

## Examination in Programming language theory

### Solutions

$$1. \frac{T_1}{\langle S, \sigma_0 \rangle \rightarrow \sigma_4} \quad \begin{array}{l} T_2 \quad \langle S, \sigma_4 \rangle \rightarrow \sigma_4 \\ \hline \langle S, \sigma_2 \rangle \rightarrow \sigma_4 \end{array} \quad \begin{array}{l} S_1 = s := s + y \text{ where} \\ S_2 = x := x - 1 \\ S_3 = S_1; S_2 \end{array}$$

$$T_1 = \frac{\langle S_1, \sigma_0 \rangle \rightarrow \sigma_1 \quad \langle S_2, \sigma_1 \rangle \rightarrow \sigma_2}{\langle S_3, \sigma_0 \rangle \rightarrow \sigma_2} \quad \begin{array}{l} \sigma_0 = [s \mapsto 0, x \mapsto 2, y \mapsto 5] \\ \sigma_1 = [s \mapsto 5, x \mapsto 2, y \mapsto 5] \\ \sigma_2 = [s \mapsto 5, x \mapsto 1, y \mapsto 5] \\ \sigma_3 = [s \mapsto 10, x \mapsto 1, y \mapsto 5] \\ \sigma_4 = [s \mapsto 10, x \mapsto 0, y \mapsto 5] \end{array}$$

$$T_2 = \frac{\langle S_1, \sigma_2 \rangle \rightarrow \sigma_3 \quad \langle S_2, \sigma_3 \rangle \rightarrow \sigma_4}{\langle S_3, \sigma_2 \rangle \rightarrow \sigma_4}$$

2.  $\langle \text{if } b \text{ then } S, \sigma \rangle \Rightarrow \langle S, \sigma \rangle$  if  $\mathcal{B}[b] = \text{tt}$ .  
 $\langle \text{if } b \text{ then } S, \sigma \rangle \Rightarrow \sigma$  if  $\mathcal{B}[b] = \text{ff}$ .
3. Let  $A = \lambda x. \lambda y. y(x \ x \ y)$ .  $A \ A \ F = (\lambda x. \lambda y. y(x \ x \ y)) \ A \ F = (\lambda y. y(A \ A \ y)) \ F = F(A \ A \ F)$ .
4. Assume that  $\mathcal{S}'_{cs}[S_i] \ c \ \sigma = c(\mathcal{S}_{ds}[S_i] \ \sigma)$  for all  $c, \sigma$  and  $i = 1, 2$ .  
 $\mathcal{S}'_{cs}[S_1; S_2] \ c \ \sigma = (\mathcal{S}'_{cs}[S_1] \circ \mathcal{S}'_{cs}[S_2]) \ c \ \sigma = (\mathcal{S}'_{cs}[S_1](\mathcal{S}'_{cs}[S_2] \ c)) \ \sigma = \mathcal{S}'_{cs}[S_1](\mathcal{S}'_{cs}[S_2] \ c) \ \sigma = (\mathcal{S}'_{cs}[S_2] \ c)(\mathcal{S}_{ds}[S_1] \ \sigma) = c(\mathcal{S}_{ds}[S_2](\mathcal{S}_{ds}[S_1] \ \sigma)) = c(\mathcal{S}_{ds}[S_1; S_2] \ \sigma)$
5. 

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data Formula = Formula Stm State State
data Rule = Rule Formula [Formula] (Bexp -> State -> Bool)
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6. There are four monotone functions.

| $x$ | $f_1(x)$ | $f_2(x)$ | $f_3(x)$ | $f_4(x)$ |
|-----|----------|----------|----------|----------|
| $a$ | $A$      | $A$      | $A$      | $B$      |
| $b$ | $A$      | $A$      | $B$      | $B$      |
| $c$ | $A$      | $B$      | $B$      | $B$      |

We show that  $f_1$  is continuous. Every subset of  $\{a, b, c\}$  is a chain. Thus there are 8 chains. For each chain  $C$  we must show that  $f_1(\bigcup C) = \bigcup \{f_1(x) \mid x \in C\}$ . In each case both members are equal to  $A$ .

(It is easy to show that if all chains are finite then monotonicity implies continuity.)