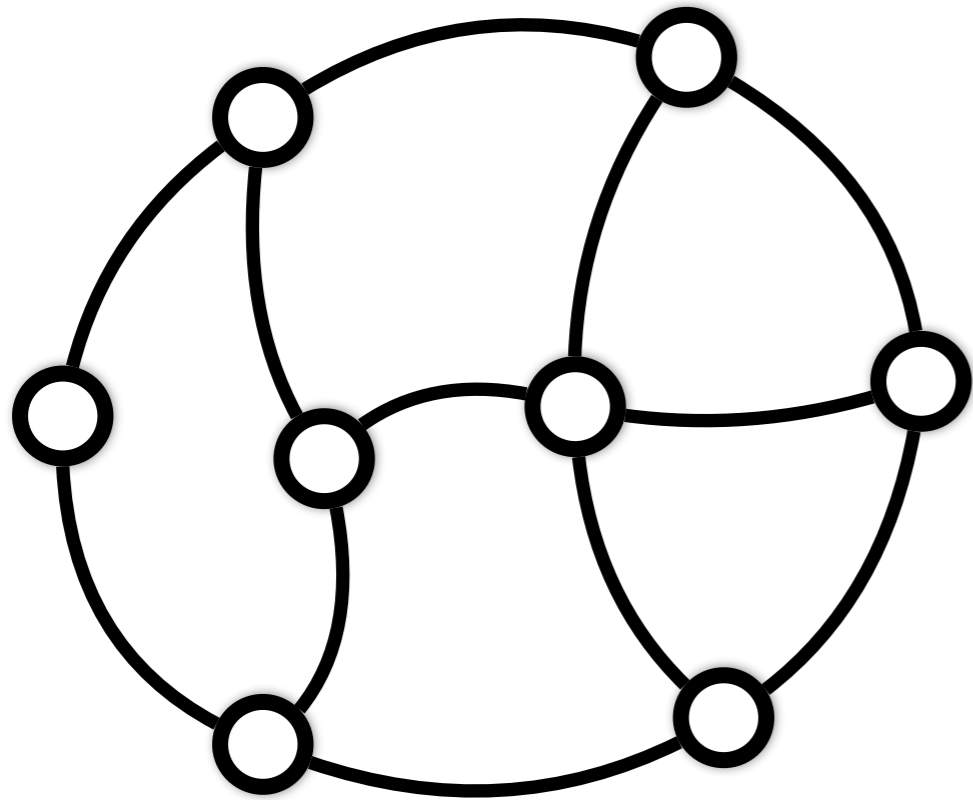
Coloring a 3-colorable Graph with $O(n^{1/2})$ Colors



If there exists vertex of degree more than $n^{1/2}$, color it and its neighbours with three fresh colors.

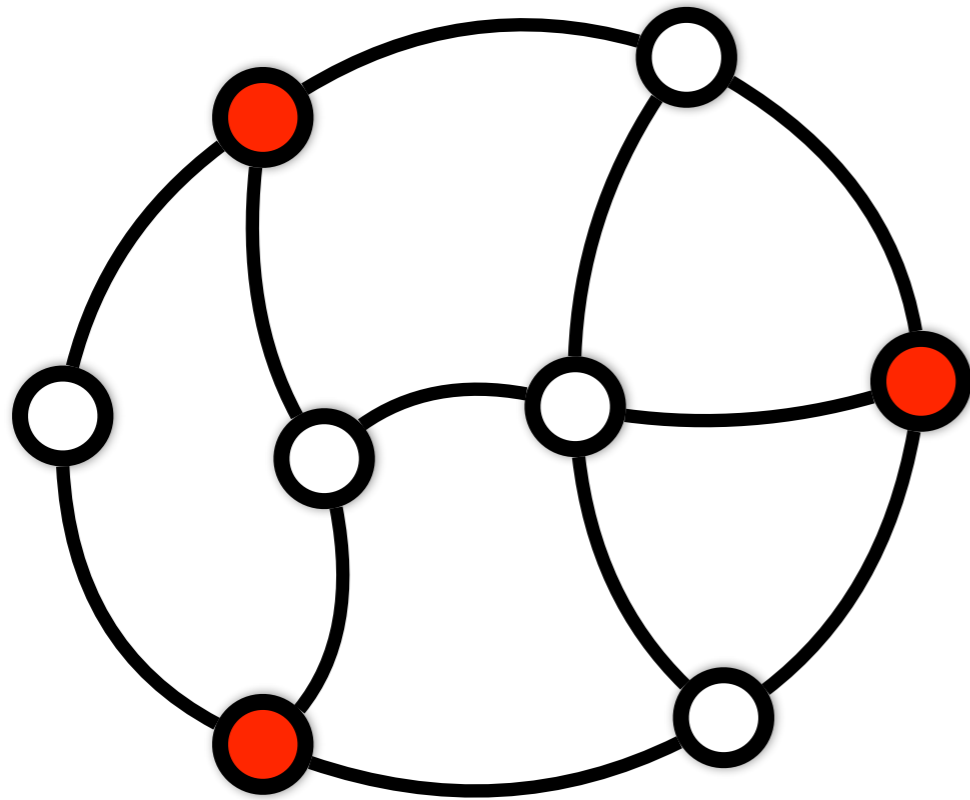
The neighbourhood must be bipartite, and hence it is easy to two-color.

Coloring a 3-colorable Graph with $O(n^{1/2})$ Colors



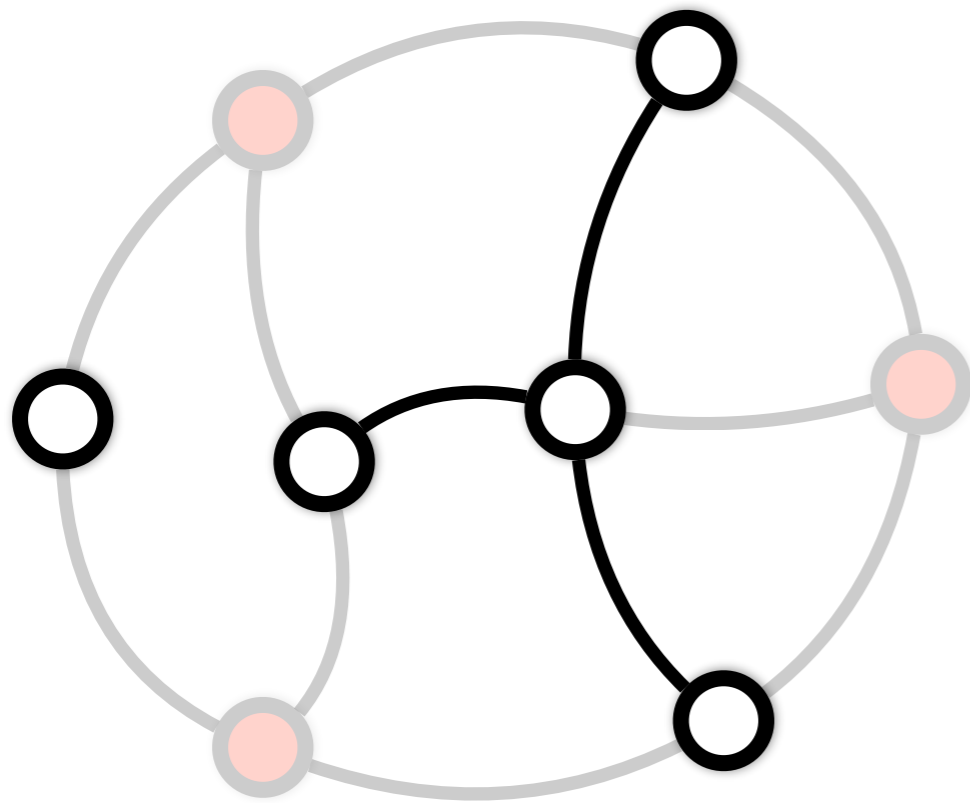
If no vertex exist with degree more than $n^{1/2}$, we can find a maximal independent set of size $n^{1/2}$. Color it with one fresh color.

Coloring a 3-colorable Graph with $O(n^{1/2})$ Colors



If no vertex exist with degree more than $n^{1/2}$, we can find a maximal independent set of size $n^{1/2}$. Color it with one fresh color.

Coloring a 3-colorable Graph with $O(n^{1/2})$ Colors



If no vertex exist with degree more than $n^{1/2}$, we can find a maximal independent set of size $n^{1/2}$. Color it with one fresh color.

Repeat strategy on remaining graph.

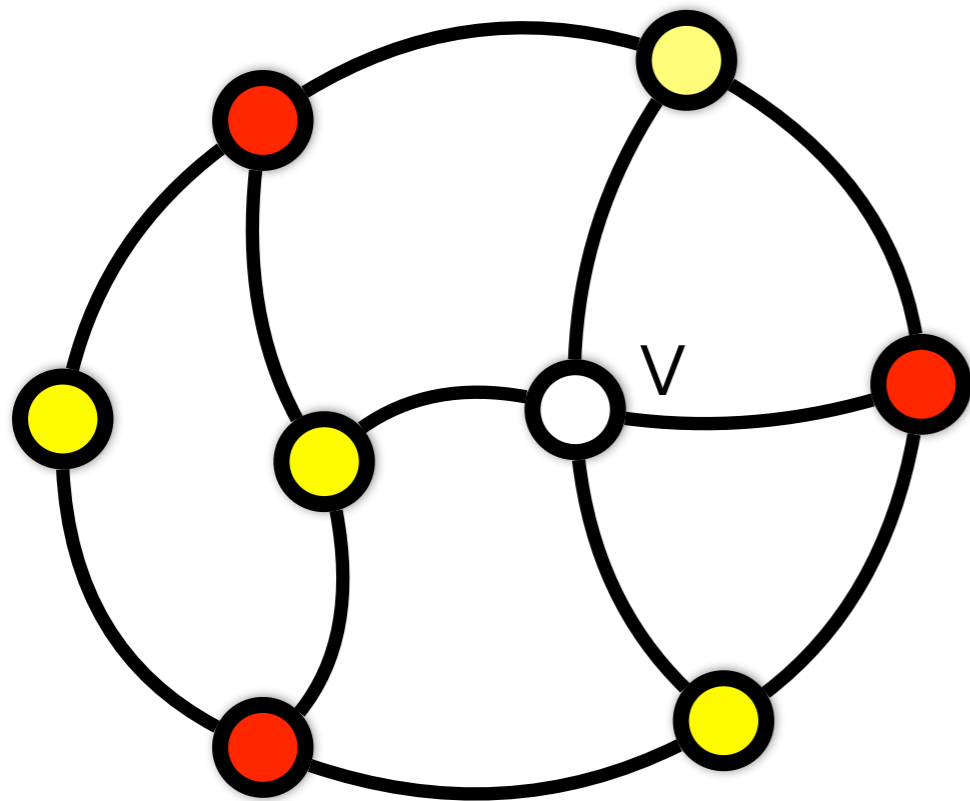
Color Usage Analysis

- In the first rule, we color $n^{1/2}$ vertices with 3 new colors.
- In the second rule, we color $n^{1/2}$ vertices with 1 new color.
- There can be at most $n^{1/2}$ steps before all vertices have been colored, in total $O(n^{1/2})$ colors used.

Exponential Time Algorithms

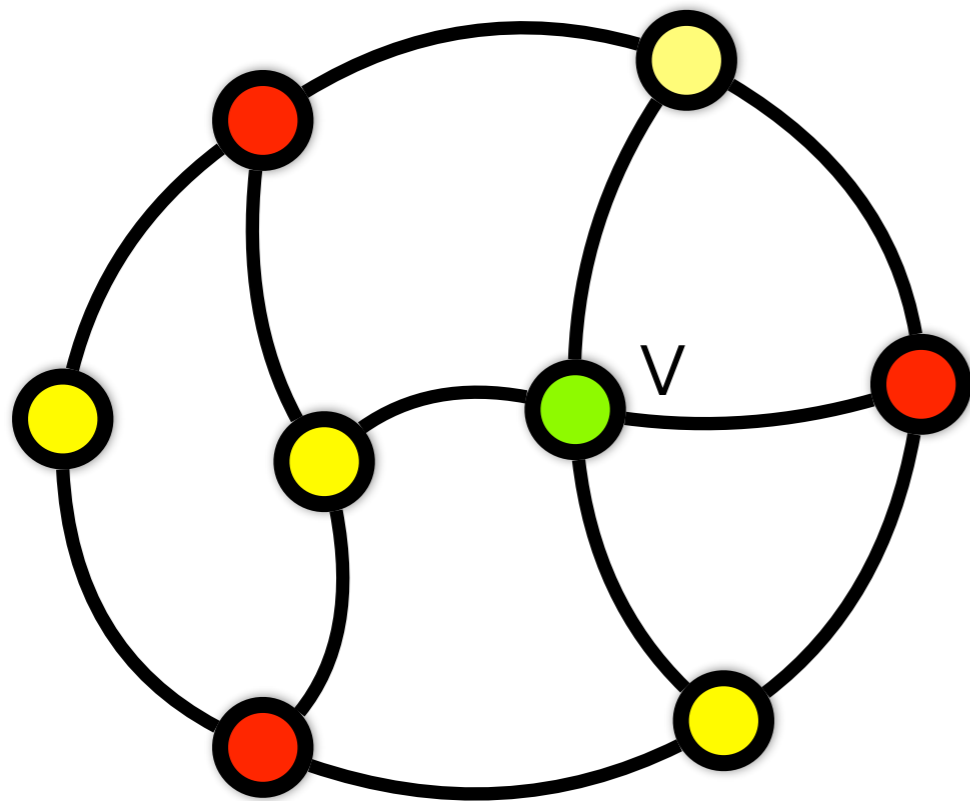
- What is the fastest worst case running time algorithm we can guarantee for *exactly* computing the chromatic number?
- We want to make sure that if a k -coloring exists the algorithm will learn about it, not just an approximation.
- We want to guarantee that the running time is bounded by as small as possible an (exponential) function in n , the number of vertices.

Naive algorithm uses k^n
time.



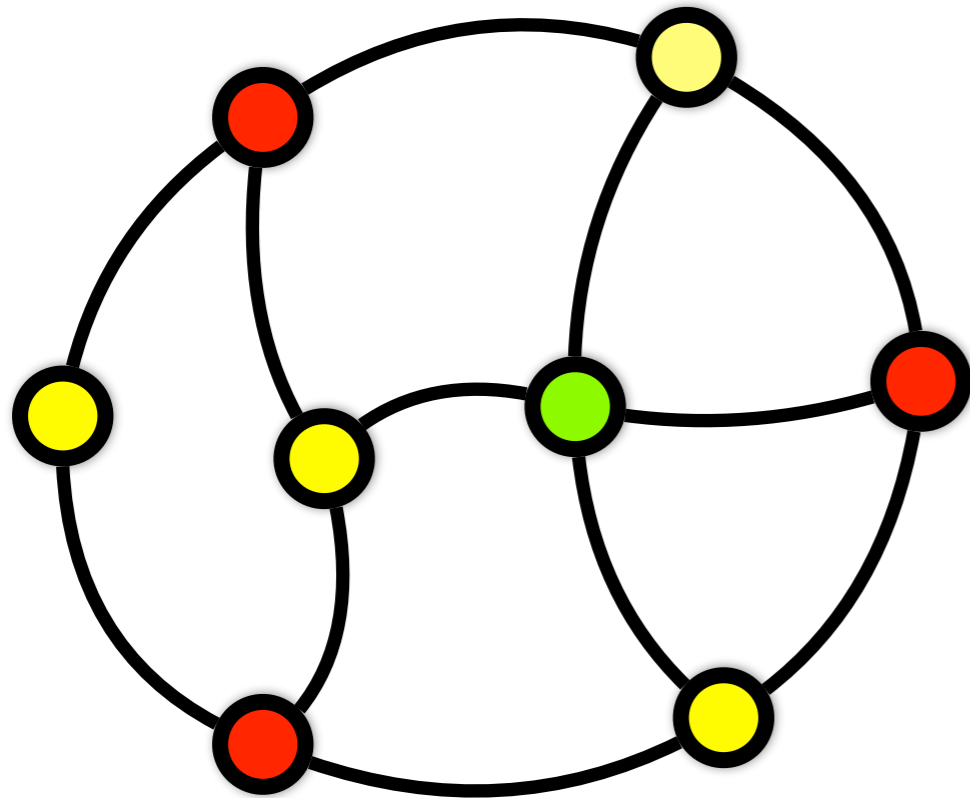
For every vertex, try each
of the k colors.

Naive algorithm uses k^n
time.



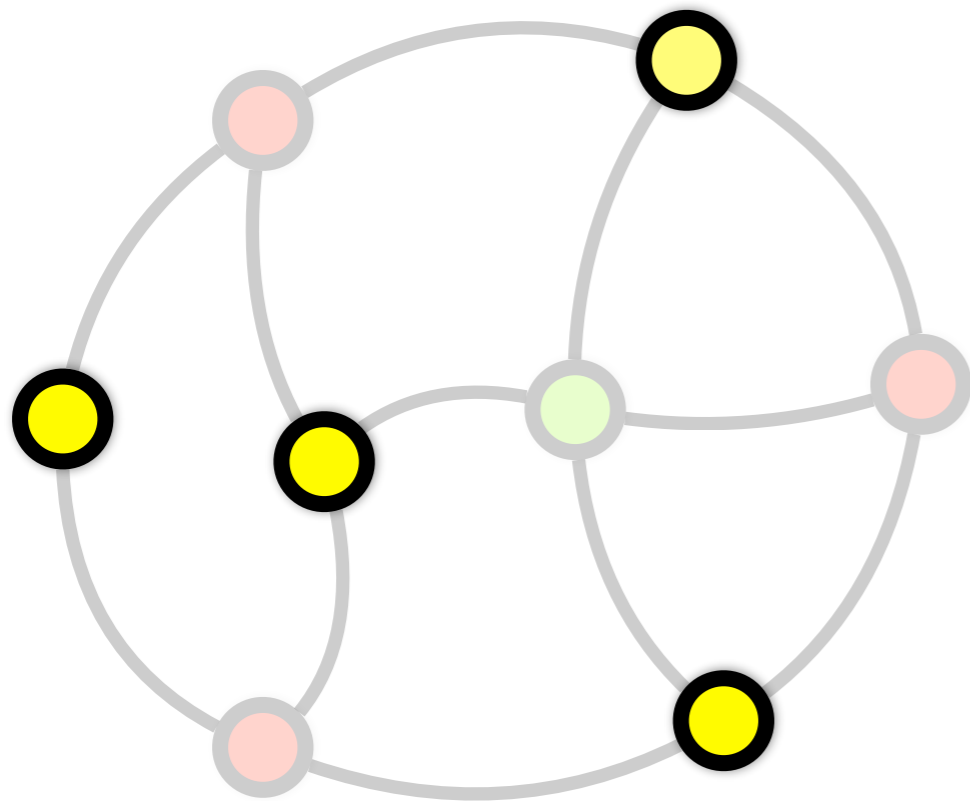
v cannot be yellow or red \Rightarrow
we might need to check k
colors for many vertices, and
backtrack many times.

Colorclasses



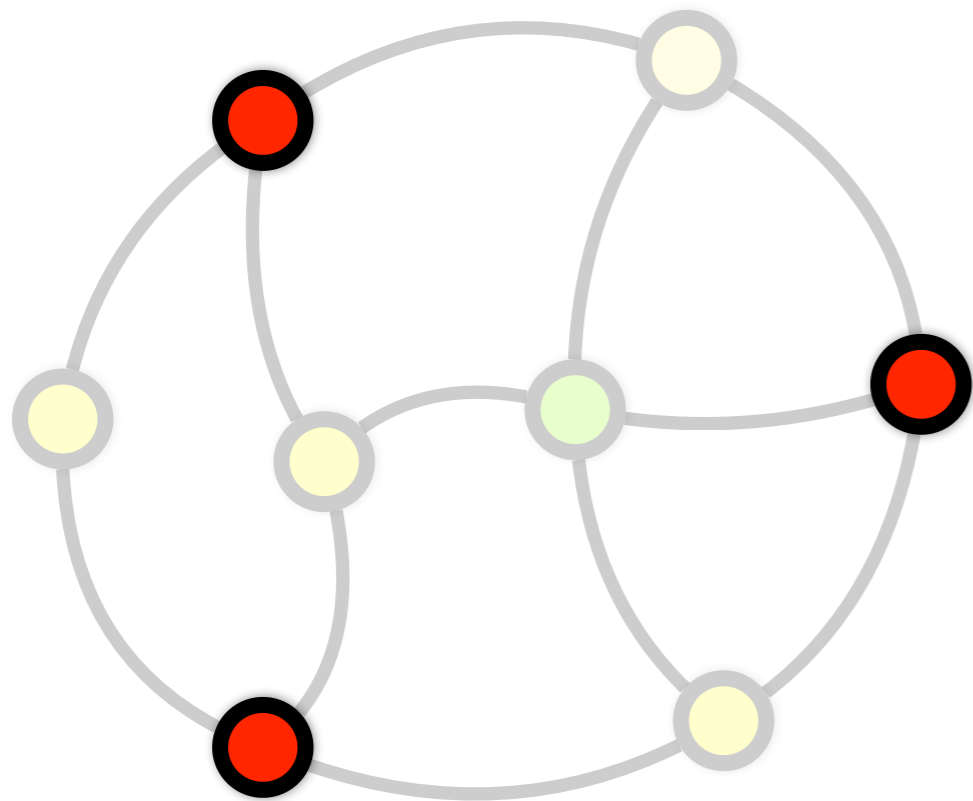
A colorclass is the set of vertices colored by the same color.
A k -coloring consists of k disjoint colorclasses.

Colorclasses



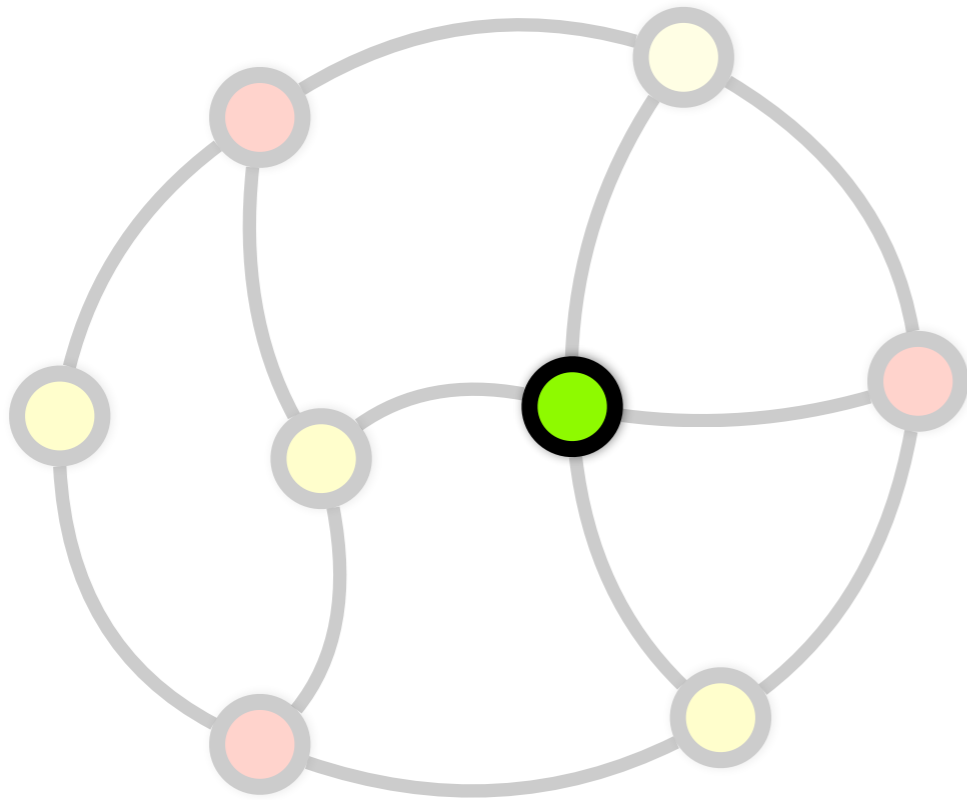
The Yellow colorclass

Colorclasses



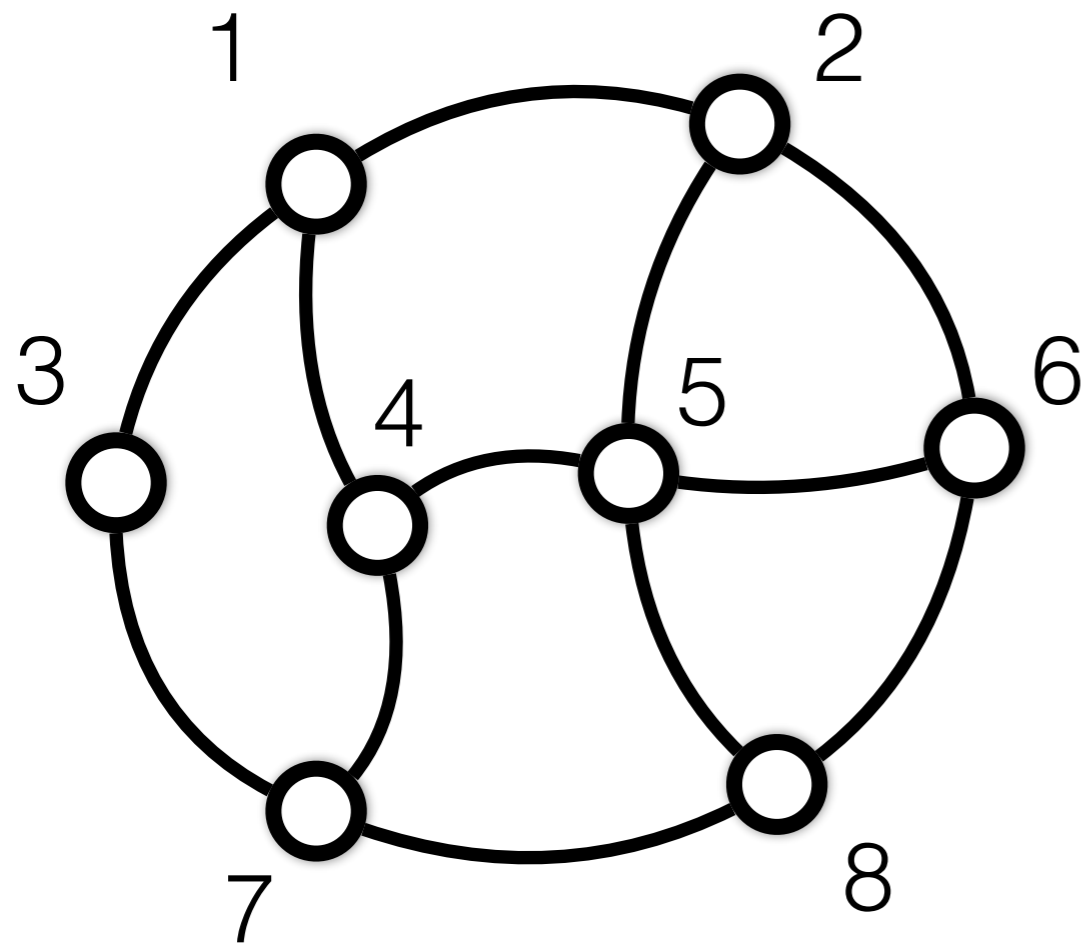
The Red colorclass.

Colorclasses



The Green colorclass.

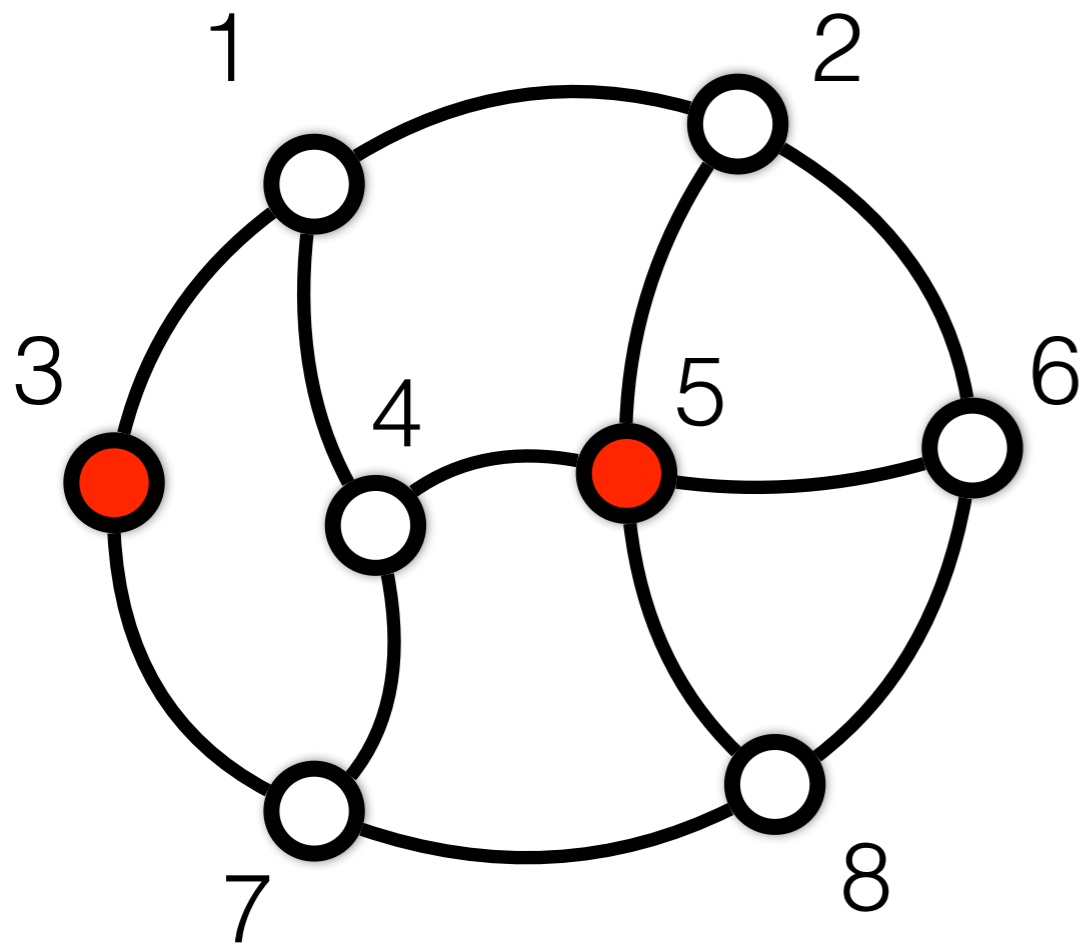
All Candidate Colorclasses



1 2 3 4 5 6 7 8 | 1 2 3 4 5 6 7 8

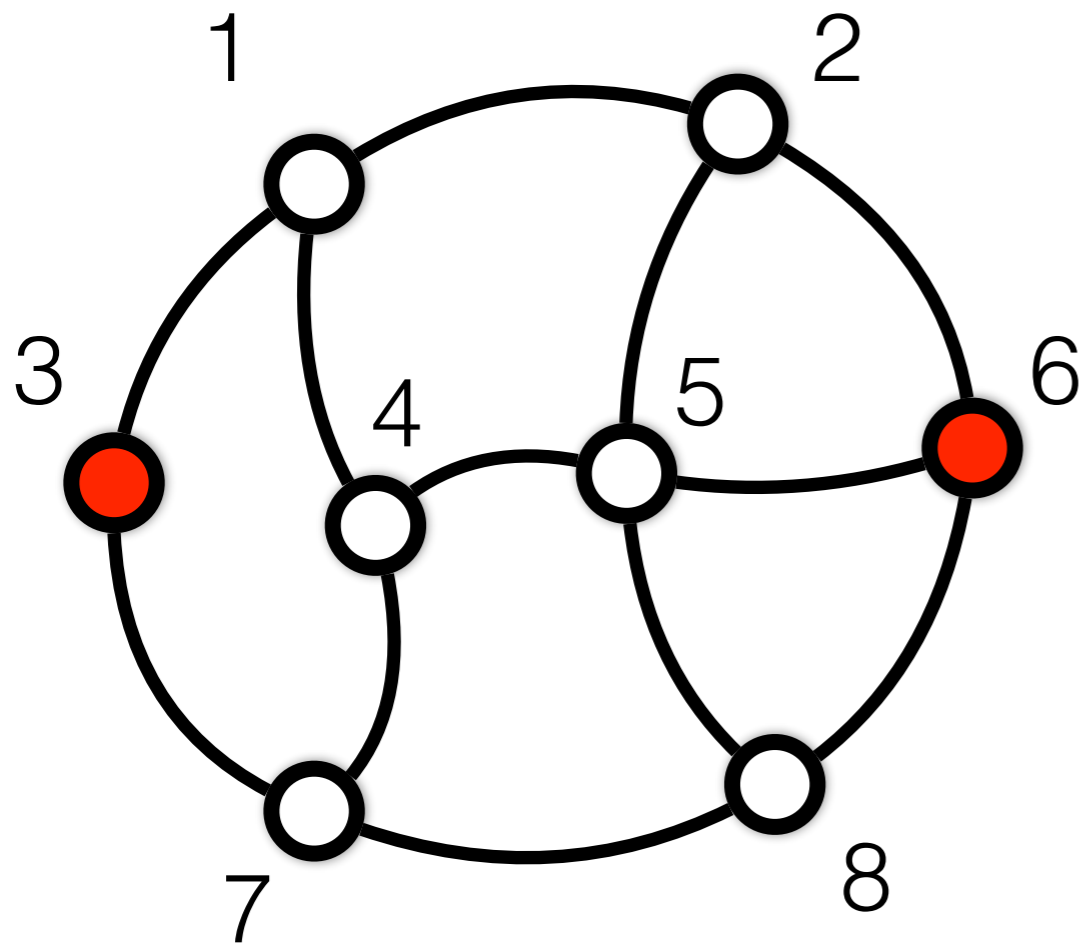


All Candidate Colorclasses



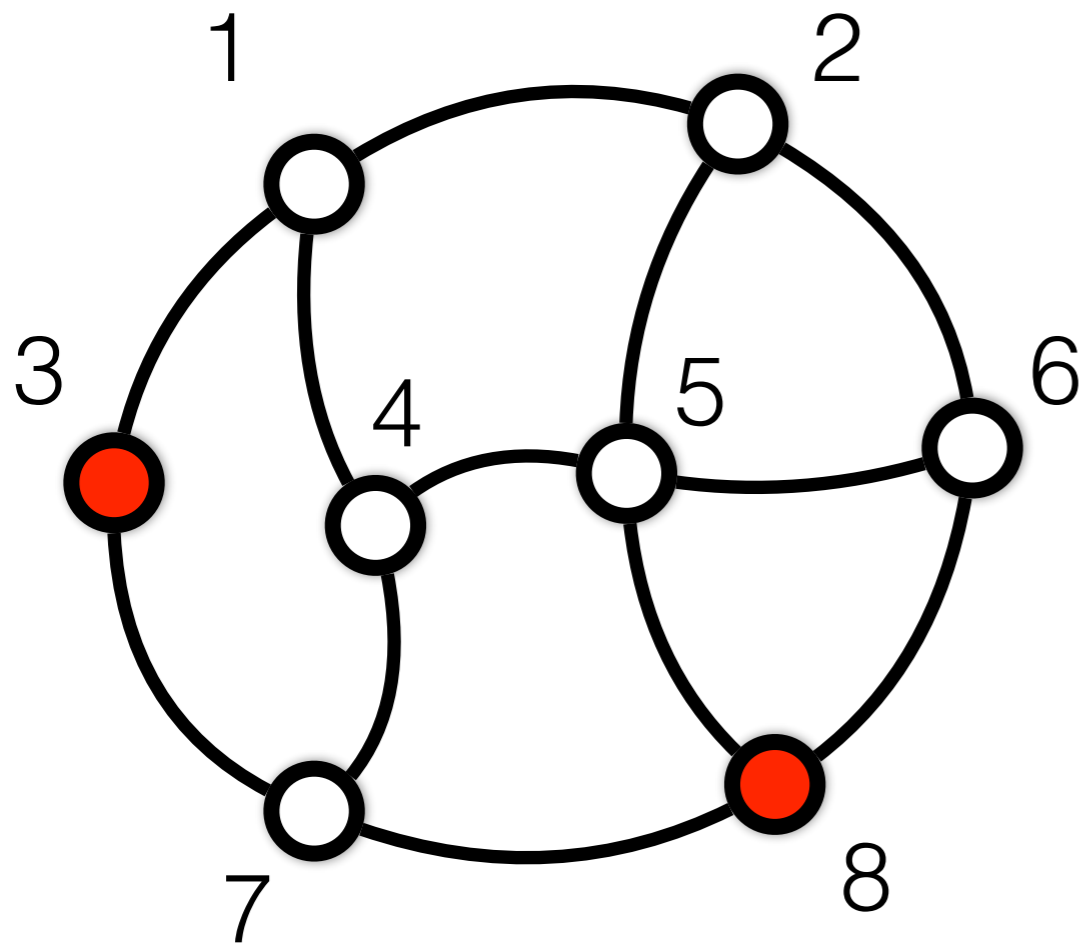
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0	0	1	0	0	0	0	0								
0	0	0	1	0	0	0	0								
0	0	0	0	1	0	0	0								
0	0	0	0	0	1	0	0								
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0	1	1	0	0	0	0	0								
0	1	0	1	0	0	0	0								
0	1	0	0	0	0	1	0								
0	1	0	0	0	0	0	1								

All Candidate Colorclasses



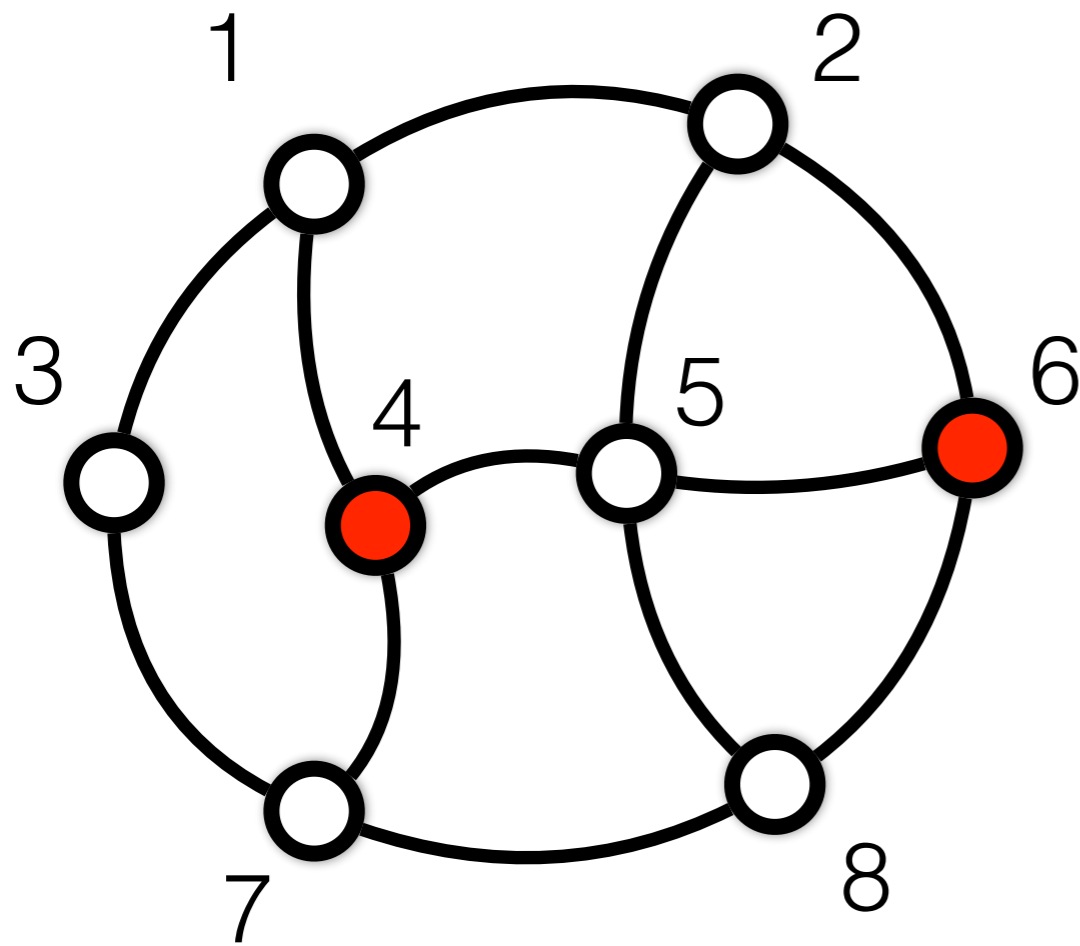
1	2	3	4	5	6	7	8	1	2	3	4	5	6	7	8
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0	1	0	0	0	0	0	1								

All Candidate Colorclasses



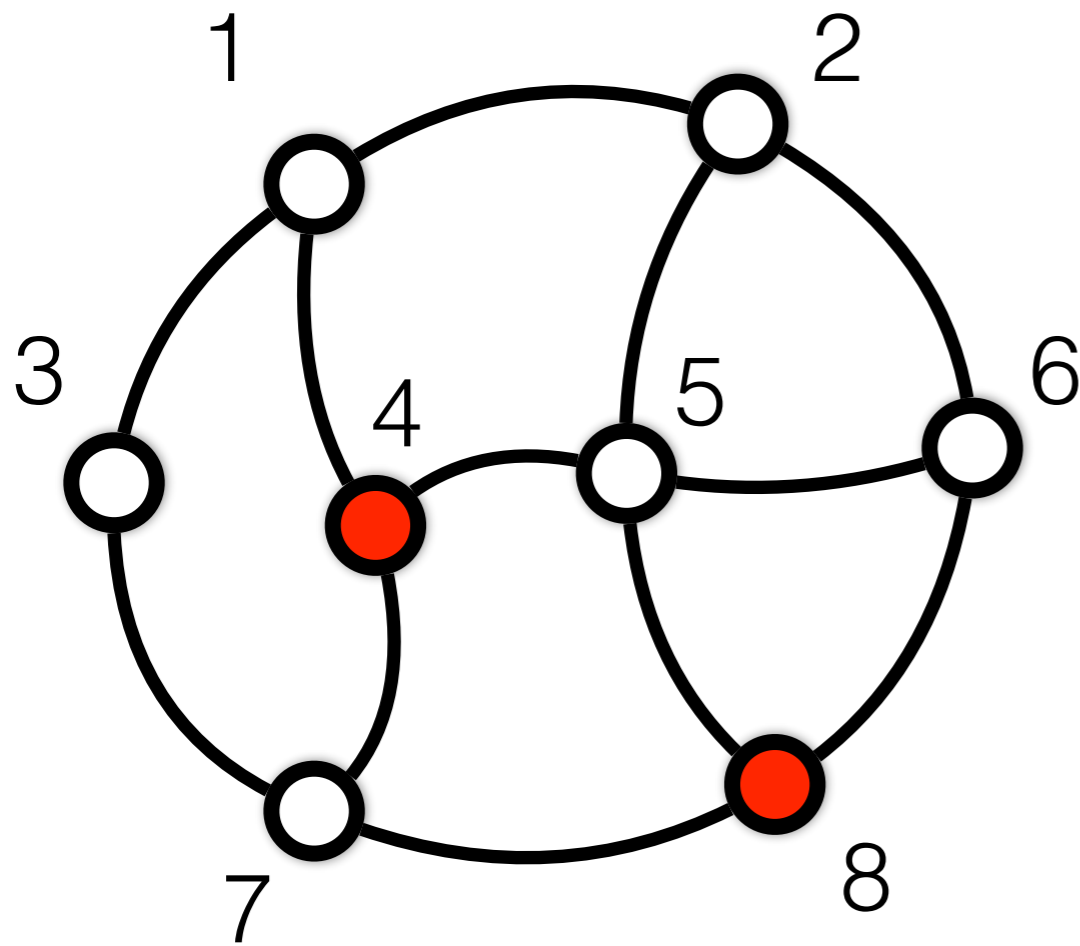
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0	0	0	0	1	0	0	0								
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0	1	0	0	0	0	0	1								

All Candidate Colorclasses



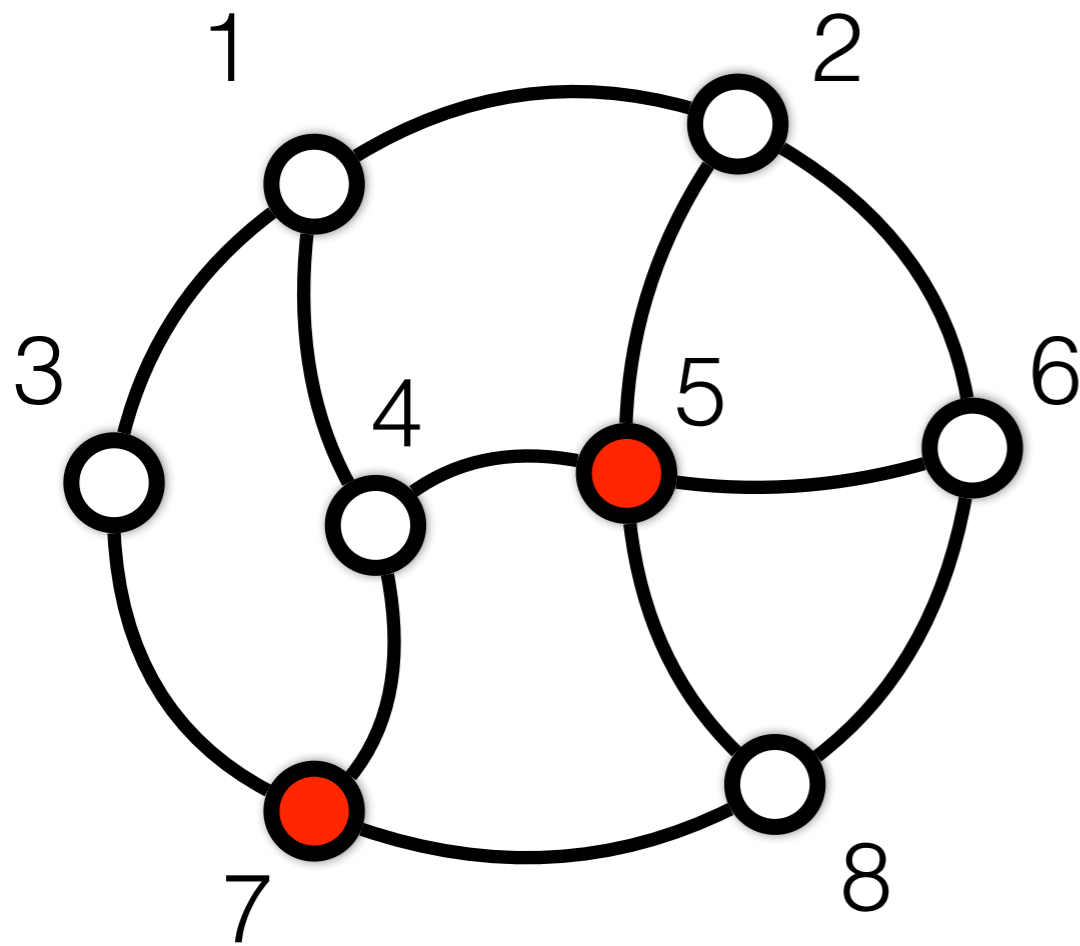
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All Candidate Colorclasses



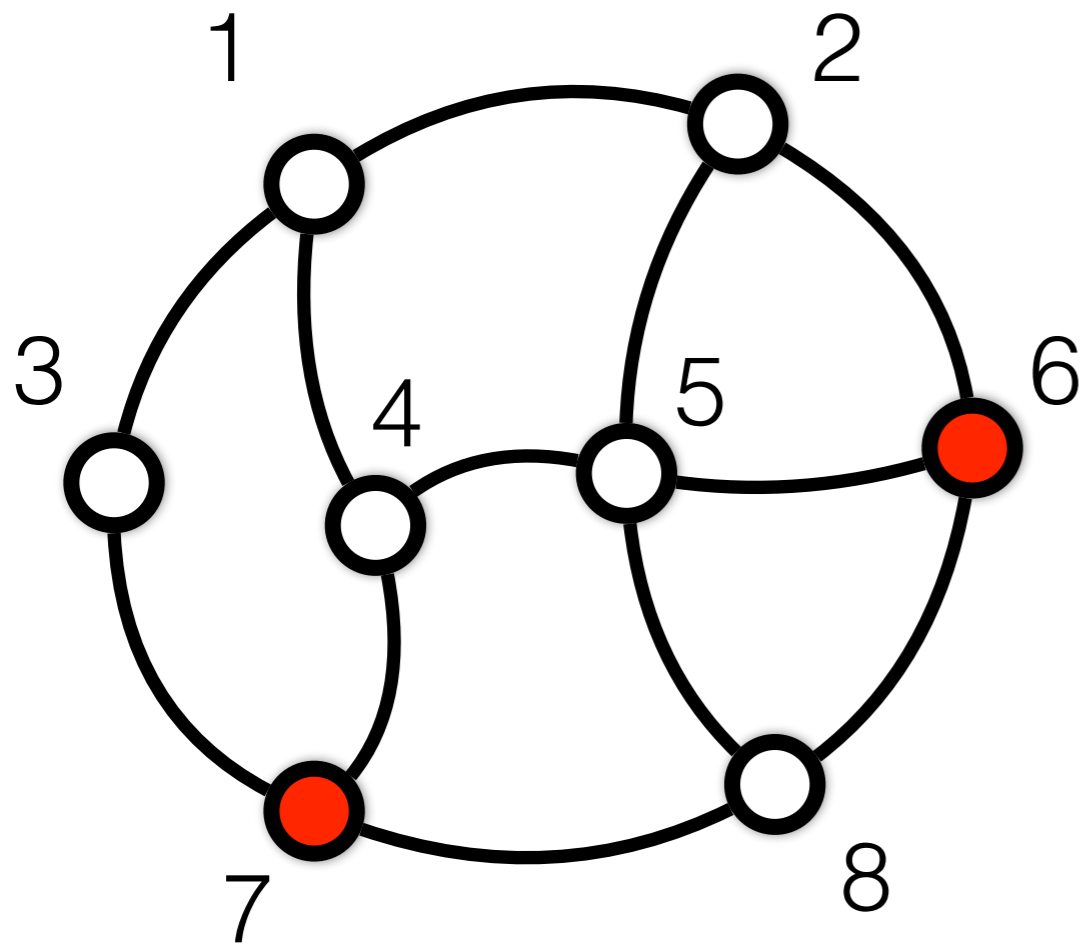
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All Candidate Colorclasses



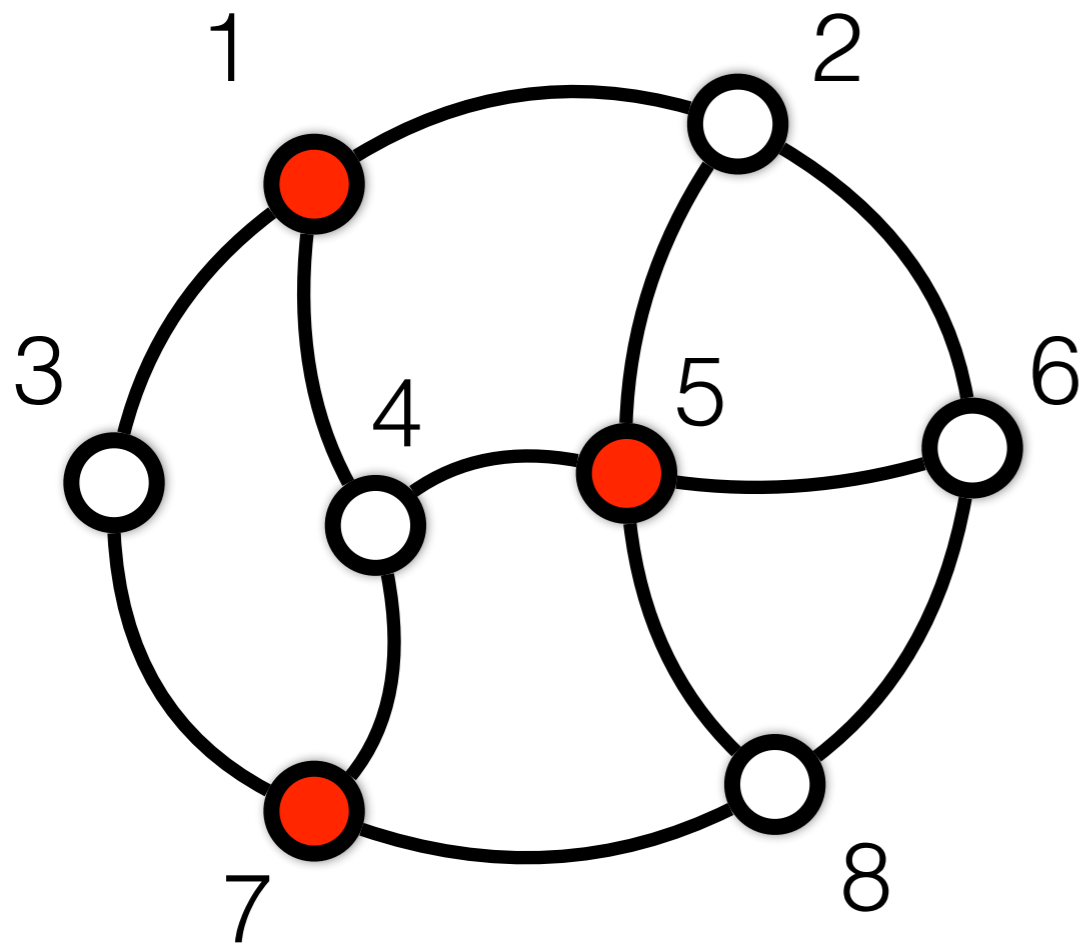
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All Candidate Colorclasses



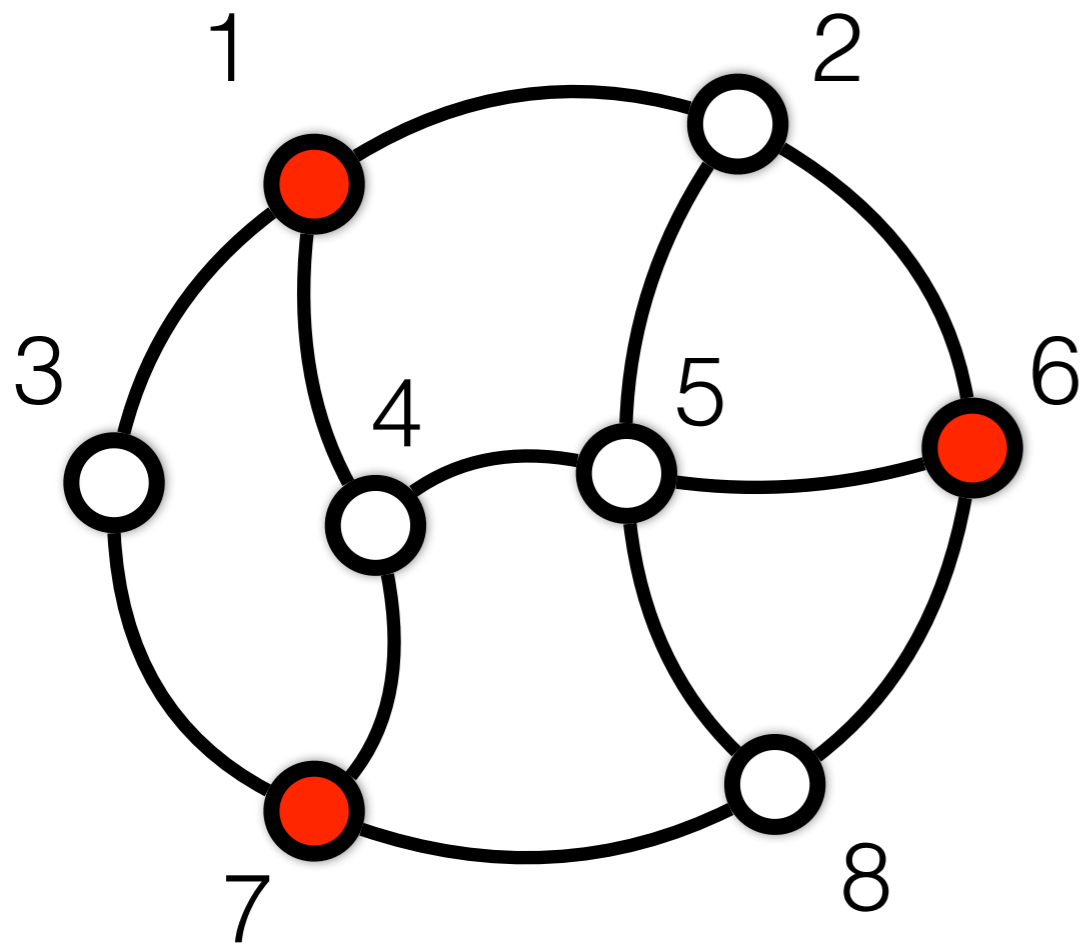
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All Candidate Colorclasses



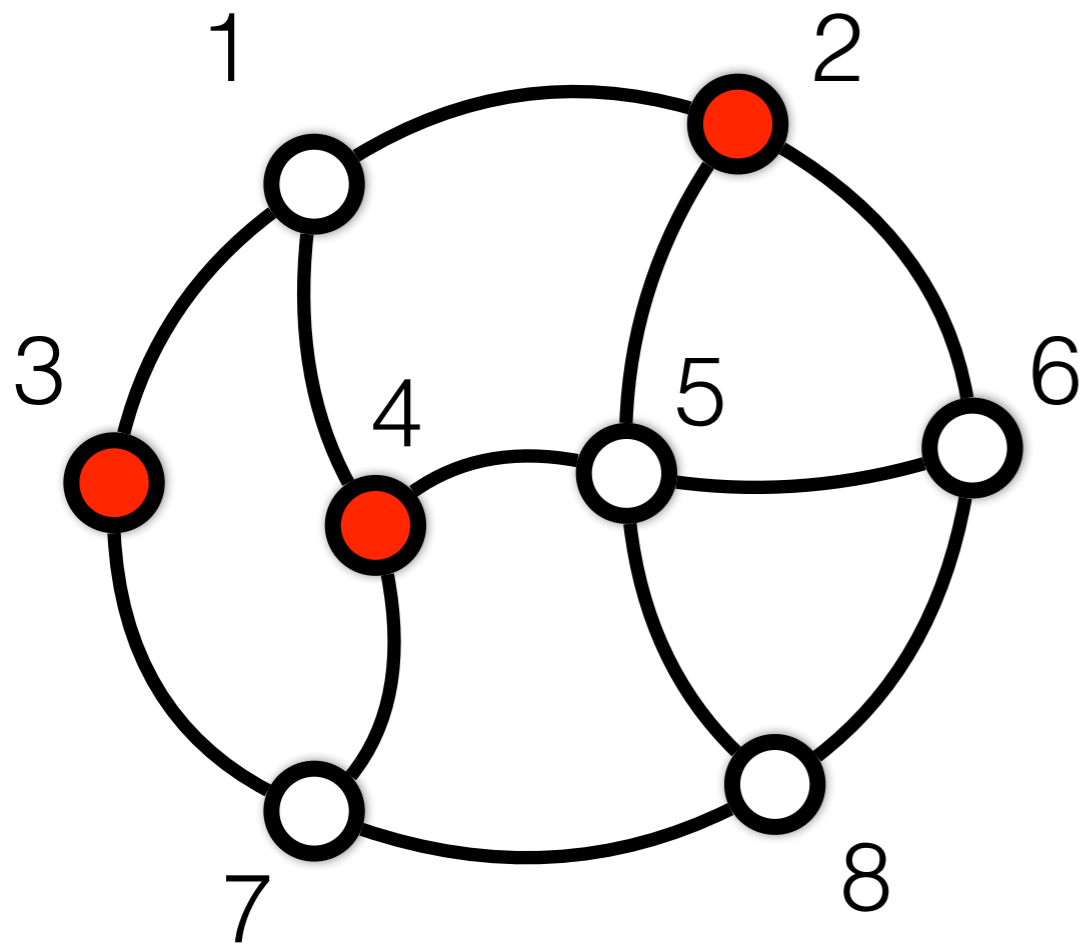
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0	1	0	1	0	0	0	0	0	1	0	0	0	0	0	0
0	1	0	0	0	0	1	0	0	1	0	0	0	0	1	0
0	1	0	0	0	0	0	1	0	1	0	0	0	0	0	1

All Candidate Colorclasses



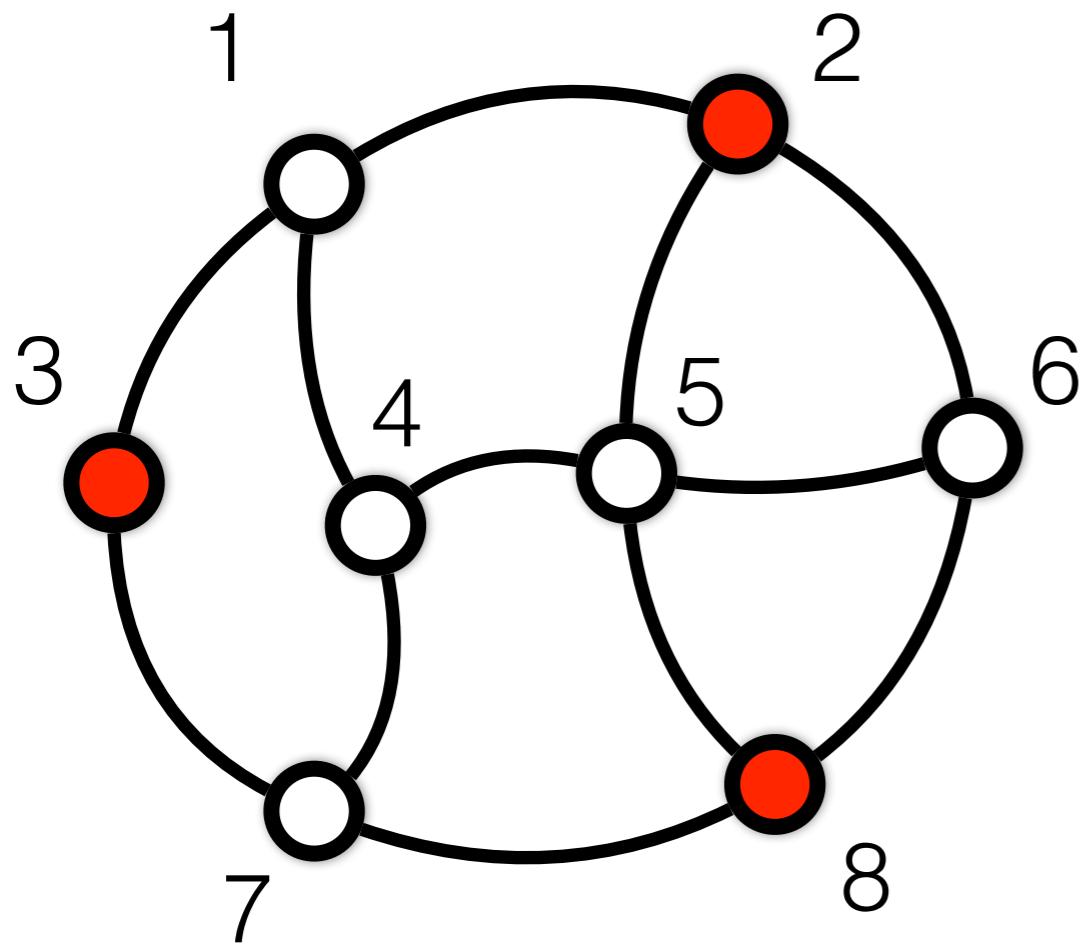
1	2	3	4	5	6	7	8	1	2	3	4	5	6	7	8
1	0	0	0	0	0	0	0	0	0	1	1	0	0	0	0
0	1	0	0	0	0	0	0	0	0	1	0	1	0	0	0
0	0	1	0	0	0	0	0	0	0	1	0	0	1	0	0
0	0	0	1	0	0	0	0	0	0	1	0	0	0	0	1
0	0	0	0	1	0	0	0	0	0	0	1	0	1	0	0
0	0	0	0	0	1	0	0	0	0	0	1	0	0	0	1
0	0	0	0	0	0	1	0	0	0	0	0	1	0	1	0
1	0	0	0	0	1	0	0	1	0	0	0	0	1	0	1
1	0	0	0	0	1	0	0	1	0	0	0	0	1	1	0
1	0	0	0	0	0	1	0	1	0	0	0	0	1	1	0
0	1	1	0	0	0	0	0	0	1	0	0	0	0	0	0
0	1	0	1	0	0	0	0	0	1	0	0	0	0	0	0
0	1	0	0	0	0	0	1	0	1	0	0	0	0	0	1

All Candidate Colorclasses



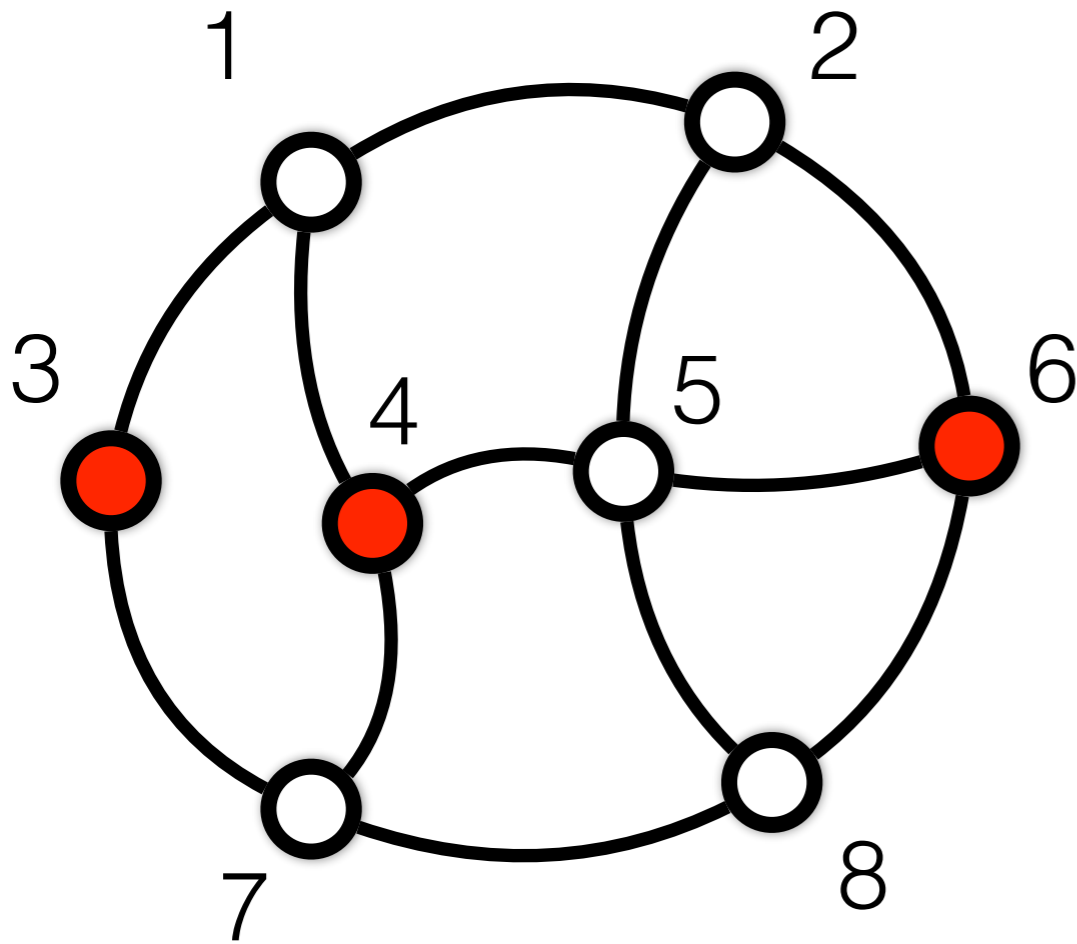
1	2	3	4	5	6	7	8	1	2	3	4	5	6	7	8
1	0	0	0	0	0	0	0	0	0	1	1	0	0	0	0
0	1	0	0	0	0	0	0	0	0	1	0	1	0	0	0
0	0	1	0	0	0	0	0	0	0	1	0	0	1	0	0
0	0	0	1	0	0	0	0	0	0	1	0	0	0	0	1
0	0	0	0	1	0	0	0	0	0	0	1	0	1	0	0
0	0	0	0	0	1	0	0	0	0	0	1	0	0	0	1
0	0	0	0	0	0	1	0	0	0	0	0	1	0	1	0
1	0	0	0	0	1	0	0	1	0	0	0	0	1	0	1
1	0	0	0	0	1	0	0	1	0	0	0	0	1	0	0
1	0	0	0	0	0	1	0	0	1	1	1	0	0	0	0
0	1	1	0	0	0	0	0								
0	1	0	1	0	0	0	0								
0	1	0	0	0	0	1	0								
0	1	0	0	0	0	0	1								

All Candidate Colorclasses



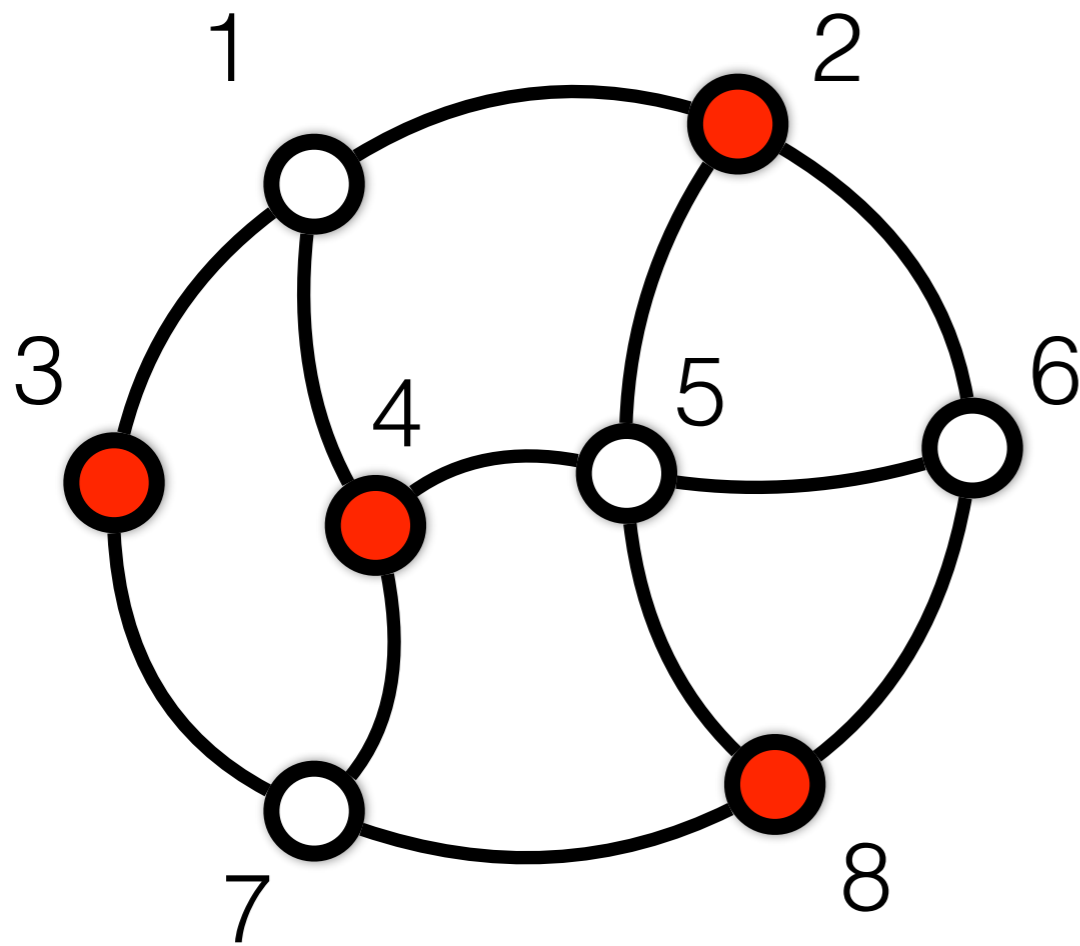
1	2	3	4	5	6	7	8	1	2	3	4	5	6	7	8
1	0	0	0	0	0	0	0	0	0	1	1	0	0	0	0
0	1	0	0	0	0	0	0	0	0	1	0	1	0	0	0
0	0	1	0	0	0	0	0	0	0	1	0	0	1	0	0
0	0	0	1	0	0	0	0	0	0	1	0	0	0	0	1
0	0	0	0	1	0	0	0	0	0	0	1	0	1	0	0
0	0	0	0	0	1	0	0	0	0	0	1	0	0	0	1
0	0	0	0	0	0	1	0	0	0	0	0	1	0	1	0
1	0	0	0	1	0	0	0	1	0	0	0	1	0	1	0
1	0	0	0	0	1	0	0	1	0	0	0	0	1	1	0
1	0	0	0	0	0	1	0	0	1	1	1	0	0	0	0
0	1	0	0	0	0	0	1	0	1	1	0	0	0	0	1
0	1	1	0	0	0	0	0								
0	1	0	1	0	0	0	0								
0	1	0	0	0	0	1	0								
0	1	0	0	0	0	0	1								

All Candidate Colorclasses



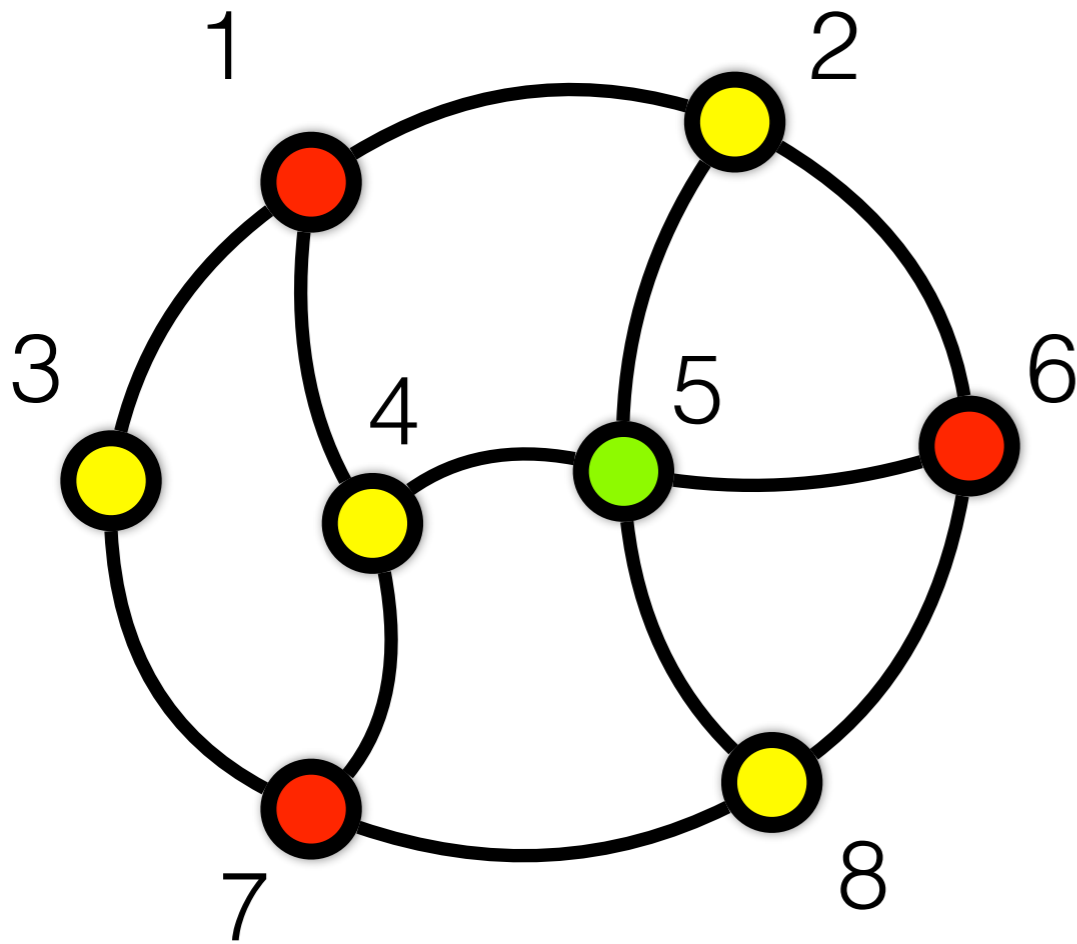
1	2	3	4	5	6	7	8	1	2	3	4	5	6	7	8
1	0	0	0	0	0	0	0	0	0	1	1	0	0	0	0
0	1	0	0	0	0	0	0	0	0	1	0	1	0	0	0
0	0	1	0	0	0	0	0	0	0	1	0	0	1	0	0
0	0	0	1	0	0	0	0	0	0	1	0	0	0	0	1
0	0	0	0	1	0	0	0	0	0	0	1	0	1	0	0
0	0	0	0	0	1	0	0	0	0	0	1	0	0	0	1
0	0	0	0	0	0	1	0	0	0	0	0	1	0	1	0
1	0	0	0	1	0	0	0	1	0	0	0	0	1	0	1
1	0	0	0	0	1	0	0	1	0	0	0	0	1	1	0
1	0	0	0	0	0	1	0	0	1	1	1	0	0	0	0
0	1	0	0	0	0	0	1	0	1	1	0	0	0	0	1
0	1	1	0	0	0	0	0	0	1	0	1	0	0	0	1
0	1	0	1	0	0	0	0	0	0	1	1	0	1	0	0
0	1	0	0	0	0	0	1	0	0	1	1	0	1	0	0

All Candidate Colorclasses



1	2	3	4	5	6	7	8	1	2	3	4	5	6	7	8
1	0	0	0	0	0	0	0	0	0	1	1	0	0	0	0
0	1	0	0	0	0	0	0	0	0	1	0	1	0	0	0
0	0	1	0	0	0	0	0	0	0	1	0	0	1	0	0
0	0	0	1	0	0	0	0	0	0	1	0	0	0	0	1
0	0	0	0	1	0	0	0	0	0	0	1	0	1	0	0
0	0	0	0	0	1	0	0	0	0	0	1	0	0	0	1
0	0	0	0	0	0	1	0	0	0	0	0	1	0	1	0
1	0	0	0	0	1	0	0	0	1	0	0	0	1	1	0
1	0	0	0	0	0	1	0	0	1	0	0	0	1	1	0
1	0	0	0	0	0	0	1	0	0	1	0	0	0	0	1
0	1	0	0	0	0	0	0	0	0	0	1	0	0	0	1
0	1	0	0	0	0	0	0	0	0	1	1	0	0	0	1
0	1	0	0	0	0	0	0	0	0	1	1	0	0	0	1
0	1	0	0	0	0	0	1	0	0	1	1	0	0	0	1

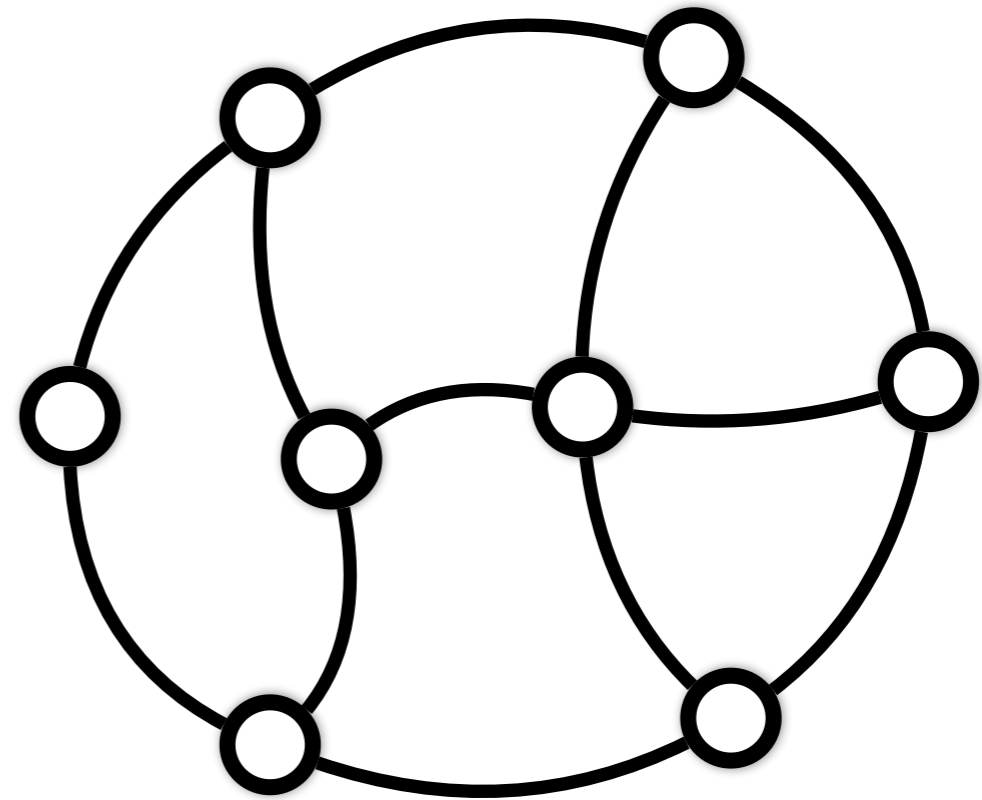
A Coloring is a Set of Disjoint Colorclasses



1	2	3	4	5	6	7	8	1	2	3	4	5	6	7	8
1	0	0	0	0	0	0	0	0	0	1	1	0	0	0	0
0	1	0	0	0	0	0	0	0	0	1	0	1	0	0	0
0	0	1	0	0	0	0	0	0	0	1	0	0	1	0	0
0	0	0	1	0	0	0	0	0	0	0	1	0	0	0	1
0	0	0	0	1	0	0	0	0	0	0	0	1	0	0	1
0	0	0	0	0	1	0	0	0	0	0	1	0	0	0	1
0	0	0	0	0	0	1	0	0	0	0	0	1	0	1	0
1	0	0	0	1	0	0	0	1	0	0	0	0	1	1	0
1	0	0	0	0	1	0	0	1	0	0	0	0	1	1	0
1	0	0	0	0	0	1	0	0	1	1	1	0	0	0	0
0	1	0	0	0	0	0	1	0	1	0	0	0	0	0	1
0	1	1	0	0	0	0	0	0	1	1	1	0	0	0	1
0	1	0	0	0	0	0	0	0	1	1	1	0	0	0	1
0	1	0	0	0	0	0	1	0	1	1	1	0	0	0	1

Dynamic Programming across Vertex Subsets

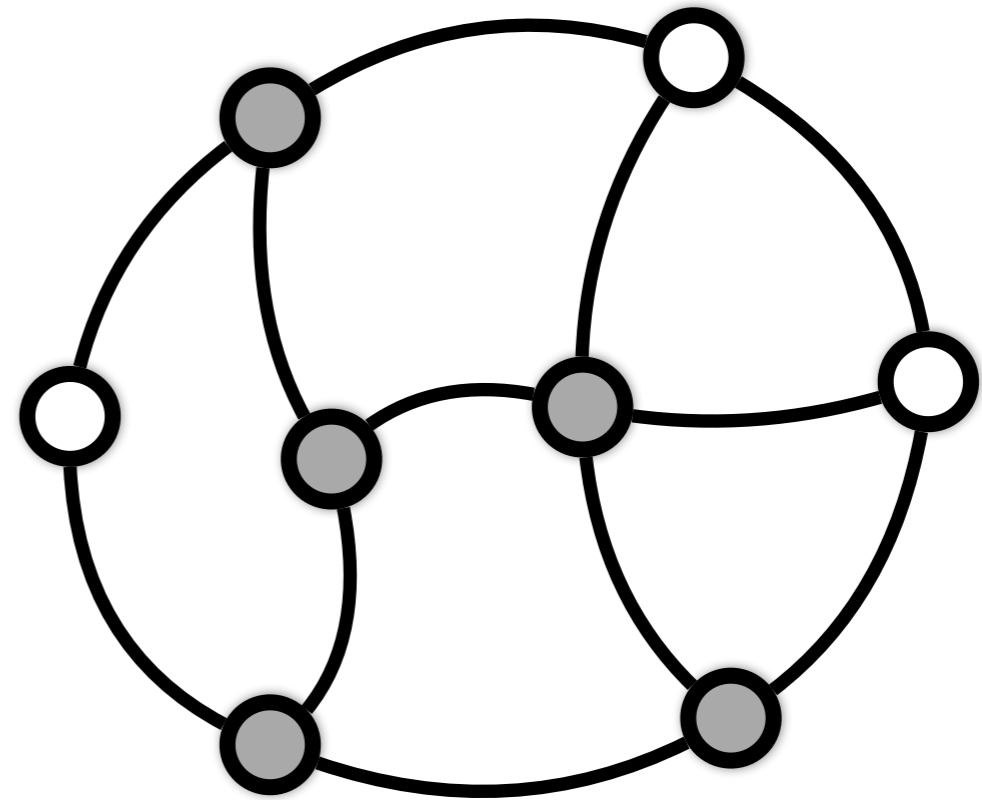
- For a vertex subset X , define $d(X)$ as the smallest number of colors needed in a proper coloring of $G[X]$, the graph *induced* by X .



The graph G .

Dynamic Programming across Vertex Subsets

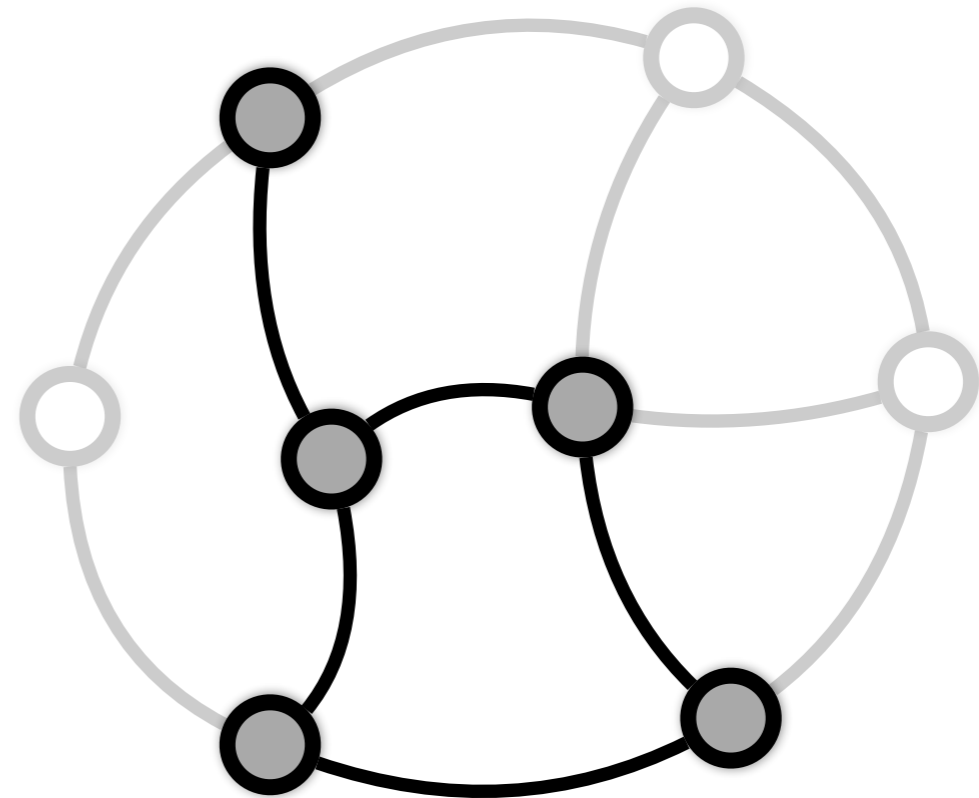
- For a vertex subset X , define $d(X)$ as the smallest number of colors needed in a proper coloring of $G[X]$, the graph *induced* by X .



X is the grey vertices.

Dynamic Programming across Vertex Subsets

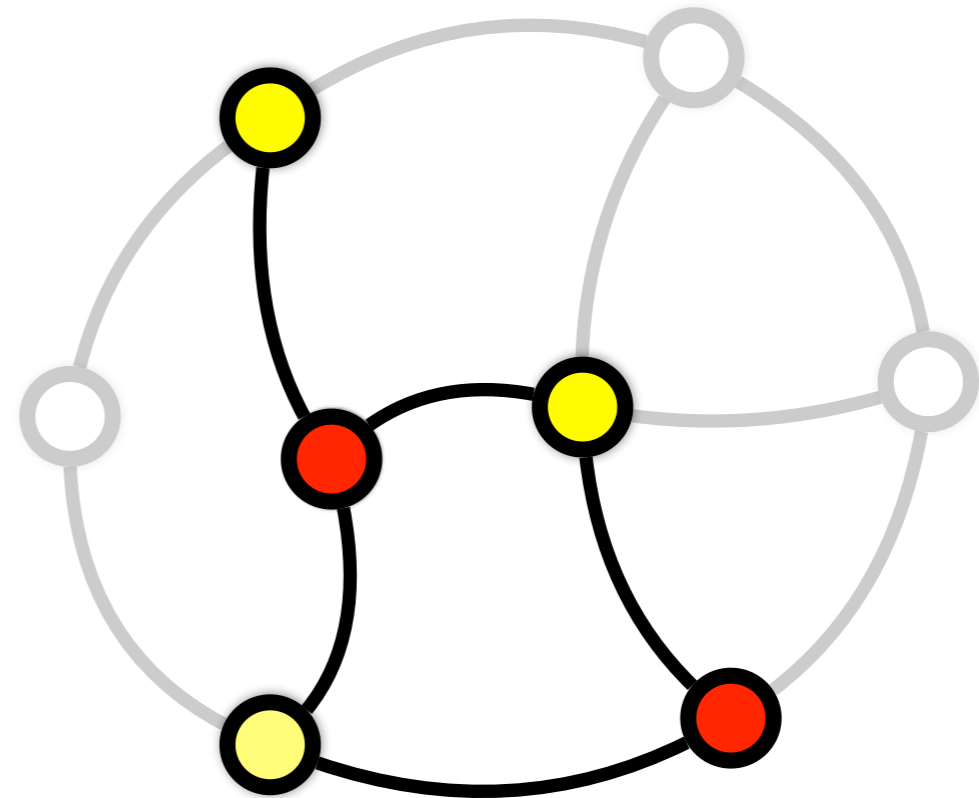
- For a vertex subset X , define $d(X)$ as the smallest number of colors needed in a proper coloring of $G[X]$, the graph *induced* by X .



The induced graph $G[X]$ is
 X and all edges between vertices
in X .

Dynamic Programming across Vertex Subsets

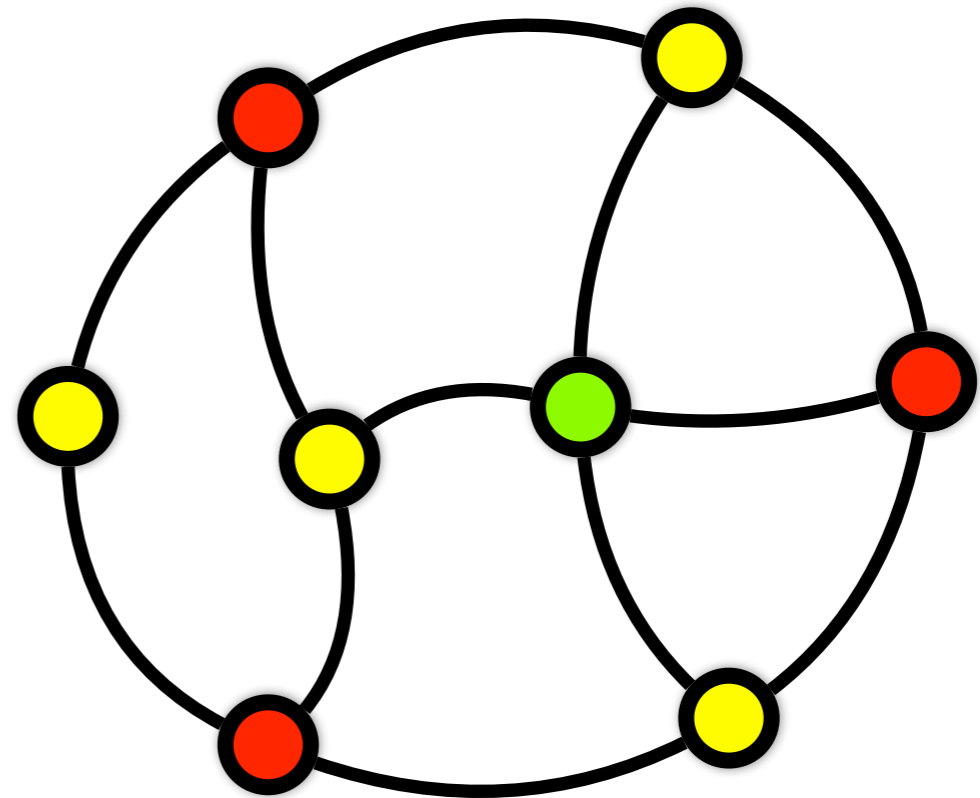
- For a vertex subset X , define $d(X)$ as the smallest number of colors needed in a proper coloring of $G[X]$, the graph *induced* by X .



$$d(X)=2.$$

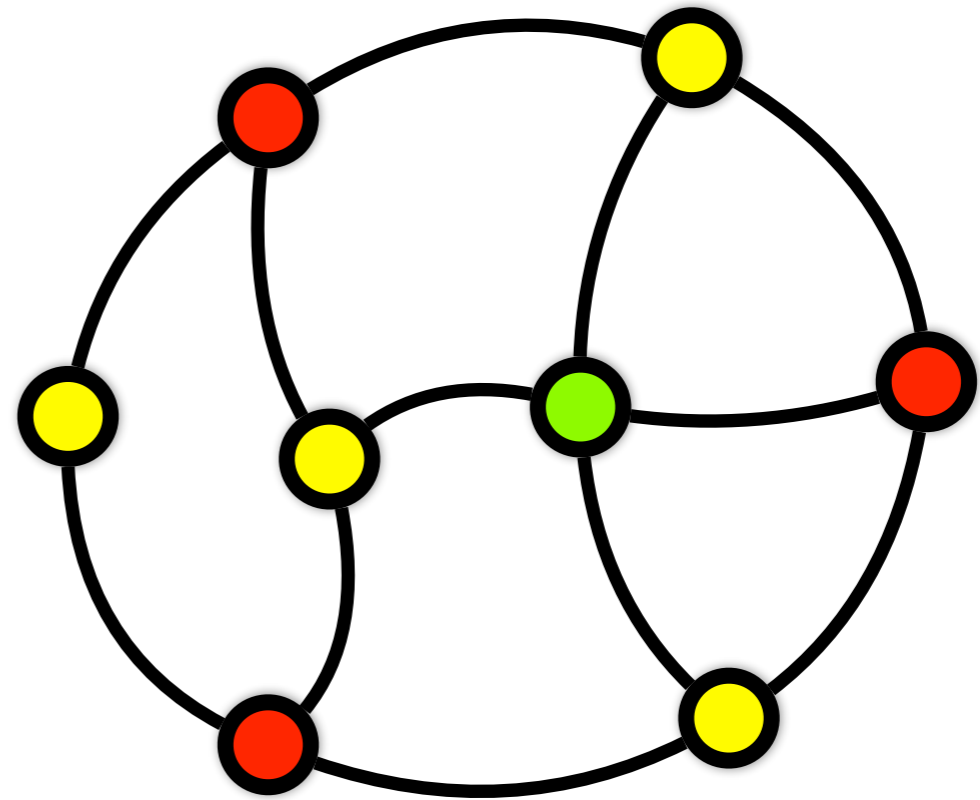
Dynamic Programming across Vertex Subsets

- $d(\emptyset) = 0$.
- $d(X) = \min_Y d(X - Y) + 1$.
Y is a colorclass
candidate in $G[X]$



Dynamic Programming across Vertex Subsets

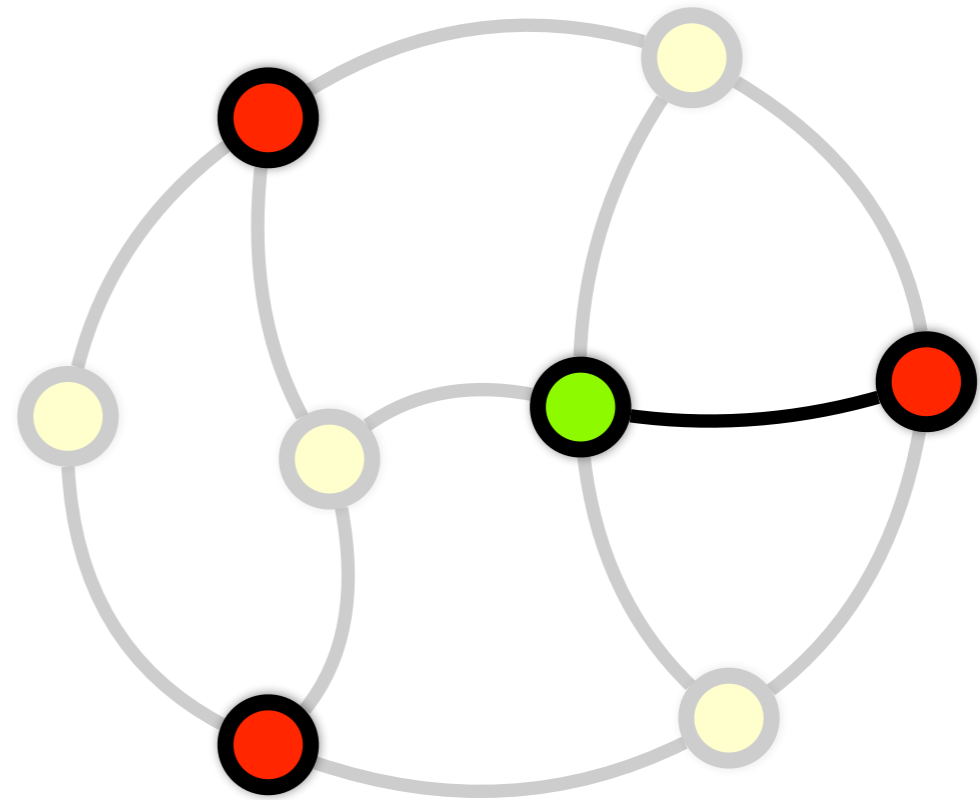
- Clearly,
 $d(V) = \chi(G)$.



$$d(V) = 3.$$

Dynamic Programming across Vertex Subsets

- $d(0)=0$.
- $d(X)=\min_Y d(X-Y)+1$.
Y is a colorclass
candidate in $G[X]$



$X=V$, $Y=Yellow$ colorclass

Running Time Analysis

$$t(n) \leq \sum_{X \subseteq V} \sum_{Y \subseteq X} 1 = \sum_{X \subseteq V} 2^{|X|} = \sum_{k=0}^n \binom{n}{k} 2^k = 3^n.$$

Overestimate of
candidate colorclasses.

Improved Bounds for the Dynamic Programming Approach

- [Lawler 1976] $O(2.44^n)$ by looping over maximal independent sets as colorclasses.
- [Byskov 2003] $O(2.40^n)$ by more careful analysis.

$2^n \text{poly}(n)$ time algorithm for chromatic number

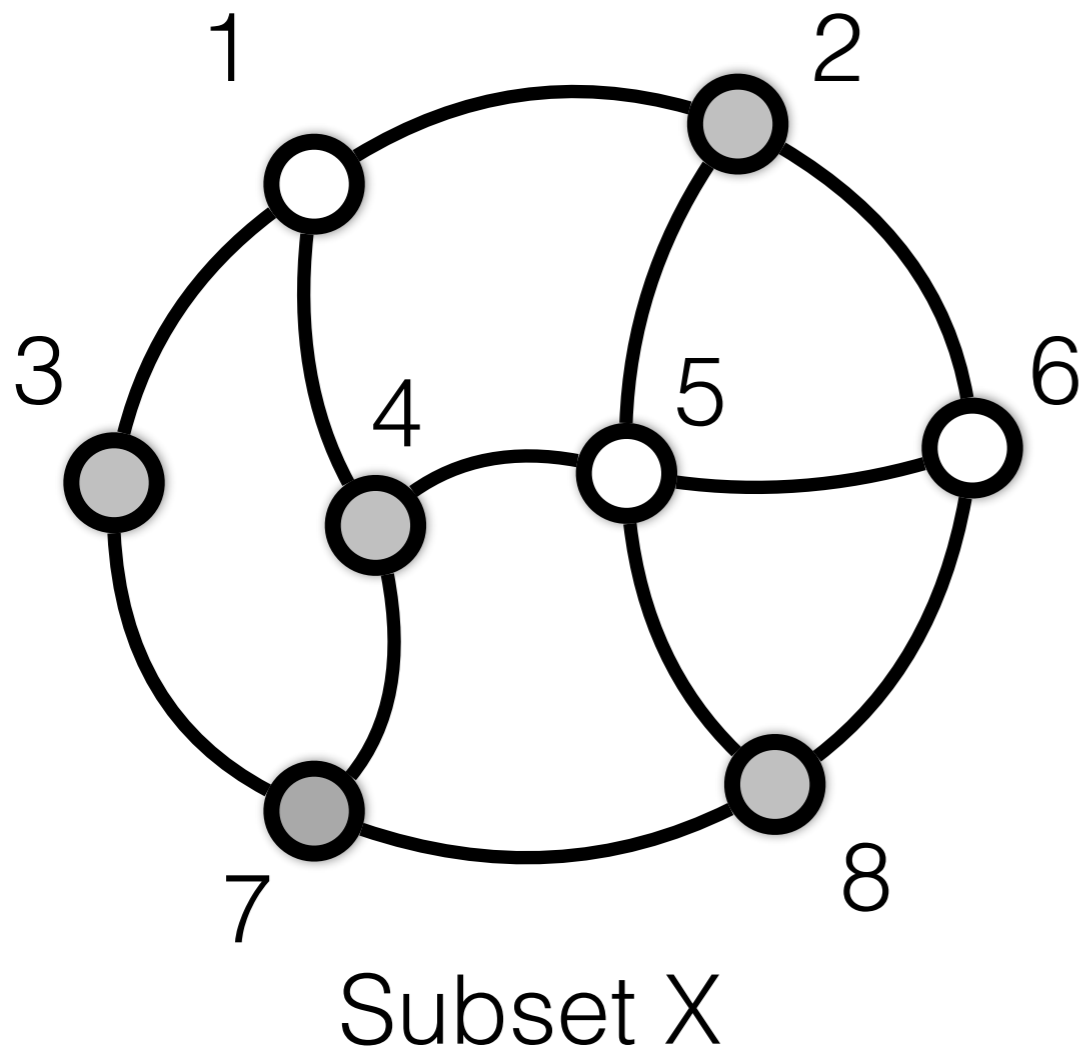
[B., Husfeldt, and Koivisto 2006]

- Use inclusion-exclusion summation to count k -colorings.
- Use the fast zeta transform (a variant of the Fast Fourier transform) to efficiently count candidate colorclasses in every induced subgraph $G[X]$ at once.

Subset Induced Colorclasses

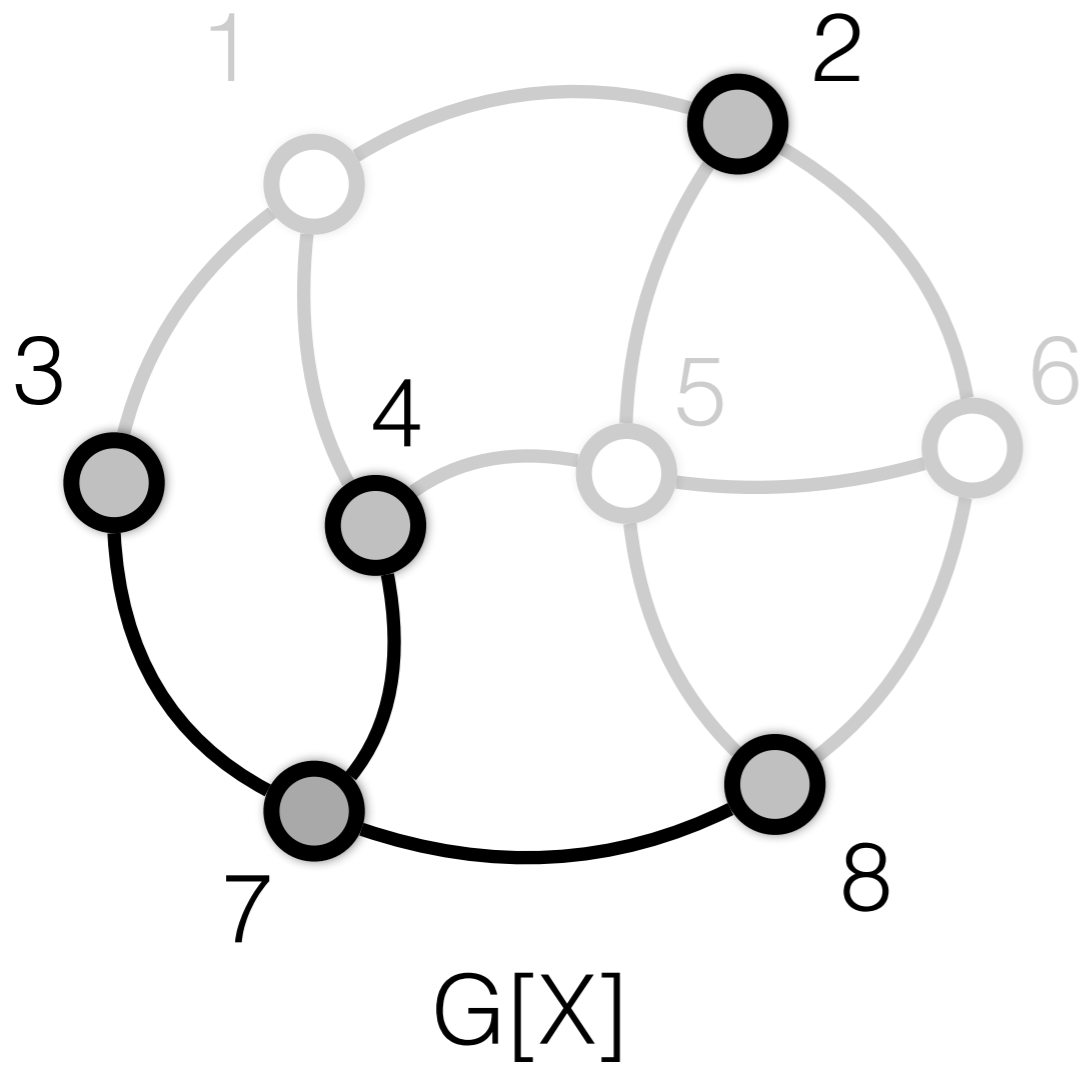
- Let $a(X)$ for X a subset of the vertices V be the number of candidate color classes in the induced graph $G[X]$.

Subset Induced Colorclasses



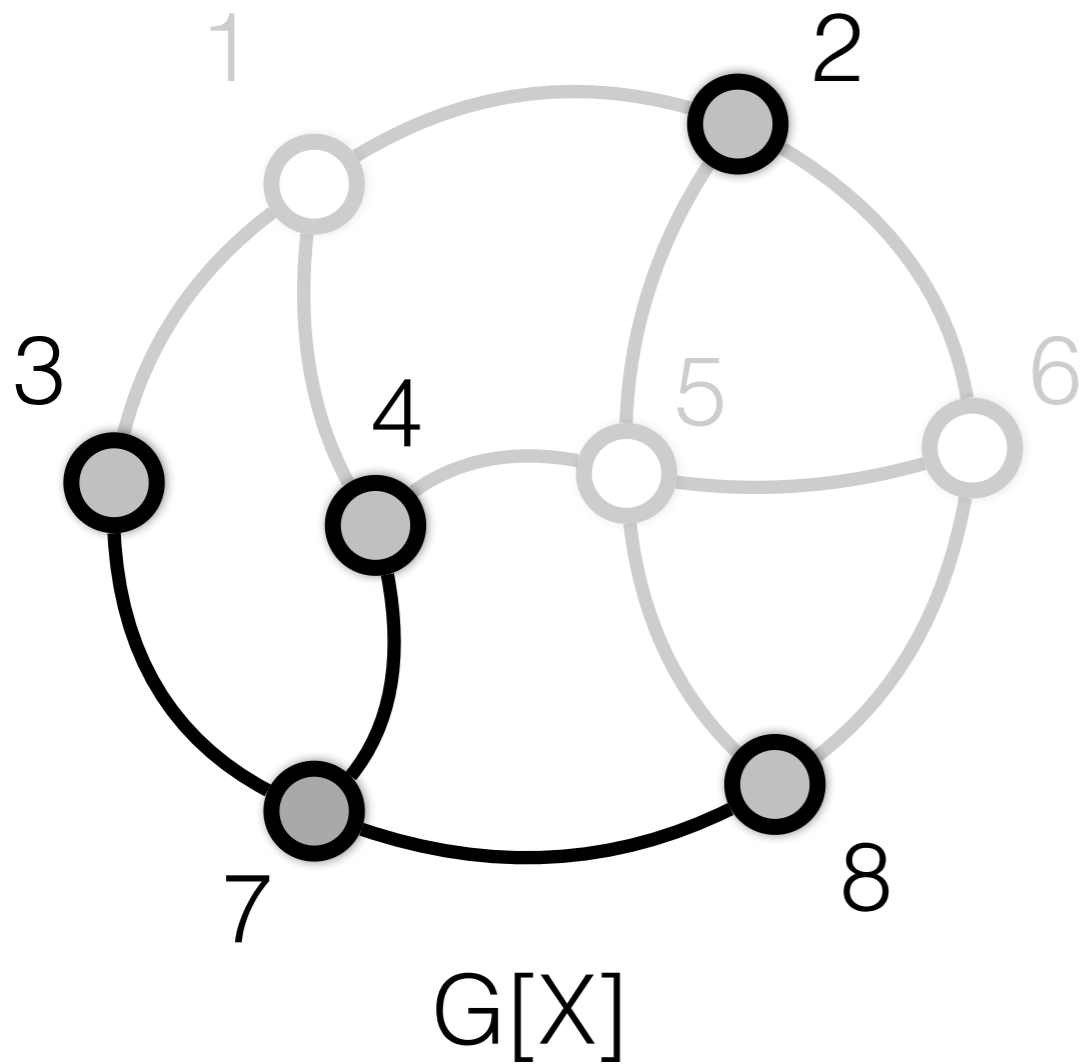
1	2	3	4	5	6	7	8	1	2	3	4	5	6	7	8
1	0	0	0	0	0	0	0	0	0	1	1	0	0	0	0
0	1	0	0	0	0	0	0	0	0	1	0	1	0	0	0
0	0	1	0	0	0	0	0	0	0	1	0	0	1	0	0
0	0	0	1	0	0	0	0	0	0	1	0	0	0	0	1
0	0	0	0	1	0	0	0	0	0	0	1	0	1	0	0
0	0	0	0	0	1	0	0	0	0	0	1	0	0	0	1
0	0	0	0	0	0	1	0	0	0	0	0	1	0	1	0
1	0	0	0	1	0	0	0	1	0	0	0	0	1	0	1
1	0	0	0	0	1	0	0	1	0	0	0	0	1	1	0
1	0	0	0	0	0	1	0	0	1	1	1	0	0	0	0
0	1	0	0	0	0	0	1	0	1	1	0	0	0	0	1
0	1	1	0	0	0	0	0	0	1	0	1	0	0	0	1
0	1	0	1	0	0	0	0	0	0	1	1	0	1	0	0
0	1	0	0	0	0	1	0	0	0	1	1	0	0	0	1
0	1	0	0	0	0	0	1	0	1	1	1	0	0	0	1

Subset Induced Colorclasses



1	2	3	4	5	6	7	8	1	2	3	4	5	6	7	8
1	0	0	0	0	0	0	0	0	0	1	1	0	0	0	0
0	1	0	0	0	0	0	0	0	0	1	0	1	0	0	0
0	0	1	0	0	0	0	0	0	0	1	0	0	1	0	0
0	0	0	1	0	0	0	0	0	0	1	0	0	0	0	1
0	0	0	0	1	0	0	0	0	0	0	1	0	1	0	0
0	0	0	0	0	1	0	0	0	0	0	1	0	0	0	1
0	0	0	0	0	0	1	0	0	0	0	0	1	0	1	0
1	0	0	0	1	0	0	0	1	0	0	0	0	1	0	1
1	0	0	0	0	1	0	0	1	0	0	0	0	1	1	0
1	0	0	0	0	0	1	0	0	1	1	1	0	0	0	0
0	1	0	0	0	0	0	1	0	1	1	0	0	0	0	1
0	1	1	0	0	0	0	0	0	1	0	1	0	0	0	1
0	1	0	1	0	0	0	0	0	0	1	1	0	1	0	0
0	1	0	0	0	0	1	0	0	0	1	1	0	0	0	1
0	1	0	0	0	0	0	1	0	1	1	1	0	0	0	1

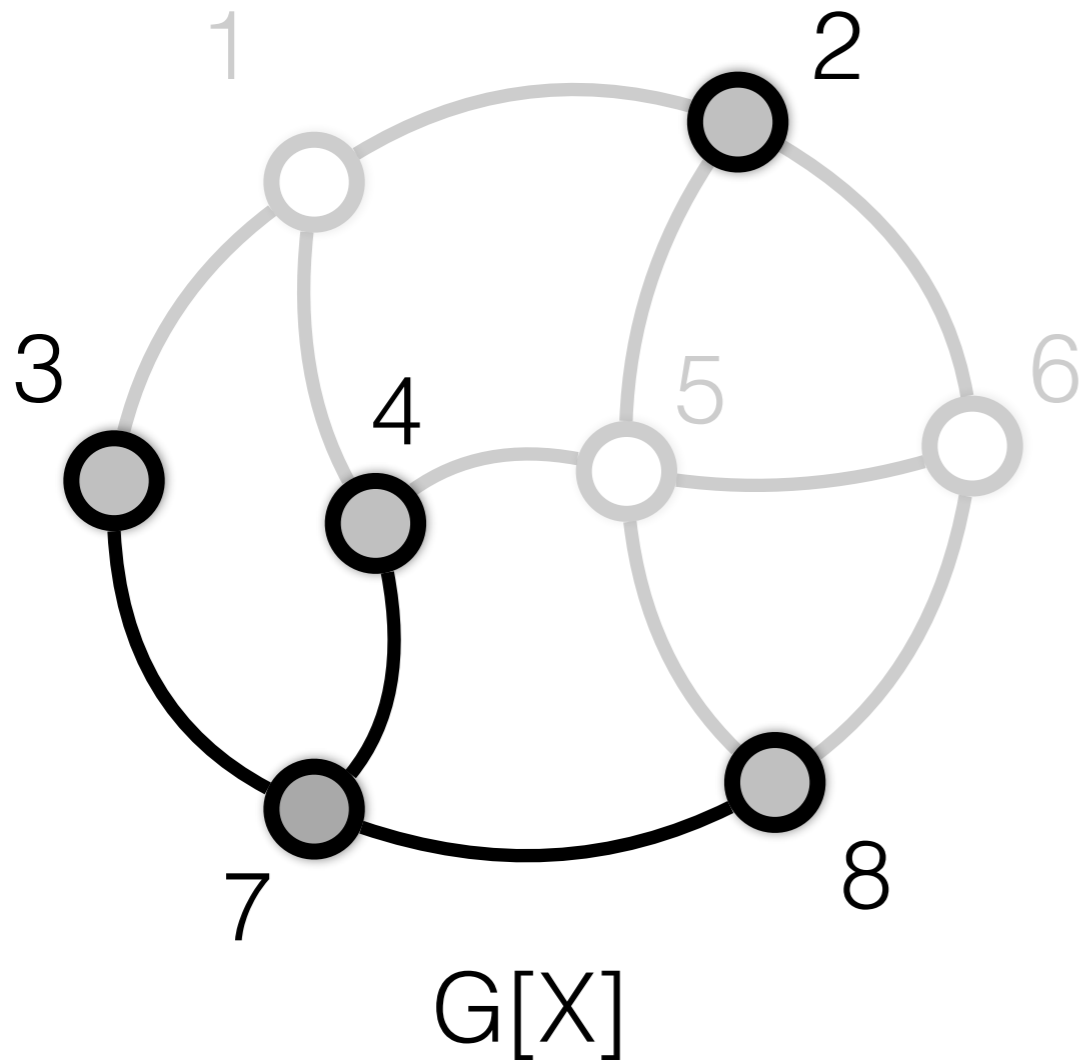
Subset Induced Colorclasses



1	2	3	4	5	6	7	8	1	2	3	4	5	6	7	8
1	0	0	0	0	0	0	0	0	0	1	1	0	0	0	0
0	1	0	0	0	0	0	0	0	0	1	0	1	0	0	0
0	0	1	0	0	0	0	0	0	0	1	0	0	1	0	0
0	0	0	1	0	0	0	0	0	0	1	0	0	0	0	1
0	0	0	0	1	0	0	0	0	0	0	1	0	1	0	0
0	0	0	0	0	1	0	0	0	0	0	1	0	0	0	1
0	0	0	0	0	0	1	0	0	0	0	0	1	0	1	0
0	0	0	0	0	0	0	1	0	0	0	0	0	1	1	0
1	0	0	0	1	0	0	0	1	0	0	0	1	0	1	0
1	0	0	0	0	1	0	0	1	0	0	0	0	1	1	0
1	0	0	0	0	0	1	0	0	1	1	1	0	0	0	0
0	1	0	1	0	0	0	0	0	1	0	1	0	0	0	1
0	1	0	0	0	0	0	0	0	0	1	1	0	1	0	0
0	1	0	0	0	0	0	1	0	1	1	1	0	0	0	1

Subset Induced Colorclasses

$$a(X) = 17$$



1	2	3	4	5	6	7	8	1	2	3	4	5	6	7	8
1	0	0	0	0	0	0	0	0	0	1	1	0	0	0	0
0	1	0	0	0	0	0	0	0	0	1	0	1	0	0	0
0	0	1	0	0	0	0	0	0	0	1	0	0	1	0	0
0	0	0	1	0	0	0	0	0	0	1	0	0	0	0	1
0	0	0	0	1	0	0	0	0	0	0	1	0	1	0	0
0	0	0	0	0	1	0	0	0	0	0	1	0	0	0	1
0	0	0	0	0	0	1	0	0	0	0	0	1	0	1	0
1	0	0	0	1	0	0	0	1	0	0	0	1	0	1	0
1	0	0	0	0	1	0	0	1	0	0	0	1	0	1	0
1	0	0	0	0	0	1	0	0	1	1	1	0	0	0	0
0	1	0	1	0	0	0	0	0	1	1	0	0	0	0	1
0	1	0	0	0	0	0	0	0	0	1	1	0	1	0	0
0	1	0	0	0	0	0	1	0	1	1	1	0	0	0	1

Inclusion-Exclusion

$$p(k) = \sum_{X \subseteq V} (-1)^{|V-X|} a(X)^k.$$

- $p(k)$ is zero if there is no k -coloring,
- $p(k)$ is non-zero if there are k -colorings.

Meaning of Powers of $a(X)$

- $a(X)^k$ counts the number of ways to pick k color classes (with repetition) in $G[X]$.

1	2	3	4	5	6	7	8	1	2	3	4	5	6	7	8
1	0	0	0	0	0	0	0	0	0	1	1	0	0	0	0
0	1	0	0	0	0	0	0	0	0	1	0	1	0	0	0
0	0	1	0	0	0	0	0	0	0	1	0	0	1	0	0
0	0	0	1	0	0	0	0	0	0	1	0	0	0	0	1
0	0	0	0	1	0	0	0	0	0	0	1	0	1	0	0
0	0	0	0	0	1	0	0	0	0	0	1	0	0	0	1
0	0	0	0	0	0	1	0	0	0	0	0	1	0	1	0
0	0	0	0	0	0	0	1	0	0	0	0	0	1	1	0
1	0	0	0	1	0	0	0	1	0	0	0	1	0	1	0
1	0	0	0	0	1	0	0	1	0	0	0	0	1	1	0
1	0	0	0	0	0	1	0	0	1	1	1	0	0	0	0
1	0	0	0	0	0	0	1	0	1	1	0	0	0	0	1
0	1	1	0	0	0	0	0	0	1	0	1	0	0	0	1
0	1	0	1	0	0	0	0	0	0	1	1	0	1	0	0
0	1	0	0	0	0	0	1	0	1	1	1	0	0	0	1

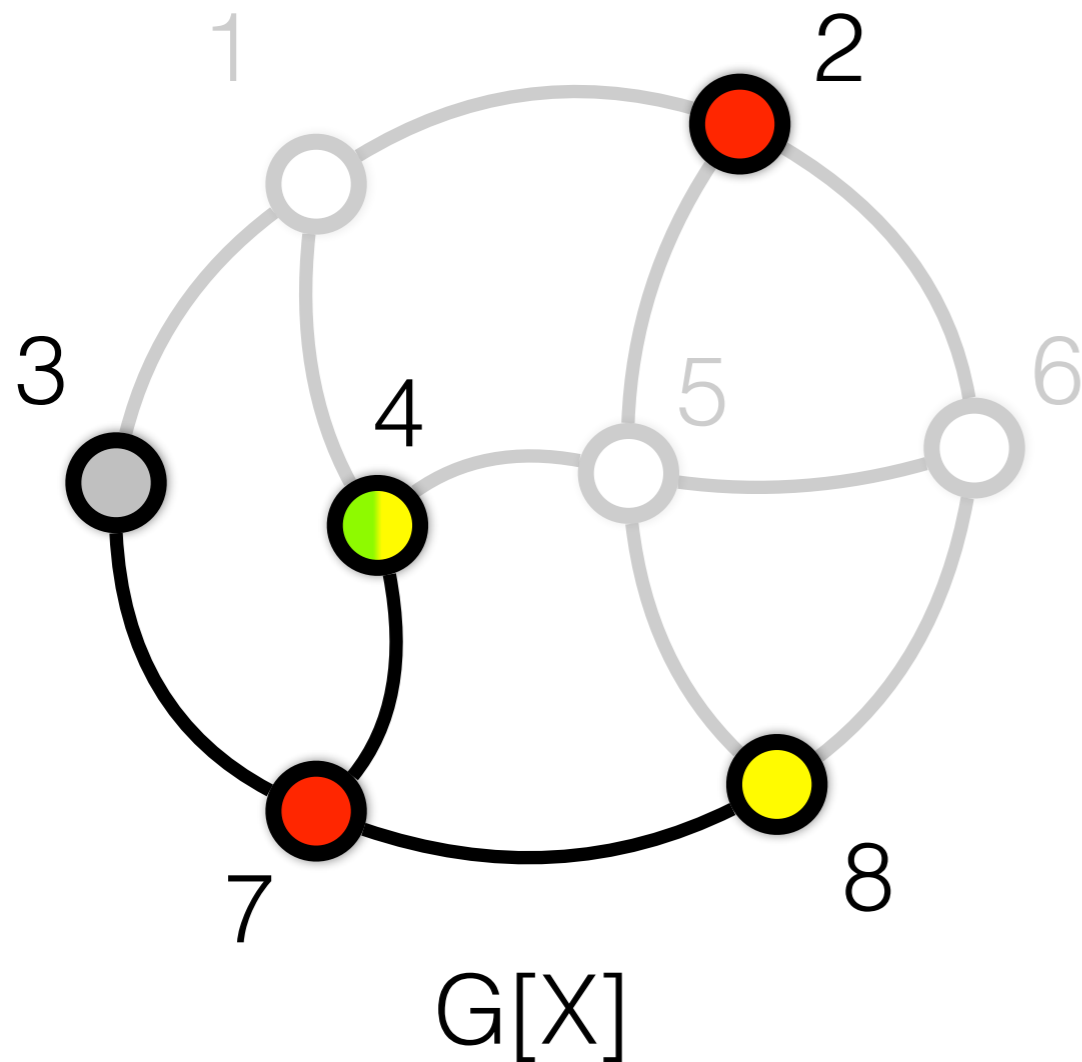
Meaning of Powers of $a(X)$

$k=3$

- $a(X)^k$ counts the number of ways to pick k color classes (with repetition) in $G[X]$.

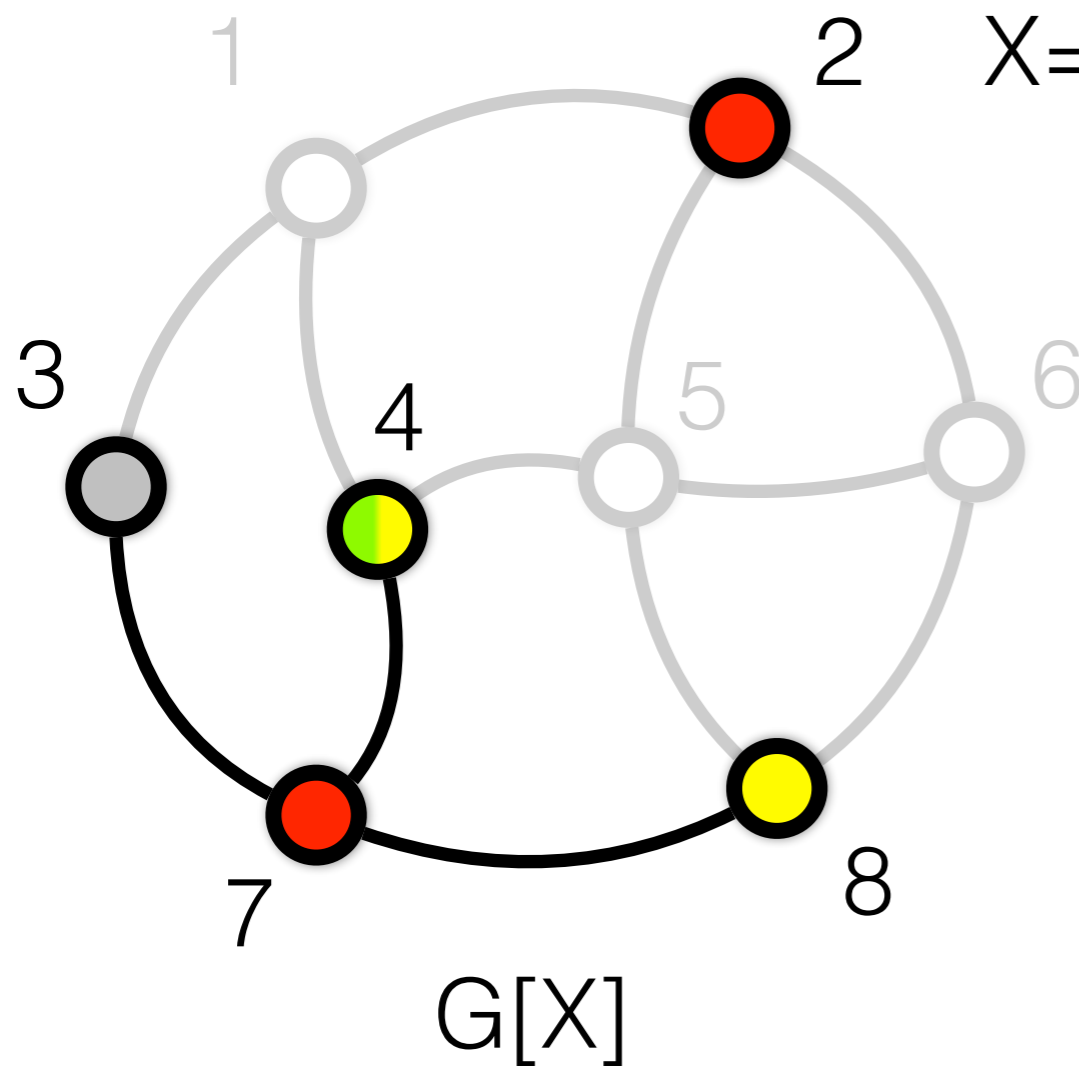
1	2	3	4	5	6	7	8	1	2	3	4	5	6	7	8
1	0	0	0	0	0	0	0	0	0	1	1	0	0	0	0
0	1	0	0	0	0	0	0	0	0	1	0	1	0	0	0
0	0	1	0	0	0	0	0	0	0	1	0	0	1	0	0
0	0	0	1	0	0	0	0	0	0	1	0	0	0	0	1
0	0	0	0	0	1	0	0	0	0	0	1	0	0	0	1
0	0	0	0	0	0	1	0	0	0	0	0	1	0	1	0
0	0	0	0	0	0	0	1	0	0	0	0	0	1	1	0
1	0	0	0	1	0	0	0	1	0	0	0	1	0	1	0
1	0	0	0	0	1	0	0	1	0	0	0	0	1	1	0
1	0	0	0	0	0	1	0	0	1	1	1	0	0	0	0
1	0	0	0	0	0	0	1	0	1	1	0	0	0	0	1
0	1	1	0	0	0	0	0	0	1	0	1	0	0	0	1
0	1	0	1	0	0	0	0	0	0	1	1	0	1	0	0
0	1	0	0	0	0	1	0	0	0	1	1	0	0	0	1
0	1	0	0	0	0	0	1	0	1	1	1	0	0	0	1

Meaning of Powers of $a(X)$



1	2	3	4	5	6	7	8	1	2	3	4	5	6	7	8
1	0	0	0	0	0	0	0	0	0	1	1	0	0	0	0
0	1	0	0	0	0	0	0	0	0	1	0	1	0	0	0
0	0	1	0	0	0	0	0	0	0	1	0	0	1	0	0
0	0	0	1	0	0	0	0	0	0	0	1	0	1	0	0
0	0	0	0	0	1	0	0	0	0	0	1	0	0	0	1
0	0	0	0	0	0	1	0	0	0	0	0	1	0	1	0
0	0	0	0	0	0	0	1	0	0	0	0	0	1	1	0
1	0	0	0	1	1	0	0	1	0	0	0	1	0	1	0
1	0	0	0	0	1	0	0	1	0	0	0	1	0	1	0
1	0	0	0	0	0	1	0	0	1	1	1	0	0	0	0
1	0	0	0	0	0	0	1	0	1	1	0	0	0	0	1
0	1	1	0	0	0	0	0	0	1	0	1	0	0	0	1
0	1	0	1	0	0	0	0	0	0	1	1	0	1	0	0
0	1	0	0	0	0	1	0	0	1	1	1	0	0	0	1
0	1	0	0	0	0	0	1	0	1	1	1	0	0	0	1

A k -tuple of Colorclasses will be Counted in Many $G[X]$'s



$X = 1\ 2\ 3\ 4\ 5\ 6\ 7\ 8$

0 0 0 1 0 0 0 0

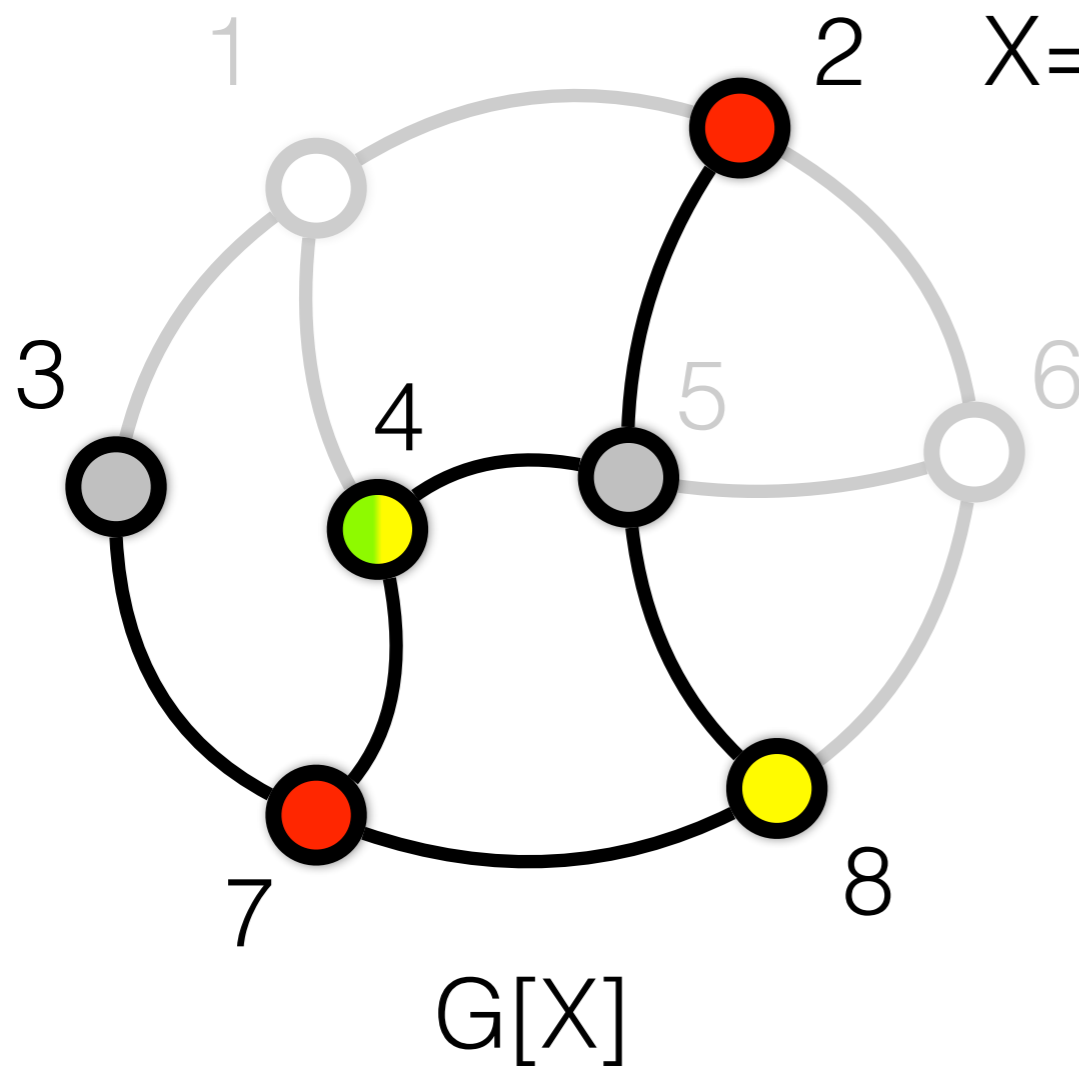
0 1 0 0 0 0 1 0

0 0 0 1 0 0 0 1

Sign = $(-1)^{8-3} = -1$

$$p(k) = \sum_{X \subseteq V} (-1)^{|V-X|} a(X)^k.$$

A k -tuple of Colorclasses will be Counted in Many $G[X]$'s



$X = 1\ 2\ 3\ 4\ 5\ 6\ 7\ 8$

0 0 0 1 0 0 0 0

0 1 0 0 0 0 1 0

0 0 0 1 0 0 0 1

Sign = $(-1)^{8-4} = +1$

$$p(k) = \sum_{X \subseteq V} (-1)^{|V-X|} a(X)^k.$$

K-tuples of Colorclasses

$$p(k) = \sum_{X \subseteq V} (-1)^{|V-X|} a(X)^k.$$

- Will be counted $2^{\#\text{of uncolored vertices}}$ times, but equally many times with sign factor -1 as $+1$. Hence, they will cancel each other in the sum unless all vertices are colored.

Fast Zeta Transform

$$a(X) = \sum_{Y \subseteq X} [Y \text{ is a candidate colorclass}].$$

- [Yates 1937] A table containing $a(X)$ for all subsets X of V can be computed in $O(n2^n)$ time.

BHK'06 $\chi(G)$ -Algorithm

- Compute by Yates's algorithm

$$a(X) = \sum_{Y \subseteq X} [Y \text{ is a candidate colorclass}].$$

$2^n \text{poly}(n)$ time.

- For $k=1:n$, evaluate

$$p(k) = \sum_{X \subseteq V} (-1)^{|V-X|} a(X)^k.$$

$2^n \text{poly}(n)$ time.

until $p(k) \neq 0$, then return k .