

# Explicit Locks

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# Synchronized

- Java incorporates a cross-platform threading model & memory model into language specification.
- Thread class
- Synchronized & Volatile
- Atomicity & Visibility

```
synchronized (lockObject) {  
    //update object state  
}
```

- So why mess with a good thing?

# Synchronized.. continued

- Can not Interrupt a Thread while Waiting to Acquire a Lock
- Might Need to Wait Forever to Acquire a Lock
- Must Release a Lock within Same Stack Frame where Acquired
- Lock Interface Provides More Extensive Locking Operations

# Package java.util.concurrent.locks

- ❑ Framework offering greater flexibility for locking and conditions from built in synchronization and monitors.
- ❑ Interfaces: Condition, Lock, ReadWriteLock
- ❑ Classes:
  - ❑ AbstractOwnableSynchronizer
  - ❑ AbstractQueuedLongSynchronizer
  - ❑ AbstractQueuedSynchronizer
  - ❑ LockSupport
  - ❑ ReentrantLock
  - ❑ ReentrantReadWriteLock
  - ❑ ReentrantReadWriteLock.ReadLock
  - ❑ ReentrantReadWriteLock.WriteLock

# Lock Interface

- Enables Access to Shared Resource by Multiple Threads
- Methods:
  - `void lock();` - Acquires the lock
  - `void lockInterruptibly();` - Acquires the lock unless the current thread is Interrupted.
  - `Condition newCondition();` - Returns a new Condition that is bound by this Lock instance.
  - `boolean tryLock();` - Acquires the lock only if it is free at the time of invocation.
  - `boolean tryLock(long time, TimeUnit unit);` - Returns true if the lock was acquired and false if waiting time expired before the lock was interrupted.
  - `void unlock();` - Releases the Lock

# Classes Implementing Lock Interface

- (1) `ReentrantLock`, (2) `ReentrantReadWriteLock.ReadLock`, (3) `ReentrantReadWriteLock.WriteLock`
- Reentrant Lock
  - Same behavior as the implicit monitor lock + some more
  - Lock owned by the thread with last successful locking and before unlocking

```
class X {  
    private final ReentrantLock lock = new ReentrantLock();  
    // ...  
  
    public void m() {  
        lock.lock(); // block until condition holds  
        try {  
            // ... method body  
        } finally {  
            lock.unlock();  
        }  
    }  
}
```

← What happens if an exception is thrown?

# ReentrantLock... continued

- Supports Fairness Policy – `public ReentrantLock(boolean fair)`
- Supports Interruptible Locks – `void lockInterruptibly()`
- Allows for Condition to be associated with this lock
- Provides Additional Methods for:
  - Queries:
    - Number of holds on this lock by the current thread
    - Whether current thread is waiting to acquire this lock
    - Whether any threads are waiting for the given condition associated with this lock
    - Whether lock is held by this thread
  - Returns a Collection of threads, the number of threads waiting for this lock (with or without the given Condition)

# ReentrantLock... continued

- Factors out the Objects monitor methods (wait, notify, notifyAll) into distinct objects.
- BoundedBuffer



```
class BoundedBuffer {
    final Lock lock = new ReentrantLock();
    final Condition notFull = lock.newCondition();
    final Condition notEmpty = lock.newCondition();

    final Object[] items = new Object[100];
    int putptr, takeptr, count;

    public void put(Object x) throws InterruptedException {
        lock.lock();
        try {
            while (count == items.length)
                notFull.await();
            items[putptr] = x;
            if (++putptr == items.length) putptr = 0;
            ++count;
            notEmpty.signal();
        } finally {
            lock.unlock();
        }
    }

    public Object take() throws InterruptedException {
        lock.lock();
        try {
            while (count == 0)
                notEmpty.await();
            Object x = items[takeptr];
            if (++takeptr == items.length) takeptr = 0;
            --count;
            notFull.signal();
            return x;
        } finally {
            lock.unlock();
        }
    }
}
```

# ReentrantReadWriteLock

- ❑ Supports Multiple Readers, but Only One Writer
- ❑ Implements ReadWriteLock Interface:
  - ❑ Lock readLock()
  - ❑ Lock writeLock()
- ❑ Encloses ReadWriteLock.ReadLock & ReadWriteLock.WriteLock classes that Implement Lock Interface
- ❑ Contains Similar Methods as ReentrantLock
- ❑ Condition Can Only Be Used with the Write Lock
- ❑ Writer can acquire a read lock, but not vice versa

# Example 1: Avoid Lock Ordering Deadlock

- Transfer money from an account A to an account B

- Using synchronized:

```
synchronized(fromAccount){  
    synchronized(toAccount){  
        //Money transfer logic  
    }  
}
```

Thread 1: Transfer from A to B

Thread 2: Transfer from B to A



- Using locked:

```
while (!expired){  
    if (fromAccount.lock.tryLock()){  
        try {  
            if(toAccount.lock.tryLock()){  
                try {  
                    //Money transfer logic...  
                    return true;  
                } finally {  
                    toAccount.lock.unlock();  
                }  
            }  
        } finally {  
            fromAccount.lock.unlock();  
        }  
    }  
    if(expired)  
        return false;  
    //Sleep a little to reduce chance of live locks  
}
```

Expired:  $currentTime \geq stopTime$

# Example 2: ReentrantReadWriteLock

- Try to Obtain Lock within Given Time Budget:

```
public boolean updateLDAPEntry(Account accountInfo, long time)
    throws InterruptedException{

    if (!lock.tryLock(time, TimeUnit.SECONDS))
        return false;
    try{
        return updateLDAP(accountInfo);
    }
    finally{
        lock.unlock();
    }
}
```

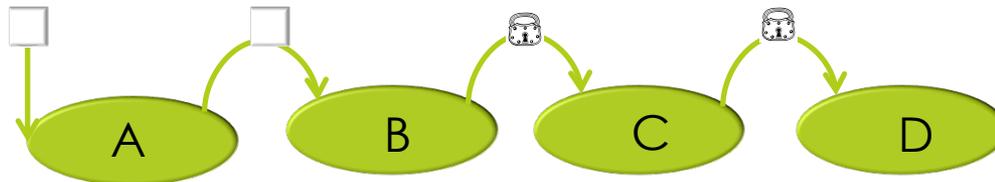
- Interruptible Locks

```
public boolean updateLDAPEntry(Account accountInfo)
    throws InterruptedException{
    lock.lockInterruptibly();
    try{
        return cancellableUpdateLDAP(accountInfo);
    }
    finally{
        lock.unlock();
    }
}
```

```
private boolean cancellableUpdateLDAP(Account accountInfo)
    throws InterruptedException{
    //implement cancellable update
    return true; //or false
}
```

# Hand-over-locking

- Intrinsic Locks Block Structured
- Reducing Lock Granularity can Enhance Scalability
- Lock interface Allows for Locks to be Acquired and Released in Different Scopes & Multiple Locks to be Acquired and Released



# Performance Considerations

- Resources Expended on Lock Management & Scheduling
- Java 5.0 (Initial Locks framework released)

Figure 1. Throughput for synchronization and Lock, single CPU

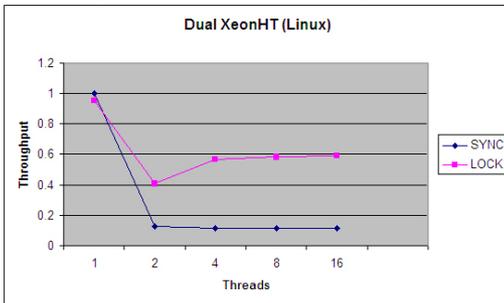
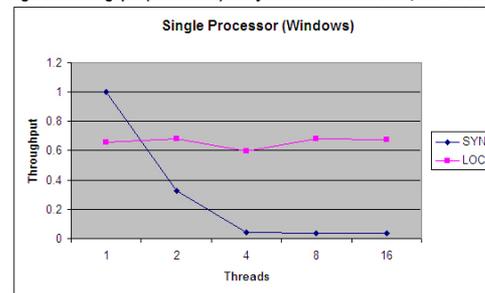


Figure 2. Throughput (normalized) for synchronization and Lock, four CPUs

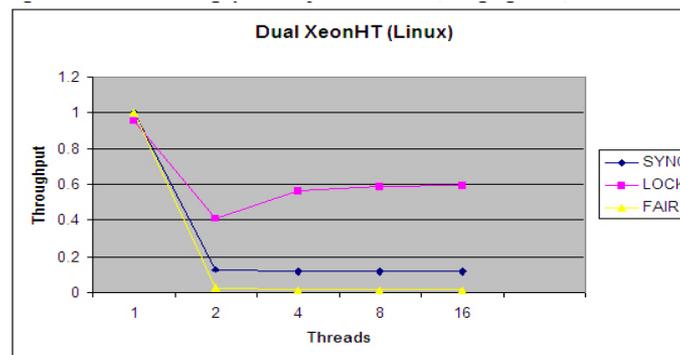


- Java 6 – Intrinsic and Explicit Scale Fairly Equally.
- Performance – Moving Target



# Fairness

- ❑ Fair vs. NonFair Locks
- ❑ Performance cost!
- ❑ High Load Can Hinder Time of Thread Resuming Time vs. its Actual Run Time.
- ❑ Long Wait Times or Mean Time Between Lock Requests.
- ❑ Java 5:



# Intrinsic (Synchronized) vs. Explicit Locks?

Feature	Intrinsic	Explicit
Timed Lock Wait	x	✓
Interruptible Lock Wait	x	✓
Fairness	x	✓
Non-block structure locking	x	✓
Familiar syntax, used extensively	✓	x
Good idea to mix the two	N	O
More dangerous	x	✓
Bright Future Awaiting ☺	✓	x

Far, Far Away,  
In the Galaxy of Java 10  
the Anticipated Performance of  
Intrinsic over Explicit Lock  
Will Be:

